

# Analytic Proof Systems for Shramko-Wansing logic **SIXTEEN**<sub>2,des</sub>

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July 02, 2024

## Abstract

We give sequent calculus, analytic tableaux, natural deduction, and clause translation systems for resolution for Shramko-Wansing logic **SIXTEEN**<sub>2,des</sub>.

## 1 Introduction

In this paper we present calculi for Shramko-Wansing logic **SIXTEEN**<sub>2,des</sub> [12]. **SIXTEEN**<sub>2,des</sub> has sixteen truth values **N**, **N**, **F**, **T**, **B**, **NF**, **NT**, **FT**, **NB**, **FB**, **TB**, **NFT**, **NFB**, **NTB**, **FTB**, **A** (with **T**, **NT**, **TB**, **NTB** designated), and connectives  $\neg_t$ ,  $\neg_f$ ,  $\wedge_t$ ,  $\vee_t$ ,  $\wedge_f$ ,  $\vee_f$ . Its syntax and semantics is detailed in section 2.

We first present a 16-sided sequent calculus in section 3. The fundamental idea for many-sided sequent calculi for finite-valued logics goes back to Schröter [11], Rousseau [8], Takahashi [14]. We follow the method given by Baaz, Fermüller, and Zach [4] and Zach [15] for constructing inference rules. This guarantees that our system automatically has soundness and completeness theorems, cut-elimination theorem and Maehara lemma (interpolation). For proofs of these results see [4, 15].

Signed tableau systems for finite-valued logics were proposed by Surma [13] and Carnielli [6], and generalized by Hähnle [7]. In section 4, we present a signed tableau system for Shramko-Wansing logic.

Many-valued natural deduction systems for finite-valued logics have been investigated by Baaz, Fermüller, and Zach [3] and Zach [15]. We give the introduction and elimination rules for the natural deduction system for **SIXTEEN**<sub>2,des</sub> in section 5.

In addition to Hähnle's work on tableaux-based theorem proving for finite-valued logic, Baaz and Fermüller [1] have studied resolution calculi for clauses

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\*Cite as: Multlog (2024), “Analytic proof systems for Shramko-Wansing logic **SIXTEEN**<sub>2,des</sub>.” [PDF generated by MULTLOG, v. 1.16a, <https://logic.at/multlog>]. Available at <https://logic.at/multlog/shramko-wansing.pdf>.

See appendices A and B for the specification of Shramko-Wansing logic.

of signed literals. In order for these calculi to be used to prove that formulas of Shramko-Wansing logic are valid or follow from some others, it is necessary to produce sets of signed clauses. In section 6, we present a translation calculus that yields a set of clauses from a set of formulas.

The rules we provide are optimal in each case, and use the algorithms developed by Salzer [9, 10].

## 2 Syntax and semantics

**Definition 1.** The propositional language  $\mathcal{L}$  for Shramko-Wansing logic consists of

1. propositional variables:  $x_0, x_1, x_2, \dots$
2. propositional connectives, arity given in parenthesis:  $\neg_t$  (1),  $\neg_f$  (1),  $\wedge_t$  (2),  $\vee_t$  (2),  $\wedge_f$  (2), and  $\vee_f$  (2)
3. auxiliary symbols: “(”, “)” and “,”

*Formulas* are defined inductively:

1. Every propositional variable is a formula.
2. If  $A$  is a formula, so is  $\neg_t A$ .
3. If  $A$  is a formula, so is  $\neg_f A$ .
4. If  $A$  and  $B$  are formulas, so is  $(A \wedge_t B)$ .
5. If  $A$  and  $B$  are formulas, so is  $(A \vee_t B)$ .
6. If  $A$  and  $B$  are formulas, so is  $(A \wedge_f B)$ .
7. If  $A$  and  $B$  are formulas, so is  $(A \vee_f B)$ .

As a notational convention, lowercase letters will be used to denote variables, possibly indexed. Uppercase letters  $A, B, C, \dots$  will stand for formulas, greek letters  $\Gamma, \Delta, \Lambda, \dots$  for sequences and sets of formulas. The symbol  $\square$  stands for general propositional connectives.

The connectives  $\wedge_t$  and  $\vee_t$  of **SIXTEEN**<sub>2,des</sub> are defined as the inf and sup of the  $\leq_t$  ordering given in fig. 1. The connectives  $\wedge_f$  and  $\vee_f$  correspond to inf and sup of  $\leq_f$ . For simplicity, we leave out the  $\leq_i$  ordering and other operators and inversions defined in [12].

**Definition 2.** The *matrix* for Shramko-Wansing logic is given by:

1. the set of *truth values*  $V = \{\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}\}$ ,
2. the set  $V^+ = \{\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}\} \subseteq V$  of *designated truth values*,

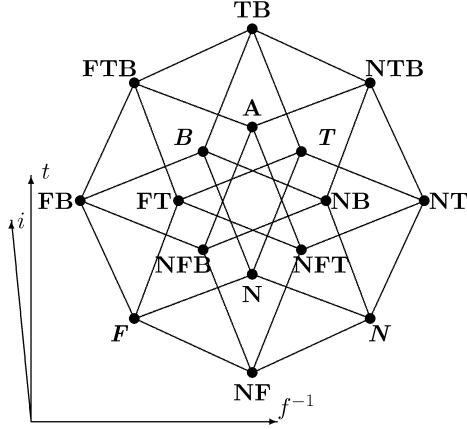


Figure 1: The  $\leq_t$  and  $\leq_f$  orderings

3. the truth functions for connectives  $\neg_t$ ,  $\neg_f$ ,  $\wedge_t$ ,  $\vee_t$ ,  $\wedge_f$  and  $\vee_f$ , as given below;

The set of *undesignated values* is  $V^- = V \setminus V^+ = \{\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}\}$ .

The truth functions for connectives  $\neg_t$ ,  $\neg_f$ ,  $\wedge_t$ ,  $\vee_t$ ,  $\wedge_f$  and  $\vee_f$  are defined by

$\widetilde{\neg}_t$		$\widetilde{\neg}_f$	
$\mathbf{N}$	$\mathbf{N}$	$\mathbf{N}$	$\mathbf{N}$
$\mathbf{N}$	$\mathbf{T}$	$\mathbf{N}$	$\mathbf{F}$
$\mathbf{F}$	$\mathbf{B}$	$\mathbf{F}$	$\mathbf{N}$
$\mathbf{T}$	$\mathbf{N}$	$\mathbf{T}$	$\mathbf{B}$
$\mathbf{B}$	$\mathbf{F}$	$\mathbf{B}$	$\mathbf{T}$
$\mathbf{NF}$	$\mathbf{TB}$	$\mathbf{NF}$	$\mathbf{NF}$
$\mathbf{NT}$	$\mathbf{NT}$	$\mathbf{NT}$	$\mathbf{FB}$
$\mathbf{FT}$	$\mathbf{NB}$	$\mathbf{FT}$	$\mathbf{NB}$
$\mathbf{NB}$	$\mathbf{FT}$	$\mathbf{NB}$	$\mathbf{FT}$
$\mathbf{FB}$	$\mathbf{FB}$	$\mathbf{FB}$	$\mathbf{NT}$
$\mathbf{TB}$	$\mathbf{NF}$	$\mathbf{TB}$	$\mathbf{TB}$
$\mathbf{NFT}$	$\mathbf{NTB}$	$\mathbf{NFT}$	$\mathbf{NFB}$
$\mathbf{NFB}$	$\mathbf{FTB}$	$\mathbf{NFB}$	$\mathbf{NFT}$
$\mathbf{NTB}$	$\mathbf{NFT}$	$\mathbf{NTB}$	$\mathbf{FTB}$
$\mathbf{FTB}$	$\mathbf{NFB}$	$\mathbf{FTB}$	$\mathbf{NTB}$
$\mathbf{A}$	$\mathbf{A}$	$\mathbf{A}$	$\mathbf{A}$

$\widetilde{\wedge}_t$	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
N	N	N	F	N	N	NF	N	F	N	F	N	NF	NF	N	F	NF
N	N	N	NF	N	N	NF	N	NF	N	NF	N	NF	NF	N	NF	NF
F	F	NF	F	F	F	NF	NF	F	NF	F	F	NF	NF	NF	F	NF
T	N	N	F	T	N	NF	NT	FT	N	F	T	NFT	NF	NT	FT	NFT
B	N	N	F	N	B	NF	N	F	NB	FB	B	NF	NFB	NB	FB	NFB
NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF
NT	N	N	NF	NT	N	NF	NT	NFT	N	NF	NT	NFT	NF	NT	NFT	NFT
FT	F	NF	F	FT	F	NF	NFT	FT	NF	F	FT	NFT	NF	NFT	FT	NFT
NB	N	N	NF	N	NB	NF	N	NF	NB	NFB	NB	NF	NFB	NB	NFB	NFB
FB	F	NF	F	F	FB	NF	NF	F	NFB	FB	FB	NF	NFB	NFB	FB	NFB
TB	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
NFT	NF	NF	NF	NFT	NF	NF	NFT	NFT	NF	NF	NFT	NFT	NF	NFT	NFT	NFT
NFB	NF	NF	NF	NFB	NF	NF	NF	NFB								
NTB	N	N	NF	NT	NB	NF	NT	NFT	NB	NFB	NTB	NFT	NFB	NTB	A	A
FTB	F	NF	F	FT	FB	NF	NFT	FT	NFB	FB	FTB	NFT	NFB	A	FTB	A
A	NF	NF	NF	NFT	NFB	NF	NFT	NFT	NFB	NFB	A	NFT	NFB	A	A	A
$\widetilde{\vee}_t$	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
N	N	N	N	T	B	N	T	T	B	B	TB	T	B	TB	TB	TB
N	N	N	N	T	B	N	NT	T	NB	B	TB	NT	NB	NTB	TB	NTB
F	N	N	F	T	B	F	T	FT	B	FB	TB	FT	FB	TB	FTB	FTB
T	T	T	T	TB	T	T	T	TB	TB	TB	TB	T	TB	TB	TB	TB
B	B	B	B	TB	B	B	TB	TB	B	TB	TB	TB	B	TB	TB	TB
NF	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
NT	T	NT	T	T	TB	NT	NT	T	NTB	TB	TB	NT	NTB	NTB	TB	NTB
FT	T	T	FT	T	TB	FT	T	FT	TB	FTB	TB	FT	FTB	TB	FTB	FTB
NB	B	NB	B	TB	B	NB	NTB	TB	NB	B	TB	NTB	NB	NTB	TB	NTB
FB	B	B	FB	TB	B	FB	TB	FTB	B	FB	TB	FTB	FB	TB	FTB	FTB
TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB
NFT	T	NT	FT	T	TB	NFT	NT	FT	NTB	FTB	TB	NFT	A	NTB	FTB	A
NFB	B	NB	FB	TB	B	NFB	NTB	FTB	NB	FB	TB	A	NFB	NTB	FTB	A
NTB	TB	NTB	TB	TB	TB	NTB	NTB	TB	NTB	TB	TB	NTB	NTB	NTB	TB	NTB
FTB	TB	TB	FTB	TB	TB	FTB	TB	FTB	TB	FTB	TB	FTB	FTB	TB	FTB	FTB
A	TB	NTB	FTB	TB	TB	A	NTB	FTB	NTB	FTB	TB	A	A	NTB	FTB	A

$\widetilde{\wedge}_f$	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
N	N	N	N	T	N	N	NT	T	N	N	T	NT	N	NT	T	NT
N	N	N	N	NT	N	N	NT	NT	N	N	NT	NT	N	NT	NT	NT
F	N	N	F	T	N	NF	NT	FT	N	F	T	NFT	NF	NT	FT	NFT
T	T	NT	T	T	T	NT	NT	T	NT	T	T	NT	NT	NT	T	NT
B	N	N	N	T	B	N	NT	T	NB	B	TB	NT	NB	NTB	TB	NTB
NF	N	N	NF	NT	N	NF	NT	NFT	N	NF	NT	NFT	NF	NT	NFT	NFT
NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT
FT	T	NT	FT	T	T	NFT	NT	FT	NT	FT	T	NFT	NFT	NT	FT	NFT
NB	N	N	N	NT	NB	N	NT	NT	NB	NB	NTB	NT	NB	NTB	NTB	NTB
FB	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
TB	T	NT	T	T	TB	NT	NT	T	NTB	TB	TB	NT	NTB	NTB	TB	NTB
NFT	NT	NT	NFT	NT	NT	NFT	NT	NFT	NT	NFT	NT	NFT	NFT	NT	NFT	NFT
NFB	N	N	NF	NT	NB	NF	NT	NFT	NB	NFB	NTB	NFT	NFB	NTB	A	A
NTB	NT	NT	NT	NT	NTB	NT	NT	NT	NTB	NTB	NTB	NT	NTB	NTB	NTB	NTB
FTB	T	NT	FT	T	TB	NFT	NT	FT	NTB	FTB	TB	NFT	A	NTB	FTB	A
A	NT	NT	NFT	NT	NTB	NFT	NT	NFT	NTB	A	NTB	NFT	A	NTB	A	A
$\widetilde{\vee}_f$	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
N	N	N	F	N	B	F	N	F	B	FB	B	F	FB	B	FB	FB
N	N	N	F	N	B	NF	N	F	NB	FB	B	NF	NFB	NB	FB	NFB
F	F	F	F	FB	F	F	F	F	FB	FB	FB	F	FB	FB	FB	FB
T	N	N	F	T	B	F	T	FT	B	FB	TB	FT	FB	TB	FTB	FTB
B	B	B	FB	B	B	FB	B	FB	B	FB	B	FB	FB	B	FB	FB
NF	F	NF	F	F	FB	NF	NF	F	NFB	FB	FB	NF	NFB	NFB	FB	NFB
NT	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
FT	F	F	F	FT	FB	F	FT	FT	FB	FB	FTB	FT	FB	FTB	FTB	FTB
NB	B	NB	FB	B	B	NFB	NB	FB	NB	FB	B	NFB	NFB	NB	FB	NFB
FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB
TB	B	B	FB	TB	B	FB	TB	FTB	B	FB	TB	FTB	FB	TB	FTB	FTB
NFT	F	NF	F	FT	FB	NF	NFT	FT	NFB	FB	FTB	NFT	NFB	A	FTB	A
NFB	FB	NFB	FB	FB	NFB	NFB	NFB	FB	NFB	FB	NFB	NFB	NFB	FB	NFB	NFB
NTB	B	NB	FB	TB	B	NFB	NTB	FTB	NB	FB	TB	A	NFB	NTB	FTB	A
FTB	FB	FB	FB	FTB	FB	FB	FTB	FTB	FB	FB	FTB	FTB	FB	FTB	FTB	FTB
A	FB	NFB	FB	FTB	FB	NFB	A	FTB	NFB	FB	FTB	A	NFB	A	FTB	A

**Definition 3.** Let  $A$  be a formula and  $x_0, x_1, \dots, x_k$  the variables occurring in  $A$ . Then an *interpretation*  $\mathcal{I}$  of  $A$  is an assignment of truth values to the variables.

**Definition 4.** Given an interpretation  $\mathcal{I}$ , we define the *valuation*  $\text{val}_{\mathcal{I}}$  for formulas  $A$  to truth values as follows:

1. If  $A$  is atomic, then  $\text{val}_{\mathcal{I}}(A)$  simply is the interpretation of  $A$ .
2. If  $A = \square(A_1, \dots, A_n)$ , where  $A_1, \dots, A_n$  are formulas, and  $\widetilde{\square}$  is the associated truth function to  $\square$ , then  $\text{val}_{\mathcal{I}}(A) = \widetilde{\square}(\text{val}_{\mathcal{I}}(A_1), \dots, \text{val}_{\mathcal{I}}(A_n))$ .

**Definition 5.** An interpretation  $\mathcal{I}$  satisfies a formula  $A$ , in symbols:  $\mathcal{I} \models A$ , iff  $\text{val}_{\mathcal{I}}(A) \in V^+$ .

**Definition 6.**  $\Delta$  entails  $A$  iff  $\mathfrak{I} \models A$  for every interpretation  $\mathfrak{I}$  such that  $\mathfrak{I} \models B$  for all  $B \in \Delta$ .  $A$  is a *tautology* iff it is satisfied by every interpretation  $\mathfrak{I}$ .

### 3 Sequent calculus for Shramko-Wansing logic

**Definition 7** (Syntax of Sequents). A *sequent*  $\Gamma$  is a 16-tuple

$$\Gamma_N \mid \Gamma_N \mid \Gamma_F \mid \Gamma_T \mid \Gamma_B \mid \Gamma_{NF} \mid \Gamma_{NT} \mid \Gamma_{FT} \mid \Gamma_{NB} \mid \Gamma_{FB} \mid \Gamma_{TB} \mid \Gamma_{NFT} \mid \Gamma_{NFB} \mid \Gamma_{NTB} \mid \Gamma_{FTB} \mid \Gamma_A$$

of finite sequences  $\Gamma_v$  of formulas, where  $v \in V$ . The  $\Gamma_v$  are called the *components* of  $\Gamma$ .

For a sequence of formulas  $\Delta$ , and  $W \subseteq V$ , let  $[W: \Delta]$  denote the sequent whose component  $\Gamma_v$  is  $\Delta$  if  $v \in W$  and empty otherwise. For  $[\{w_1, \dots, w_k\}: \Delta]$  we also write  $[w_1, \dots, w_k: \Delta]$ . If  $\Gamma$  and  $\Gamma'$  are sequents, then  $\Gamma, \Gamma'$  denotes the component-wise union, i.e., the  $v$ -component of  $\Gamma, \Gamma'$  is  $\Gamma_v, \Gamma'_v$ .

**Definition 8.** Let  $\mathfrak{I}$  be an interpretation.  $\mathfrak{I}$  satisfies a sequent  $\Gamma$  iff there is a  $v \in V$  so that for some formula  $A \in \Gamma_v$ ,  $\text{val}_{\mathfrak{I}}(F) = v$ .  $\mathfrak{I}$  is called a *model* of  $\Gamma$ , in symbols  $\mathfrak{I} \models \Gamma$ .

$\Gamma$  is called *satisfiable* iff there is an interpretation  $\mathfrak{I}$  so that  $\mathfrak{I} \models \Gamma$  and *valid* iff for every interpretation  $\mathfrak{I}$ ,  $\mathfrak{I} \models \Gamma$ .

**Proposition 9.**  $\Delta \models A$  iff the sequent  $[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \Delta], [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A]$  is valid.

**Definition 10.** The *sequent calculus* for Shramko-Wansing logic is given by:

1. axiom schemas of the form  $[V: A]$ ,
2. weakening rules for every truth value  $v$ :

$$\frac{\Gamma}{\Gamma, [v: A]} \text{ W: } v$$

3. exchange rules for every truth value  $v$ :

$$\frac{\Gamma, [v: A, B], \Delta}{\Gamma, [v: B, A], \Delta} \text{ X: } v$$

4. contraction rules for every truth value  $v$ :

$$\frac{\Gamma, [v: A, A]}{\Gamma, [v: A]} \text{ C: } v$$

5. cut rules for every two truth values  $v \neq w$ :

$$\frac{\Gamma, [v: A] \quad \Delta, [w: A]}{\Gamma, \Delta} \text{ CUT: } vw$$

6. an introduction rule  $\square : v$  for every connective  $\square$  and every truth value  $v$ , as specified below.

(2)–(5) are called *structural rules*. (6) are called *logical rules*.

The introduction rules for connective  $\neg_t$  are given by

$$\begin{array}{c}
 \frac{\Gamma, [\mathbf{N}: A]}{\Gamma, [\mathbf{N}: \neg_t A]} \neg_t:\mathbf{N} \quad \frac{\Gamma, [\mathbf{T}: A]}{\Gamma, [\mathbf{N}: \neg_t A]} \neg_t:\mathbf{N} \quad \frac{\Gamma, [\mathbf{B}: A]}{\Gamma, [\mathbf{F}: \neg_t A]} \neg_t:\mathbf{F} \\
 \frac{\Gamma, [\mathbf{N}: A]}{\Gamma, [\mathbf{T}: \neg_t A]} \neg_t:\mathbf{T} \quad \frac{\Gamma, [\mathbf{F}: A]}{\Gamma, [\mathbf{B}: \neg_t A]} \neg_t:\mathbf{B} \quad \frac{\Gamma, [\mathbf{TB}: A]}{\Gamma, [\mathbf{NF}: \neg_t A]} \neg_t:\mathbf{NF} \\
 \frac{\Gamma, [\mathbf{NT}: A]}{\Gamma, [\mathbf{NT}: \neg_t A]} \neg_t:\mathbf{NT} \quad \frac{\Gamma, [\mathbf{NB}: A]}{\Gamma, [\mathbf{FT}: \neg_t A]} \neg_t:\mathbf{FT} \quad \frac{\Gamma, [\mathbf{FT}: A]}{\Gamma, [\mathbf{NB}: \neg_t A]} \neg_t:\mathbf{NB} \\
 \frac{\Gamma, [\mathbf{FB}: A]}{\Gamma, [\mathbf{FB}: \neg_t A]} \neg_t:\mathbf{FB} \quad \frac{\Gamma, [\mathbf{NF}: A]}{\Gamma, [\mathbf{TB}: \neg_t A]} \neg_t:\mathbf{TB} \quad \frac{\Gamma, [\mathbf{NTB}: A]}{\Gamma, [\mathbf{NFT}: \neg_t A]} \neg_t:\mathbf{NFT} \\
 \frac{\Gamma, [\mathbf{FTB}: A]}{\Gamma, [\mathbf{NFB}: \neg_t A]} \neg_t:\mathbf{NFB} \quad \frac{\Gamma, [\mathbf{NFT}: A]}{\Gamma, [\mathbf{NTB}: \neg_t A]} \neg_t:\mathbf{NTB} \\
 \frac{\Gamma, [\mathbf{NFB}: A]}{\Gamma, [\mathbf{FTB}: \neg_t A]} \neg_t:\mathbf{FTB} \quad \frac{\Gamma, [\mathbf{A}: A]}{\Gamma, [\mathbf{A}: \neg_t A]} \neg_t:\mathbf{A}
 \end{array}$$

The introduction rules for connective  $\neg_f$  are given by

$$\begin{array}{c}
 \frac{\Gamma, [\mathbf{N}: A]}{\Gamma, [\mathbf{N}: \neg_f A]} \neg_f:\mathbf{N} \quad \frac{\Gamma, [\mathbf{F}: A]}{\Gamma, [\mathbf{N}: \neg_f A]} \neg_f:\mathbf{N} \quad \frac{\Gamma, [\mathbf{N}: A]}{\Gamma, [\mathbf{F}: \neg_f A]} \neg_f:\mathbf{F} \\
 \frac{\Gamma, [\mathbf{B}: A]}{\Gamma, [\mathbf{T}: \neg_f A]} \neg_f:\mathbf{T} \quad \frac{\Gamma, [\mathbf{T}: A]}{\Gamma, [\mathbf{B}: \neg_f A]} \neg_f:\mathbf{B} \quad \frac{\Gamma, [\mathbf{NF}: A]}{\Gamma, [\mathbf{NF}: \neg_f A]} \neg_f:\mathbf{NF} \\
 \frac{\Gamma, [\mathbf{FB}: A]}{\Gamma, [\mathbf{NT}: \neg_f A]} \neg_f:\mathbf{NT} \quad \frac{\Gamma, [\mathbf{NB}: A]}{\Gamma, [\mathbf{FT}: \neg_f A]} \neg_f:\mathbf{FT} \quad \frac{\Gamma, [\mathbf{FT}: A]}{\Gamma, [\mathbf{NB}: \neg_f A]} \neg_f:\mathbf{NB} \\
 \frac{\Gamma, [\mathbf{NT}: A]}{\Gamma, [\mathbf{FB}: \neg_f A]} \neg_f:\mathbf{FB} \quad \frac{\Gamma, [\mathbf{TB}: A]}{\Gamma, [\mathbf{TB}: \neg_f A]} \neg_f:\mathbf{TB} \quad \frac{\Gamma, [\mathbf{NFB}: A]}{\Gamma, [\mathbf{NFT}: \neg_f A]} \neg_f:\mathbf{NFT} \\
 \frac{\Gamma, [\mathbf{NFT}: A]}{\Gamma, [\mathbf{NFB}: \neg_f A]} \neg_f:\mathbf{NFB} \quad \frac{\Gamma, [\mathbf{FTB}: A]}{\Gamma, [\mathbf{NTB}: \neg_f A]} \neg_f:\mathbf{NTB} \\
 \frac{\Gamma, [\mathbf{NTB}: A]}{\Gamma, [\mathbf{FTB}: \neg_f A]} \neg_f:\mathbf{FTB} \quad \frac{\Gamma, [\mathbf{A}: A]}{\Gamma, [\mathbf{A}: \neg_f A]} \neg_f:\mathbf{A}
 \end{array}$$

The introduction rules for connective  $\wedge_t$  are given by

$$\frac{\Gamma, [\mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{TB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{T}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{B}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{TB}: A]}{\Gamma, [\mathbf{N}: A \wedge_t B]} \wedge_t:\mathbf{N}$$

$$\begin{array}{c}
\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{NT}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{NT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NT}, \mathbf{NB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{NT}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: A]}{\Gamma, [\mathbf{N}: A \wedge_t B]} \wedge_t: \mathbf{N} \\[10pt]
\frac{\Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FB}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{F}: A \wedge_t B]} \wedge_t: \mathbf{F} \\[10pt]
\frac{\Gamma, [\mathbf{T}, \mathbf{TB}: B] \quad \Gamma, [\mathbf{T}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{TB}: A]}{\Gamma, [\mathbf{T}: A \wedge_t B]} \wedge_t: \mathbf{T} \\[10pt]
\frac{\Gamma, [\mathbf{B}, \mathbf{TB}: B] \quad \Gamma, [\mathbf{B}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{TB}: A]}{\Gamma, [\mathbf{B}: A \wedge_t B]} \wedge_t: \mathbf{B} \\[10pt]
\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B]}{\Gamma, [\mathbf{NF}: A \wedge_t B]} \wedge_t: \mathbf{NF} \\[10pt]
\frac{\Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}: A, B] \quad \Gamma, [\mathbf{NT}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A]}{\Gamma, [\mathbf{NT}: A \wedge_t B]} \wedge_t: \mathbf{NT} \\[10pt]
\frac{\Gamma, [\mathbf{T}, \mathbf{FT}, \mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{T}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{FT}, \mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{FT}: A \wedge_t B]} \wedge_t: \mathbf{FT} \\[10pt]
\frac{\Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: A]}{\Gamma, [\mathbf{NB}: A \wedge_t B]} \wedge_t: \mathbf{NB} \\[10pt]
\frac{\Gamma, [\mathbf{B}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{B}, \mathbf{FB}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{FB}: A \wedge_t B]} \wedge_t: \mathbf{FB} \\[10pt]
\frac{\Gamma, [\mathbf{TB}: B] \quad \Gamma, [\mathbf{TB}: A]}{\Gamma, [\mathbf{TB}: A \wedge_t B]} \wedge_t: \mathbf{TB} \\[10pt]
\frac{\Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{TB}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{NT}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{NFT}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{TB}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{NFT}: A \wedge_t B]} \wedge_t: \mathbf{NFT} \\[10pt]
\frac{\Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{NFB}: A \wedge_t B]} \wedge_t: \mathbf{NFB} \\[10pt]
\frac{\Gamma, [\mathbf{TB}, \mathbf{NTB}: B] \quad \Gamma, [\mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{TB}, \mathbf{NTB}: A]}{\Gamma, [\mathbf{NTB}: A \wedge_t B]} \wedge_t: \mathbf{NTB} \\[10pt]
\frac{\Gamma, [\mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{FTB}: A \wedge_t B]} \wedge_t: \mathbf{FTB} \\[10pt]
\frac{\Gamma, [\mathbf{TB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: B] \quad \Gamma, [\mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{TB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{A}: A \wedge_t B]} \wedge_t: \mathbf{A}
\end{array}$$

The introduction rules for connective  $\vee_t$  are given by

$$\begin{array}{c}
\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}: A]}{\Gamma, [\mathbf{N}: A \vee_t B]} \vee_t: \mathbf{N} \\[10pt]
\frac{\Gamma, [\mathbf{N}, \mathbf{NF}: B] \quad \Gamma, [\mathbf{N}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}: A]}{\Gamma, [\mathbf{N}: A \vee_t B]} \vee_t: \mathbf{N}
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma, [F, \mathbf{NF}: B] \quad \Gamma, [F: A, B] \quad \Gamma, [F, \mathbf{NF}: A]}{\Gamma, [F: A \vee_t B]} \vee_t: \mathbf{F} \\
\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{NT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A]}{\Gamma, [T: A \vee_t B]} \vee_t: \mathbf{T} \\
\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FB}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [B: A \vee_t B]} \vee_t: \mathbf{B} \\
\frac{\Gamma, [\mathbf{NF}: B] \quad \Gamma, [\mathbf{NF}: A]}{\Gamma, [\mathbf{NF}: A \vee_t B]} \vee_t: \mathbf{NF} \\
\frac{\Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NFT}: B] \quad \Gamma, [\mathbf{N}, \mathbf{NT}: A, B] \quad \Gamma, [\mathbf{NT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NFT}: A]}{\Gamma, [\mathbf{NT}: A \vee_t B]} \vee_t: \mathbf{NT} \\
\frac{\Gamma, [F, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: B] \quad \Gamma, [F, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [F, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: A]}{\Gamma, [\mathbf{FT}: A \vee_t B]} \vee_t: \mathbf{FT} \\
\frac{\Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NB}: A \vee_t B]} \vee_t: \mathbf{NB} \\
\frac{\Gamma, [F, \mathbf{NF}, \mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [F, \mathbf{FB}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{NFB}: A, B] \quad \Gamma, [F, \mathbf{NF}, \mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{FB}: A \vee_t B]} \vee_t: \mathbf{FB} \\
\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{NT}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{TB}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B]}{\Gamma, [\mathbf{TB}: A \vee_t B]} \vee_t: \mathbf{TB} \\
\frac{\Gamma, [\mathbf{NF}, \mathbf{NFT}: B] \quad \Gamma, [\mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{NF}, \mathbf{NFT}: A]}{\Gamma, [\mathbf{NFT}: A \vee_t B]} \vee_t: \mathbf{NFT} \\
\frac{\Gamma, [\mathbf{NF}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{NF}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NFB}: A \vee_t B]} \vee_t: \mathbf{NFB} \\
\frac{\Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: B] \quad \Gamma, [\mathbf{N}, \mathbf{NT}, \mathbf{NB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{NT}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{NTB}: A \vee_t B]} \vee_t: \mathbf{NTB} \\
\frac{\Gamma, [F, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: B] \quad \Gamma, [F, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{NFT}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [F, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{FTB}: A \vee_t B]} \vee_t: \mathbf{FTB} \\
\frac{\Gamma, [\mathbf{NF}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: B] \quad \Gamma, [\mathbf{NFT}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{NFB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{NF}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: A]}{\Gamma, [\mathbf{A}: A \vee_t B]} \vee_t: \mathbf{A}
\end{array}$$

The introduction rules for connective  $\wedge_f$  are given by

$$\begin{array}{c}
\frac{\Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{F}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{B}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FB}: A]}{\Gamma, [\mathbf{N}: A \wedge_f B]} \wedge_f: \mathbf{N} \\
\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{N}: A \wedge_f B]} \wedge_f: \mathbf{N} \\
\frac{\Gamma, [F, \mathbf{FB}: B] \quad \Gamma, [F: A, B] \quad \Gamma, [F, \mathbf{FB}: A]}{\Gamma, [F: A \wedge_f B]} \wedge_f: \mathbf{F}
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{TB}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{FT}, \mathbf{TB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{T}: A \wedge_f B]} \wedge_f : \mathbf{T} \\[10pt]
\frac{\Gamma, [\mathbf{B}, \mathbf{FB}: B] \quad \Gamma, [\mathbf{B}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{FB}: A]}{\Gamma, [\mathbf{B}: A \wedge_f B]} \wedge_f : \mathbf{B} \\[10pt]
\frac{\Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}: A, B] \quad \Gamma, [\mathbf{NF}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NF}: A \wedge_f B]} \wedge_f : \mathbf{NF} \\[10pt]
\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{NT}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{TB}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B]}{\Gamma, [\mathbf{NT}: A \wedge_f B]} \wedge_f : \mathbf{NT} \\[10pt]
\frac{\Gamma, [\mathbf{F}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{F}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{FT}: A \wedge_f B]} \wedge_f : \mathbf{FT} \\[10pt]
\frac{\Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NB}: A \wedge_f B]} \wedge_f : \mathbf{NB} \\[10pt]
\frac{\Gamma, [\mathbf{FB}: B] \quad \Gamma, [\mathbf{FB}: A]}{\Gamma, [\mathbf{FB}: A \wedge_f B]} \wedge_f : \mathbf{FB} \\[10pt]
\frac{\Gamma, [\mathbf{B}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{B}, \mathbf{TB}: A, B] \quad \Gamma, [\mathbf{TB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{TB}: A \wedge_f B]} \wedge_f : \mathbf{TB} \\[10pt]
\frac{\Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{NF}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{NFT}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{NFT}: A \wedge_f B]} \wedge_f : \mathbf{NFT} \\[10pt]
\frac{\Gamma, [\mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NFB}: A \wedge_f B]} \wedge_f : \mathbf{NFB} \\[10pt]
\frac{\Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{TB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{NTB}: A \wedge_f B]} \wedge_f : \mathbf{NTB} \\[10pt]
\frac{\Gamma, [\mathbf{FB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{FTB}: A \wedge_f B]} \wedge_f : \mathbf{FTB} \\[10pt]
\frac{\Gamma, [\mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: B] \quad \Gamma, [\mathbf{NFB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{A}: A \wedge_f B]} \wedge_f : \mathbf{A}
\end{array}$$

The introduction rules for connective  $\vee_f$  are given by

$$\begin{array}{c}
\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{NT}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{T}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{NT}: A]}{\Gamma, [\mathbf{N}: A \vee_f B]} \vee_f : \mathbf{N} \\[10pt]
\frac{\Gamma, [\mathbf{N}, \mathbf{NT}: B] \quad \Gamma, [\mathbf{N}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NT}: A]}{\Gamma, [\mathbf{N}: A \vee_f B]} \vee_f : \mathbf{N} \\[10pt]
\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A]}{\Gamma, [\mathbf{F}: A \vee_f B]} \vee_f : \mathbf{F} \\[10pt]
\frac{\Gamma, [\mathbf{T}, \mathbf{NT}: B] \quad \Gamma, [\mathbf{T}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}: A]}{\Gamma, [\mathbf{T}: A \vee_f B]} \vee_f : \mathbf{T}
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma, [N, N, T, B, NT, NB, TB, NTB: B] \quad \Gamma, [N, N, B, NB: A, B] \quad \Gamma, [N, T, B, TB: A, B] \quad \Gamma, [B, NB, TB, NTB: A, B] \quad \Gamma, [N, N, T, B, NT, NB, TB, NTB: A]}{\Gamma, [B: A \vee_f B]} \vee_f: B \\
\frac{\Gamma, [N, NF, NT, NFT: B] \quad \Gamma, [N, NF: A, B] \quad \Gamma, [NF, NFT: A, B] \quad \Gamma, [N, NF, NT, NFT: A]}{\Gamma, [NF: A \vee_f B]} \vee_f: NF \\
\frac{\Gamma, [NT: B] \quad \Gamma, [NT: A]}{\Gamma, [NT: A \vee_f B]} \vee_f: NT \\
\frac{\Gamma, [T, NT, FT, NFT: B] \quad \Gamma, [T, FT: A, B] \quad \Gamma, [FT, NFT: A, B] \quad \Gamma, [T, NT, FT, NFT: A]}{\Gamma, [FT: A \vee_f B]} \vee_f: FT \\
\frac{\Gamma, [N, NT, NB, NTB: B] \quad \Gamma, [N, NB: A, B] \quad \Gamma, [NB, NTB: A, B] \quad \Gamma, [N, NT, NB, NTB: A]}{\Gamma, [NB: A \vee_f B]} \vee_f: NB \\
\frac{\Gamma, [N, N, F, B, NF, NB, FB, NFB: A, B] \quad \Gamma, [N, F, T, B, FT, FB, TB, FTB: A, B] \quad \Gamma, [F, NF, FT, FB, NFT, NFB, FTB, A: A, B] \quad \Gamma, [B, NB, FB, TB, NFB, NTB, FTB, A: A, B]}{\Gamma, [FB: A \vee_f B]} \vee_f: FB \\
\frac{\Gamma, [T, NT, TB, NTB: B] \quad \Gamma, [T, TB: A, B] \quad \Gamma, [TB, NTB: A, B] \quad \Gamma, [T, NT, TB, NTB: A]}{\Gamma, [TB: A \vee_f B]} \vee_f: TB \\
\frac{\Gamma, [NT, NFT: B] \quad \Gamma, [NFT: A, B] \quad \Gamma, [NT, NFT: A]}{\Gamma, [NFT: A \vee_f B]} \vee_f: NFT \\
\frac{\Gamma, [N, NF, NT, NB, NFT, NFB, NTB, A: B] \quad \Gamma, [N, NF, NB, NFB: A, B] \quad \Gamma, [NF, NFT, NFB, A: A, B] \quad \Gamma, [NB, NFB, NTB, A: A, B] \quad \Gamma, [N, NF, NT, NB, NFT, NFB, NTB, A: A]}{\Gamma, [NFB: A \vee_f B]} \vee_f: NFB \\
\frac{\Gamma, [NT, NTB: B] \quad \Gamma, [NTB: A, B] \quad \Gamma, [NT, NTB: A]}{\Gamma, [NTB: A \vee_f B]} \vee_f: NTB \\
\frac{\Gamma, [T, NT, FT, TB, NFT, NTB, FTB, A: B] \quad \Gamma, [T, FT, TB, FTB: A, B] \quad \Gamma, [FT, NFT, FTB, A: A, B] \quad \Gamma, [TB, NTB, FTB, A: A, B] \quad \Gamma, [T, NT, FT, TB, NFT, NTB, FTB, A: A]}{\Gamma, [FTB: A \vee_f B]} \vee_f: FTB \\
\frac{\Gamma, [NT, NFT, NTB, A: B] \quad \Gamma, [NFT, A: A, B] \quad \Gamma, [NTB, A: A, B] \quad \Gamma, [NT, NFT, NTB, A: A]}{\Gamma, [A: A \vee_f B]} \vee_f: A
\end{array}$$

**Definition 11.** An upward tree of sequents is called a *proof* in the sequent calculus iff every leaf is an axiom, and all other sequents in it are obtained from the ones standing immediately above it by applying one of the rules. The sequent at the root of  $P$  is called its *end-sequent*. A sequent  $\Gamma$  is called *provable* iff it is the end-sequent of some proof.

**Theorem 12** (Soundness and Completeness). *A sequent is provable if and only if it is valid.*

*Proof.* See Theorems 3.1 and 3.2 of Baaz et al. [4] or Theorems 3.3.8 and 3.3.10 of Zach [15].  $\square$

**Corollary 13.** *In Shramko-Wansing logic,  $\models A$  iff  $[T, NT, TB, NTB: A]$  has a sequent proof, and  $\Delta \models A$  iff  $[N, N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: \Delta], [T, NT, TB, NTB: A]$  has a proof.*

**Theorem 14** (Cut-elimination). *The cut rule is eliminable in the sequent calculus for Shramko-Wansing logic.*

*Proof.* See Theorem 4.1 of Baaz et al. [4] or Theorem 3.5.3 of Zach [15].  $\square$

**Theorem 15** (Maehara lemma). *The Maehara lemma holds for the sequent calculus for Shramko-Wansing logic.*

*Proof.* See Theorem 3.8.1 of Zach [15].  $\square$

## 4 Tableaux for Shramko-Wansing logic

Although the method of Surma [13] and Carnielli [6] for obtaining signed analytic tableaux systems applies to Shramko-Wansing logic, it has a drawback. As Hähnle [7] pointed out, to show that a formula is valid, it is required to provide as many closed tableaux as there are non-designated values. This is usually not desirable; the generalized approach by Hähnle [7] solves this problem. Below we give a tableau system for Shramko-Wansing logic using the sets of signs  $V \setminus \{v\}$ , i.e., the tableau system exactly dual to that of Carnielli (in the sense of [2]).

**Definition 16.** A *signed formula* is an expression of the form  $v: A$  where  $v \in V$  and  $A$  is a formula.

**Definition 17.** A *tableau* for a set of signed formulas  $\Delta$  is a downward rooted tree of signed formulas where each one is either an element of  $\Delta$  or results from a signed formula in the branch above it by a branch expansion rule. A tableau is *closed* if every branch contains, for some formula  $A$ , the signed formulas  $v: A$  for all  $v \in V$ , or a signed formula  $v: A$  with a branch expansion rule that explicitly closes the branch ( $\otimes$ ).

The branch expansion rules for connective  $\neg_t$  are given by

$$\begin{array}{cccccc} \frac{\mathbf{N}: \neg_t A}{\mathbf{N}: A} & \frac{\mathbf{N}: \neg_t A}{\mathbf{T}: A} & \frac{\mathbf{F}: \neg_t A}{\mathbf{B}: A} & \frac{\mathbf{T}: \neg_t A}{\mathbf{N}: A} & \frac{\mathbf{B}: \neg_t A}{\mathbf{F}: A} & \frac{\mathbf{NF}: \neg_t A}{\mathbf{TB}: A} \\[10pt] \frac{\mathbf{NT}: \neg_t A}{\mathbf{NT}: A} & \frac{\mathbf{FT}: \neg_t A}{\mathbf{NB}: A} & \frac{\mathbf{NB}: \neg_t A}{\mathbf{FT}: A} & \frac{\mathbf{FB}: \neg_t A}{\mathbf{FB}: A} & \frac{\mathbf{TB}: \neg_t A}{\mathbf{NF}: A} & \frac{\mathbf{NFT}: \neg_t A}{\mathbf{NTB}: A} \\[10pt] \frac{\mathbf{NFB}: \neg_t A}{\mathbf{FTB}: A} & \frac{\mathbf{NTB}: \neg_t A}{\mathbf{NFT}: A} & \frac{\mathbf{FTB}: \neg_t A}{\mathbf{NFB}: A} & \frac{\mathbf{A}: \neg_t A}{\mathbf{A}: A} & & \end{array}$$

The branch expansion rules for connective  $\neg_f$  are given by

$$\begin{array}{cccccc} \frac{\mathbf{N}: \neg_f A}{\mathbf{N}: A} & \frac{\mathbf{N}: \neg_f A}{\mathbf{F}: A} & \frac{\mathbf{F}: \neg_f A}{\mathbf{N}: A} & \frac{\mathbf{T}: \neg_f A}{\mathbf{B}: A} & \frac{\mathbf{B}: \neg_f A}{\mathbf{T}: A} & \frac{\mathbf{NF}: \neg_f A}{\mathbf{NF}: A} \\[10pt] \frac{\mathbf{NT}: \neg_f A}{\mathbf{FB}: A} & \frac{\mathbf{FT}: \neg_f A}{\mathbf{NB}: A} & \frac{\mathbf{NB}: \neg_f A}{\mathbf{FT}: A} & \frac{\mathbf{FB}: \neg_f A}{\mathbf{NT}: A} & \frac{\mathbf{TB}: \neg_f A}{\mathbf{TB}: A} & \\[10pt] \frac{\mathbf{NFT}: \neg_f A}{\mathbf{NFB}: A} & \frac{\mathbf{NFB}: \neg_f A}{\mathbf{NFT}: A} & \frac{\mathbf{NTB}: \neg_f A}{\mathbf{FTB}: A} & \frac{\mathbf{FTB}: \neg_f A}{\mathbf{NTB}: A} & \frac{\mathbf{A}: \neg_f A}{\mathbf{A}: A} & \end{array}$$

The branch expansion rules for connective  $\wedge_t$  are given by

$$\begin{array}{c}
\frac{\mathbf{N}: A \wedge_t B}{\begin{array}{ccccc} \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A \\ \mathbf{T}: B & \mathbf{T}: A & \mathbf{B}: A & \mathbf{T}: A \\ \mathbf{B}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{B}: A \\ \mathbf{TB}: B & \mathbf{T}: B & \mathbf{B}: B & \mathbf{TB}: A \end{array}}
\\[10pt]
\frac{\mathbf{N}: A \wedge_t B}{\begin{array}{ccccc} \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A \\ \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{NT}: A & \mathbf{N}: A \\ \mathbf{T}: B & \mathbf{T}: A & \mathbf{B}: A & \mathbf{NB}: A & \mathbf{T}: A \\ \mathbf{B}: B & \mathbf{NT}: A & \mathbf{NB}: A & \mathbf{NTB}: A & \mathbf{B}: A \\ \mathbf{NT}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{NT}: A \\ \mathbf{NB}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{NT}: B & \mathbf{NB}: A \\ \mathbf{TB}: B & \mathbf{T}: B & \mathbf{B}: B & \mathbf{NB}: B & \mathbf{TB}: A \\ \mathbf{NTB}: B & \mathbf{NT}: B & \mathbf{NB}: B & \mathbf{NTB}: B & \mathbf{NTB}: A \end{array}}
\\[10pt]
\frac{\mathbf{F}: A \wedge_t B}{\begin{array}{ccccc} \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{F}: A & \mathbf{N}: A \\ \mathbf{F}: B & \mathbf{F}: A & \mathbf{F}: A & \mathbf{FT}: A & \mathbf{F}: A \\ \mathbf{T}: B & \mathbf{T}: A & \mathbf{B}: A & \mathbf{FB}: A & \mathbf{T}: A \\ \mathbf{B}: B & \mathbf{FT}: A & \mathbf{FB}: A & \mathbf{FTB}: A & \mathbf{B}: A \\ \mathbf{FT}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{F}: B & \mathbf{FT}: A \\ \mathbf{FB}: B & \mathbf{F}: B & \mathbf{F}: B & \mathbf{FT}: B & \mathbf{FB}: A \\ \mathbf{TB}: B & \mathbf{T}: B & \mathbf{B}: B & \mathbf{FB}: B & \mathbf{TB}: A \\ \mathbf{FTB}: B & \mathbf{FT}: B & \mathbf{FB}: B & \mathbf{FTB}: B & \mathbf{FTB}: A \end{array}}
\\[10pt]
\frac{\mathbf{T}: A \wedge_t B}{\mathbf{T}: B \quad \mathbf{T}: A \quad \mathbf{T}: A} \quad \frac{\mathbf{B}: A \wedge_t B}{\mathbf{B}: B \quad \mathbf{B}: A \quad \mathbf{B}: A} \quad \frac{}{\mathbf{TB}: B \quad \mathbf{TB}: A \quad \mathbf{B}: B \quad \mathbf{TB}: A}
\end{array}$$

$\text{NF}: A \wedge_t B$			
$\mathbf{N}: A$	$\mathbf{N}: A$	$\mathbf{N}: A$	$F: A$
$N: A$	$N: A$	$\mathbf{NF}: A$	$\mathbf{NF}: A$
$F: A$	$F: A$	$\mathbf{NT}: A$	$\mathbf{FT}: A$
$T: A$	$B: A$	$\mathbf{NB}: A$	$\mathbf{FB}: A$
$\mathbf{NF}: A$	$\mathbf{NF}: A$	$\mathbf{NFT}: A$	$\mathbf{NFT}: A$
$\mathbf{NT}: A$	$\mathbf{NB}: A$	$\mathbf{NFB}: A$	$\mathbf{NFB}: A$
$\mathbf{FT}: A$	$\mathbf{FB}: A$	$\mathbf{NTB}: A$	$\mathbf{FTB}: A$
$\mathbf{NFT}: A$	$\mathbf{NFB}: A$	$A: A$	$A: A$
$N: B$	$N: B$	$N: B$	$F: B$
$N: B$	$N: B$	$\mathbf{NF}: B$	$\mathbf{NF}: B$
$F: B$	$F: B$	$\mathbf{NT}: B$	$\mathbf{FT}: B$
$T: B$	$B: B$	$\mathbf{NB}: B$	$\mathbf{FB}: B$
$\mathbf{NF}: B$	$\mathbf{NF}: B$	$\mathbf{NFT}: B$	$\mathbf{NFT}: B$
$\mathbf{NT}: B$	$\mathbf{NB}: B$	$\mathbf{NFB}: B$	$\mathbf{NFB}: B$
$\mathbf{FT}: B$	$\mathbf{FB}: B$	$\mathbf{NTB}: B$	$\mathbf{FTB}: B$
$\mathbf{NFT}: B$	$\mathbf{NFB}: B$	$A: B$	$A: B$
$\mathbf{NT}: A \wedge_t B$			
$T: B$	$T: A$	$\mathbf{NT}: A$	$T: A$
$\mathbf{NT}: B$	$\mathbf{NT}: A$	$\mathbf{NTB}: A$	$\mathbf{NT}: A$
$\mathbf{TB}: B$	$T: B$	$\mathbf{NT}: B$	$\mathbf{TB}: A$
$\mathbf{NTB}: B$	$\mathbf{NT}: B$	$\mathbf{NTB}: B$	$\mathbf{NTB}: A$
$\mathbf{FT}: A \wedge_t B$			
$T: B$	$T: A$	$\mathbf{FT}: A$	$T: A$
$\mathbf{FT}: B$	$\mathbf{FT}: A$	$\mathbf{FTB}: A$	$\mathbf{FT}: A$
$\mathbf{TB}: B$	$T: B$	$\mathbf{FT}: B$	$\mathbf{TB}: A$
$\mathbf{FTB}: B$	$\mathbf{FT}: B$	$\mathbf{FTB}: B$	$\mathbf{FTB}: A$
$\mathbf{NB}: A \wedge_t B$			
$B: B$	$B: A$	$\mathbf{NB}: A$	$B: A$
$\mathbf{NB}: B$	$\mathbf{NB}: A$	$\mathbf{NTB}: A$	$\mathbf{NB}: A$
$\mathbf{TB}: B$	$B: B$	$\mathbf{NB}: B$	$\mathbf{TB}: A$
$\mathbf{NTB}: B$	$\mathbf{NB}: B$	$\mathbf{NTB}: B$	$\mathbf{NTB}: A$
$\mathbf{FB}: A \wedge_t B$			
$B: B$	$B: A$	$\mathbf{FB}: A$	$B: A$
$\mathbf{FB}: B$	$\mathbf{FB}: A$	$\mathbf{FTB}: A$	$\mathbf{FB}: A$
$\mathbf{TB}: B$	$B: B$	$\mathbf{FB}: B$	$\mathbf{TB}: A$
$\mathbf{FTB}: B$	$\mathbf{FB}: B$	$\mathbf{FTB}: B$	$\mathbf{FTB}: A$
$\mathbf{TB}: A \wedge_t B$			
		$\mathbf{TB}: B$	$\mathbf{TB}: A$

$$\begin{array}{ccccc}
& & \text{NFT: } A \wedge_t B & & \\
\hline
T: B & T: A & NT: A & FT: A & T: A \\
NT: B & NT: A & NFT: A & NFT: A & NT: A \\
FT: B & FT: A & NTB: A & FTB: A & FT: A \\
TB: B & NFT: A & A: A & A: A & TB: A \\
NFT: B & T: B & NT: B & FT: B & NFT: A \\
NTB: B & NT: B & NFT: B & NFT: B & NTB: A \\
FTB: B & FT: B & NTB: B & FTB: B & FTB: A \\
A: B & NFT: B & A: B & A: B & A: A
\end{array}$$
  

$$\begin{array}{ccccc}
& & \text{NFB: } A \wedge_t B & & \\
\hline
B: B & B: A & NB: A & FB: A & B: A \\
NB: B & NB: A & NFB: A & NFB: A & NB: A \\
FB: B & FB: A & NTB: A & FTB: A & FB: A \\
TB: B & NFB: A & A: A & A: A & TB: A \\
NFB: B & B: B & NB: B & FB: B & NFB: A \\
NTB: B & NB: B & NFB: B & NFB: B & NTB: A \\
FTB: B & FB: B & NTB: B & FTB: B & FTB: A \\
A: B & NFB: B & A: B & A: B & A: A
\end{array}$$
  

$$\begin{array}{ccc}
\text{NTB: } A \wedge_t B & & \text{FTB: } A \wedge_t B \\
\hline
TB: B & NTB: A & TB: A & TB: B & FTB: A & TB: A \\
NTB: B & NTB: B & NTB: A & FTB: B & FTB: B & FTB: A
\end{array}$$
  

$$\begin{array}{ccccc}
& & \text{A: } A \wedge_t B & & \\
\hline
& TB: B & NTB: A & FTB: A & TB: A \\
& NTB: B & A: A & A: A & NTB: A \\
& FTB: B & NTB: B & FTB: B & FTB: A \\
& A: B & A: B & A: B & A: A
\end{array}$$

The branch expansion rules for connective  $\vee_t$  are given by

$$\begin{array}{ccccccc}
& & \text{N: } A \vee_t B & & & \text{N: } A \vee_t B & \\
\hline
N: B & N: A & N: A & N: A & N: B & N: A & N: A \\
N: B & N: A & F: A & N: A & NF: B & N: B & NF: A \\
F: B & N: B & N: B & F: A & & & \\
NF: B & N: B & F: B & NF: A & & &
\end{array}$$
  

$$\begin{array}{ccc}
& & \text{F: } A \vee_t B \\
\hline
& F: B & F: A & F: A \\
& NF: B & F: B & NF: A
\end{array}$$

$T: A \vee_t B$				
$N: B$	$N: A$	$N: A$	$T: A$	$N: A$
$N: B$	$N: A$	$F: A$	$NT: A$	$N: A$
$F: B$	$T: A$	$T: A$	$FT: A$	$F: A$
$T: B$	$NT: A$	$FT: A$	$NFT: A$	$T: A$
$NF: B$	$N: B$	$N: B$	$T: B$	$NF: A$
$NT: B$	$N: B$	$F: B$	$NT: B$	$NT: A$
$FT: B$	$T: B$	$T: B$	$FT: B$	$FT: A$
$NFT: B$	$NT: B$	$FT: B$	$NFT: B$	$NFT: A$
$B: A \vee_t B$				
$N: B$	$N: A$	$N: A$	$B: A$	$N: A$
$N: B$	$N: A$	$F: A$	$NB: A$	$N: A$
$F: B$	$B: A$	$B: A$	$FB: A$	$F: A$
$B: B$	$NB: A$	$FB: A$	$NFB: A$	$B: A$
$NF: B$	$N: B$	$N: B$	$B: B$	$NF: A$
$NB: B$	$N: B$	$F: B$	$NB: B$	$NB: A$
$FB: B$	$B: B$	$B: B$	$FB: B$	$FB: A$
$NFB: B$	$NB: B$	$FB: B$	$NFB: B$	$NFB: A$
$NT: A \vee_t B$				
$N: B$	$N: A$	$NT: A$	$N: A$	
$NF: B$	$NT: A$	$NFT: A$	$NF: A$	
$NT: B$	$N: B$	$NT: B$	$NT: A$	
$NFT: B$	$NT: B$	$NFT: B$	$NFT: A$	
$FT: A \vee_t B$				
$F: B$	$F: A$	$FT: A$	$F: A$	
$NF: B$	$FT: A$	$NFT: A$	$NF: A$	
$FT: B$	$F: B$	$FT: B$	$FT: A$	
$NFT: B$	$FT: B$	$NFT: B$	$NFT: A$	
$NB: A \vee_t B$				
$N: B$	$N: A$	$NB: A$	$N: A$	
$NF: B$	$NB: A$	$NFB: A$	$NF: A$	
$NB: B$	$N: B$	$NB: B$	$NB: A$	
$NFB: B$	$NB: B$	$NFB: B$	$NFB: A$	
$FB: A \vee_t B$				
$F: B$	$F: A$	$FB: A$	$F: A$	
$NF: B$	$FB: A$	$NFB: A$	$NF: A$	
$FB: B$	$F: B$	$FB: B$	$FB: A$	
$NFB: B$	$FB: B$	$NFB: B$	$NFB: A$	

$\mathbf{TB}: A \vee_t B$				
$N: A$	$N: A$	$T: A$	$B: A$	
$N: A$	$F: A$	$NT: A$	$NB: A$	
$T: A$	$T: A$	$FT: A$	$FB: A$	
$B: A$	$B: A$	$TB: A$	$TB: A$	
$NT: A$	$FT: A$	$NFT: A$	$NFB: A$	
$NB: A$	$FB: A$	$NTB: A$	$NTB: A$	
$TB: A$	$TB: A$	$FTB: A$	$FTB: A$	
$NTB: A$	$FTB: A$	$A: A$	$A: A$	
$N: B$	$N: B$	$T: B$	$B: B$	
$N: B$	$F: B$	$NT: B$	$NB: B$	
$T: B$	$T: B$	$FT: B$	$FB: B$	
$B: B$	$B: B$	$TB: B$	$TB: B$	
$NT: B$	$FT: B$	$NFT: B$	$NFB: B$	
$NB: B$	$FB: B$	$NTB: B$	$NTB: B$	
$TB: B$	$TB: B$	$FTB: B$	$FTB: B$	
$NTB: B$	$FTB: B$	$A: B$	$A: B$	
$\mathbf{NFT}: A \vee_t B$				
$NF: B$	$NFT: A$	$NF: A$	$NF: B$	$NFB: A$
$NFT: B$	$NFT: B$	$NFT: A$	$NFB: B$	$NFB: A$
$\mathbf{NTB}: A \vee_t B$				
$N: B$	$N: A$	$NT: A$	$NB: A$	$N: A$
$NF: B$	$NT: A$	$NFT: A$	$NFB: A$	$NF: A$
$NT: B$	$NB: A$	$NTB: A$	$NTB: A$	$NT: A$
$NB: B$	$NTB: A$	$A: A$	$A: A$	$NB: A$
$NFT: B$	$N: B$	$NT: B$	$NB: B$	$NFT: A$
$NFB: B$	$NT: B$	$NFT: B$	$NFB: B$	$NFB: A$
$NTB: B$	$NB: B$	$NTB: B$	$NTB: B$	$NTB: A$
$A: B$	$NTB: B$	$A: B$	$A: B$	$A: A$
$\mathbf{FTB}: A \vee_t B$				
$F: B$	$F: A$	$FT: A$	$FB: A$	$F: A$
$NF: B$	$FT: A$	$NFT: A$	$NFB: A$	$NF: A$
$FT: B$	$FB: A$	$FTB: A$	$FTB: A$	$FT: A$
$FB: B$	$FTB: A$	$A: A$	$A: A$	$FB: A$
$NFT: B$	$F: B$	$FT: B$	$FB: B$	$NFT: A$
$NFB: B$	$FT: B$	$NFT: B$	$NFB: B$	$NFB: A$
$FTB: B$	$FB: B$	$FTB: B$	$FTB: B$	$FTB: A$
$A: B$	$FTB: B$	$A: B$	$A: B$	$A: A$

$$\begin{array}{cccc}
& \mathbf{A}: A \vee_t B & & \\
\hline
\mathbf{NF}: B & \mathbf{NFT}: A & \mathbf{NFB}: A & \mathbf{NF}: A \\
\mathbf{NFT}: B & \mathbf{A}: A & \mathbf{A}: A & \mathbf{NFT}: A \\
\mathbf{NFB}: B & \mathbf{NFT}: B & \mathbf{NFB}: B & \mathbf{NFB}: A \\
\mathbf{A}: B & \mathbf{A}: B & \mathbf{A}: B & \mathbf{A}: A
\end{array}$$

The branch expansion rules for connective  $\wedge_f$  are given by

$$\begin{array}{ccccc}
& \mathbf{N}: A \wedge_f B & & & \\
\hline
\mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A & \\
\mathbf{F}: B & \mathbf{F}: A & \mathbf{B}: A & \mathbf{F}: A & \\
\mathbf{B}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{B}: A & \\
\mathbf{FB}: B & \mathbf{F}: B & \mathbf{B}: B & \mathbf{FB}: A &
\end{array}$$
  

$$\begin{array}{ccccc}
& \mathbf{N}: A \wedge_f B & & & \\
\hline
\mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A \\
\mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{NF}: A & \mathbf{N}: A \\
\mathbf{F}: B & \mathbf{F}: A & \mathbf{B}: A & \mathbf{NB}: A & \mathbf{F}: A \\
\mathbf{B}: B & \mathbf{NF}: A & \mathbf{NB}: A & \mathbf{NFB}: A & \mathbf{B}: A \\
\mathbf{NF}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{NF}: A \\
\mathbf{NB}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{NF}: B & \mathbf{NB}: A \\
\mathbf{FB}: B & \mathbf{F}: B & \mathbf{B}: B & \mathbf{NB}: B & \mathbf{FB}: A \\
\mathbf{NFB}: B & \mathbf{NF}: B & \mathbf{NB}: B & \mathbf{NFB}: B & \mathbf{NFB}: A
\end{array}$$
  

$$\begin{array}{ccc}
& \mathbf{F}: A \wedge_f B & \\
\hline
& \mathbf{F}: B & \mathbf{F}: A \\
& \mathbf{FB}: B & \mathbf{F}: B & \mathbf{FB}: A
\end{array}$$
  

$$\begin{array}{ccccc}
& \mathbf{T}: A \wedge_f B & & & \\
\hline
\mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{T}: A & \mathbf{N}: A \\
\mathbf{F}: B & \mathbf{F}: A & \mathbf{T}: A & \mathbf{FT}: A & \mathbf{F}: A \\
\mathbf{T}: B & \mathbf{T}: A & \mathbf{B}: A & \mathbf{TB}: A & \mathbf{T}: A \\
\mathbf{B}: B & \mathbf{FT}: A & \mathbf{TB}: A & \mathbf{FTB}: A & \mathbf{B}: A \\
\mathbf{FT}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{T}: B & \mathbf{FT}: A \\
\mathbf{FB}: B & \mathbf{F}: B & \mathbf{T}: B & \mathbf{FT}: B & \mathbf{FB}: A \\
\mathbf{TB}: B & \mathbf{T}: B & \mathbf{B}: B & \mathbf{TB}: B & \mathbf{TB}: A \\
\mathbf{FTB}: B & \mathbf{FT}: B & \mathbf{TB}: B & \mathbf{FTB}: B & \mathbf{FTB}: A
\end{array}$$
  

$$\begin{array}{ccc}
& \mathbf{B}: A \wedge_f B & \\
\hline
\mathbf{B}: B & \mathbf{B}: A & \mathbf{B}: A \\
\mathbf{FB}: B & \mathbf{B}: B & \mathbf{FB}: A
\end{array}
\quad
\begin{array}{ccccc}
& \mathbf{NF}: A \wedge_f B & & & \\
\hline
\mathbf{F}: B & \mathbf{F}: A & \mathbf{NF}: A & \mathbf{F}: A & \\
\mathbf{NF}: B & \mathbf{NF}: A & \mathbf{NFB}: A & \mathbf{NF}: A & \\
& \mathbf{FB}: B & \mathbf{F}: B & \mathbf{NF}: B & \mathbf{FB}: A \\
& \mathbf{NFB}: B & \mathbf{NF}: B & \mathbf{NFB}: B & \mathbf{NFB}: A
\end{array}$$

$\text{NT}: A \wedge_f B$				
$N: A$	$N: A$	$N: A$	$T: A$	
$N: A$	$N: A$	$NF: A$	$NT: A$	
$F: A$	$T: A$	$NT: A$	$FT: A$	
$T: A$	$B: A$	$NB: A$	$TB: A$	
$NF: A$	$NT: A$	$NFT: A$	$NFT: A$	
$NT: A$	$NB: A$	$NFB: A$	$NTB: A$	
$FT: A$	$TB: A$	$NTB: A$	$FTB: A$	
$NFT: A$	$NTB: A$	$A: A$	$A: A$	
$N: B$	$N: B$	$N: B$	$T: B$	
$N: B$	$N: B$	$NF: B$	$NT: B$	
$F: B$	$T: B$	$NT: B$	$FT: B$	
$T: B$	$B: B$	$NB: B$	$TB: B$	
$NF: B$	$NT: B$	$NFT: B$	$NFT: B$	
$NT: B$	$NB: B$	$NFB: B$	$NTB: B$	
$FT: B$	$TB: B$	$NTB: B$	$FTB: B$	
$NFT: B$	$NTB: B$	$A: B$	$A: B$	
$FT: A \wedge_f B$				
$F: B$	$F: A$	$FT: A$	$F: A$	
$FT: B$	$FT: A$	$FTB: A$	$FT: A$	
$FB: B$	$F: B$	$FT: B$	$FB: A$	
$FTB: B$	$FT: B$	$FTB: B$	$FTB: A$	
$NB: A \wedge_f B$				
$B: B$	$B: A$	$NB: A$	$B: A$	$FB: B$
$NB: B$	$NB: A$	$NFB: A$	$NB: A$	$FB: A$
$FB: B$	$B: B$	$NB: B$	$FB: A$	
$NFB: B$	$NB: B$	$NFB: B$	$NFB: A$	
$TB: A \wedge_f B$				
$B: B$	$B: A$	$TB: A$	$B: A$	
$FB: B$	$TB: A$	$FTB: A$	$FB: A$	
$TB: B$	$B: B$	$TB: B$	$TB: A$	
$FTB: B$	$TB: B$	$FTB: B$	$FTB: A$	
$NFT: A \wedge_f B$				
$F: B$	$F: A$	$NF: A$	$FT: A$	$F: A$
$NF: B$	$NF: A$	$NFT: A$	$NFT: A$	$NF: A$
$FT: B$	$FT: A$	$NFB: A$	$FTB: A$	$FT: A$
$FB: B$	$NFT: A$	$A: A$	$A: A$	$FB: A$
$NFT: B$	$F: B$	$NF: B$	$FT: B$	$NFT: A$
$NFB: B$	$NF: B$	$NFT: B$	$NFT: B$	$NFB: A$
$FTB: B$	$FT: B$	$NFB: B$	$FTB: B$	$FTB: A$
$A: B$	$NFT: B$	$A: B$	$A: B$	$A: A$

$$\begin{array}{c}
\frac{\text{NFB}: A \wedge_f B}{\begin{array}{ccc} \text{FB}: B & \text{NFB}: A & \text{FB}: A \\ \text{NFB}: B & \text{NFB}: B & \text{NFB}: A \end{array}} \\
\frac{\text{NTB}: A \wedge_f B}{\begin{array}{ccccc} B: B & B: A & \text{NB}: A & \text{TB}: A & B: A \\ \text{NB}: B & \text{NB}: A & \text{NFB}: A & \text{NTB}: A & \text{NB}: A \\ \text{FB}: B & \text{TB}: A & \text{NTB}: A & \text{FTB}: A & \text{FB}: A \\ \text{TB}: B & \text{NTB}: A & \text{A}: A & \text{A}: A & \text{TB}: A \\ \text{NFB}: B & B: B & \text{NB}: B & \text{TB}: B & \text{NFB}: A \\ \text{NTB}: B & \text{NB}: B & \text{NFB}: B & \text{NTB}: B & \text{NTB}: A \\ \text{FTB}: B & \text{TB}: B & \text{NTB}: B & \text{FTB}: B & \text{FTB}: A \\ \text{A}: B & \text{NTB}: B & \text{A}: B & \text{A}: B & \text{A}: A \end{array}} \\
\frac{\text{FTB}: A \wedge_f B}{\begin{array}{ccc} \text{FB}: B & \text{FTB}: A & \text{FB}: A \\ \text{FTB}: B & \text{FTB}: B & \text{FTB}: A \end{array}} \\
\frac{\text{A}: A \wedge_f B}{\begin{array}{cccc} \text{FB}: B & \text{NFB}: A & \text{FTB}: A & \text{FB}: A \\ \text{NFB}: B & \text{A}: A & \text{A}: A & \text{NFB}: A \\ \text{FTB}: B & \text{NFB}: B & \text{FTB}: B & \text{FTB}: A \\ \text{A}: B & \text{A}: B & \text{A}: B & \text{A}: A \end{array}}
\end{array}$$

The branch expansion rules for connective  $\vee_f$  are given by

$$\begin{array}{c}
\frac{\text{N}: A \vee_f B}{\begin{array}{cccc} \text{N}: B & \text{N}: A & \text{N}: A & \text{N}: A \\ \text{N}: B & \text{N}: A & \text{T}: A & \text{N}: A \\ \text{T}: B & \text{N}: B & \text{N}: B & \text{T}: A \\ \text{NT}: B & \text{N}: B & \text{T}: B & \text{NT}: A \end{array}} \quad \frac{\text{N}: A \vee_f B}{\begin{array}{ccc} \text{N}: B & \text{N}: A & \text{N}: A \\ \text{NT}: B & \text{N}: B & \text{NT}: B \\ \text{N}: A & \text{N}: A & \text{NT}: A \end{array}} \\
\frac{\text{F}: A \vee_f B}{\begin{array}{ccccc} \text{N}: B & \text{N}: A & \text{N}: A & \text{F}: A & \text{N}: A \\ \text{N}: B & \text{N}: A & \text{F}: A & \text{NF}: A & \text{N}: A \\ \text{F}: B & \text{F}: A & \text{T}: A & \text{FT}: A & \text{F}: A \\ \text{T}: B & \text{NF}: A & \text{FT}: A & \text{NFT}: A & \text{T}: A \\ \text{NF}: B & \text{N}: B & \text{N}: B & \text{F}: B & \text{NF}: A \\ \text{NT}: B & \text{N}: B & \text{F}: B & \text{NF}: B & \text{NT}: A \\ \text{FT}: B & \text{F}: B & \text{T}: B & \text{FT}: B & \text{FT}: A \\ \text{NFT}: B & \text{NF}: B & \text{FT}: B & \text{NFT}: B & \text{NFT}: A \end{array}} \\
\frac{\text{T}: A \vee_f B}{\begin{array}{ccc} \text{T}: B & \text{T}: A & \text{T}: A \\ \text{NT}: B & \text{T}: B & \text{NT}: A \end{array}}
\end{array}$$

$B: A \vee_f B$				
$N:B$	$N:A$	$N:A$	$B:A$	$N:A$
$N:B$	$N:A$	$T:A$	$NB:A$	$N:A$
$T:B$	$B:A$	$B:A$	$TB:A$	$T:A$
$B:B$	$NB:A$	$TB:A$	$NTB:A$	$B:A$
$NT:B$	$N:B$	$N:B$	$B:B$	$NT:A$
$NB:B$	$N:B$	$T:B$	$NB:B$	$NB:A$
$TB:B$	$B:B$	$B:B$	$TB:B$	$TB:A$
$NTB:B$	$NB:B$	$TB:B$	$NTB:B$	$NTB:A$
$\text{NF}: A \vee_f B$				
$N:B$	$N:A$	$NF:A$	$N:A$	$NT:B$
$NF:B$	$NF:A$	$NFT:A$	$NF:A$	$NT:A$
$NT:B$	$N:B$	$NF:B$	$NT:A$	
$NFT:B$	$NF:B$	$NFT:B$	$NFT:A$	
$FT: A \vee_f B$				
$T:B$	$T:A$	$FT:A$	$T:A$	
$NT:B$	$FT:A$	$NFT:A$	$NT:A$	
$FT:B$	$T:B$	$FT:B$	$FT:A$	
$NFT:B$	$FT:B$	$NFT:B$	$NFT:A$	
$NB: A \vee_f B$				
$N:B$	$N:A$	$NB:A$	$N:A$	
$NT:B$	$NB:A$	$NTB:A$	$NT:A$	
$NB:B$	$N:B$	$NB:B$	$NB:A$	
$NTB:B$	$NB:B$	$NTB:B$	$NTB:A$	
$FB: A \vee_f B$				
$N:A$	$N:A$	$F:A$	$B:A$	
$N:A$	$F:A$	$NF:A$	$NB:A$	
$F:A$	$T:A$	$FT:A$	$FB:A$	
$B:A$	$B:A$	$FB:A$	$TB:A$	
$NF:A$	$FT:A$	$NFT:A$	$NFB:A$	
$NB:A$	$FB:A$	$NFB:A$	$NTB:A$	
$FB:A$	$TB:A$	$FTB:A$	$FTB:A$	
$NFB:A$	$FTB:A$	$A:A$	$A:A$	
$N:B$	$N:B$	$F:B$	$B:B$	
$N:B$	$F:B$	$NF:B$	$NB:B$	
$F:B$	$T:B$	$FT:B$	$FB:B$	
$B:B$	$B:B$	$FB:B$	$TB:B$	
$NF:B$	$FT:B$	$NFT:B$	$NFB:B$	
$NB:B$	$FB:B$	$NFB:B$	$NTB:B$	
$FB:B$	$TB:B$	$FTB:B$	$FTB:B$	
$NFB:B$	$FTB:B$	$A:B$	$A:B$	

$\text{TB}: A \vee_f B$				
$T: B$	$T: A$	$\text{TB}: A$	$T: A$	
$\text{NT}: B$	$\text{TB}: A$	$\text{NTB}: A$	$\text{NT}: A$	
$\text{TB}: B$	$T: B$	$\text{TB}: B$	$\text{TB}: A$	
$\text{NTB}: B$	$\text{TB}: B$	$\text{NTB}: B$	$\text{NTB}: A$	
$\text{NFT}: A \vee_f B$				
$\text{NT}: B$	$\text{NFT}: A$	$\text{NT}: A$		
$\text{NFT}: B$	$\text{NFT}: B$	$\text{NFT}: A$		
$\text{NFT}: B$	$N: B$	$\text{NF}: B$	$\text{NB}: B$	$\text{NFT}: A$
$\text{NFB}: B$	$\text{NF}: B$	$\text{NFT}: B$	$\text{NFB}: B$	$\text{NFB}: A$
$\text{NFB}: B$	$N: B$	$\text{NF}: B$	$\text{NB}: B$	$\text{NFT}: A$
$\text{NFB}: B$	$\text{NF}: B$	$\text{NFT}: B$	$\text{NFB}: B$	$\text{NFB}: A$
$\text{NTB}: B$	$\text{NB}: B$	$\text{NFB}: B$	$\text{NTB}: B$	$\text{NTB}: A$
$A: B$	$\text{NFB}: B$	$A: B$	$A: B$	$A: A$
$\text{NFB}: A \vee_f B$				
$N: B$	$N: A$	$\text{NF}: A$	$\text{NB}: A$	$N: A$
$\text{NF}: B$	$\text{NF}: A$	$\text{NFT}: A$	$\text{NFB}: A$	$\text{NF}: A$
$\text{NT}: B$	$\text{NB}: A$	$\text{NFB}: A$	$\text{NTB}: A$	$\text{NT}: A$
$\text{NB}: B$	$\text{NFB}: A$	$A: A$	$A: A$	$\text{NB}: A$
$\text{NFT}: B$	$N: B$	$\text{NF}: B$	$\text{NB}: B$	$\text{NFT}: A$
$\text{NFB}: B$	$\text{NF}: B$	$\text{NFT}: B$	$\text{NFB}: B$	$\text{NFB}: A$
$\text{NTB}: B$	$\text{NB}: B$	$\text{NFB}: B$	$\text{NTB}: B$	$\text{NTB}: A$
$A: B$	$\text{NFB}: B$	$A: B$	$A: B$	$A: A$
$\text{NTB}: A \vee_f B$				
$\text{NT}: B$	$\text{NTB}: A$	$\text{NT}: A$		
$\text{NTB}: B$	$\text{NTB}: B$	$\text{NTB}: A$		
$\text{FTB}: A \vee_f B$				
$T: B$	$T: A$	$\text{FT}: A$	$\text{TB}: A$	$T: A$
$\text{NT}: B$	$\text{FT}: A$	$\text{NFT}: A$	$\text{NTB}: A$	$\text{NT}: A$
$\text{FT}: B$	$\text{TB}: A$	$\text{FTB}: A$	$\text{FTB}: A$	$\text{FT}: A$
$\text{TB}: B$	$\text{FTB}: A$	$A: A$	$A: A$	$\text{TB}: A$
$\text{NFT}: B$	$T: B$	$\text{FT}: B$	$\text{TB}: B$	$\text{NFT}: A$
$\text{NTB}: B$	$\text{FT}: B$	$\text{NFT}: B$	$\text{NTB}: B$	$\text{NTB}: A$
$\text{FTB}: B$	$\text{TB}: B$	$\text{FTB}: B$	$\text{FTB}: B$	$\text{FTB}: A$
$A: B$	$\text{FTB}: B$	$A: B$	$A: B$	$A: A$
$A: A \vee_f B$				
$\text{NT}: B$	$\text{NFT}: A$	$\text{NTB}: A$	$\text{NT}: A$	
$\text{NFT}: B$	$A: A$	$A: A$	$\text{NFT}: A$	
$\text{NTB}: B$	$\text{NFT}: B$	$\text{NTB}: B$	$\text{NTB}: A$	
$A: B$	$A: B$	$A: B$	$A: A$	

**Definition 18.** An interpretation  $\mathfrak{I}$  satisfies a signed formula  $v: A$  iff  $\text{val}_{\mathfrak{I}}(A) \neq v$ . A set of signed formulas is satisfiable if some interpretation  $\mathfrak{I}$  satisfies all signed formulas in it.

**Theorem 19.** A set of signed formulas is unsatisfiable iff it has a closed tableau.

*Proof.* Apply Theorems 4.14 and 4.21 of Hähnle [7]; interpreting  $v: A$  as  $S A$  where  $S = V \setminus \{v\}$ .  $\square$

**Corollary 20.** In Shramko-Wansing logic,  $\models A$  iff  $\{v: A \mid v \in V^+\}$  has a closed tableau.  $\Delta \models A$  iff  $\{v: B \mid v \in V^-, B \in \Delta\} \cup \{v: A \mid v \in V^+\}$  has a closed tableau.

## 5 Natural deduction for Shramko-Wansing logic

Let  $\Gamma$  be a (set) sequent,  $V^+ \subseteq V$  the set of *designated truth values*. The set of non-designated truth values is then  $V^- = V \setminus V^+$ . We divide the sequent  $\Gamma$  into its designated part  $\Gamma^+$  and its non-designated part  $\Gamma^-$  in the obvious way:

$$\begin{aligned}\Gamma^+ &:= \langle \Gamma_v \mid v \in V^+ \rangle \\ \Gamma^- &:= \langle \Gamma_v \mid v \in V^- \rangle\end{aligned}$$

**Definition 21.** The *natural deduction rules* for Shramko-Wansing logic are given by:

1. A weakening rule for all  $v \in V^+$ :

$$\frac{\Gamma^+}{\Gamma^+, [v: A]} \text{ W}:v$$

2. For every connective  $\square$  and every truth value  $v$  an introduction rule  $\square I:v$  (if  $v \in V^+$ ) or an elimination rule  $\square E:v$  (if  $v \in V^-$ ).

The introduction and elimination rules for connective  $\neg_t$  are given by

$$\begin{array}{c} \Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{N}: A]] \\ \frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \qquad \Gamma_1^+, [\mathbf{T}: A]}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{N} \\ \\ \Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{B}: A]] \\ \frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \qquad \Gamma_1^+, [\mathbf{T}: A]}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{F} \\ \\ \Gamma_1^-, [[\mathbf{N}: A]] \\ \frac{\Gamma_1^+}{\Gamma_1^+, [\mathbf{T}: \neg_t A]} \neg_t I: \mathbf{T} \\ \\ \Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{F}: A]] \\ \frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \qquad \Gamma_1^+, [\mathbf{T}: A]}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{B} \end{array}$$

$$\begin{array}{c}
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-}{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+, [\mathbf{TB}: A]} \quad \frac{\Gamma_1^-}{\Gamma_1^+, [\mathbf{TB}: A]} \quad \neg_t E: \mathbf{NF} \\
\\
\frac{\Gamma_1^-}{\Gamma_1^+, [\mathbf{NT}: A]} \quad \frac{\Gamma_1^+, [\mathbf{NT}: A]}{\Gamma_1^+, [\mathbf{NT}: \neg_t A]} \quad \neg_t I: \mathbf{NT} \\
\\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{NB}: A]]}{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+} \quad \frac{\Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \quad \neg_t E: \mathbf{FT} \\
\\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{FT}: A]]}{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+} \quad \frac{\Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \quad \neg_t E: \mathbf{NB} \\
\\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{FB}: A]]}{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+} \quad \frac{\Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \quad \neg_t E: \mathbf{FB} \\
\\
\frac{\Gamma_1^-, [[\mathbf{NF}: A]]}{\Gamma_1^+, [\mathbf{TB}: \neg_t A]} \quad \frac{\Gamma_1^+}{\Gamma_1^+, [\mathbf{TB}: \neg_t A]} \quad \neg_t I: \mathbf{TB} \\
\\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-}{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+, [\mathbf{NTB}: A]} \quad \frac{\Gamma_1^-}{\Gamma_0^+, \Gamma_1^+} \quad \neg_t E: \mathbf{NFT} \\
\\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{FTB}: A]]}{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+} \quad \frac{\Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \quad \neg_t E: \mathbf{NFB} \\
\\
\frac{\Gamma_1^-, [[\mathbf{NFT}: A]]}{\Gamma_1^+, [\mathbf{NTB}: \neg_t A]} \quad \frac{\Gamma_1^+}{\Gamma_1^+, [\mathbf{NTB}: \neg_t A]} \quad \neg_t I: \mathbf{NTB} \\
\\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{NFB}: A]]}{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+} \quad \frac{\Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \quad \neg_t E: \mathbf{FTB} \\
\\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{A}: A]]}{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+} \quad \frac{\Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \quad \neg_t E: \mathbf{A}
\end{array}$$

The introduction and elimination rules for connective  $\neg_f$  are given by

$$\begin{array}{c}
\Gamma_0^-, [[N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: \neg_f A]] \quad \Gamma_1^-, [[N: A]] \\
\frac{\Gamma_0^+, [T, NT, TB, NTB: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: N
\end{array}$$

$$\begin{array}{c}
\Gamma_0^-, [[N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: \neg_f A]] \quad \Gamma_1^-, [[F: A]] \\
\frac{\Gamma_0^+, [T, NT, TB, NTB: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: N
\end{array}$$

$$\begin{array}{c}
\Gamma_0^-, [[N, N, B, NF, FT, NB, FB, NFT, NFB, FTB, A: \neg_f A]] \quad \Gamma_1^-, [[N: A]] \\
\frac{\Gamma_0^+, [T, NT, TB, NTB: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: F
\end{array}$$

$$\begin{array}{c}
\Gamma_1^-, [[B: A]] \\
\frac{\Gamma_1^+}{\Gamma_1^+, [T: \neg_f A]} \neg_f I: T
\end{array}$$

$$\begin{array}{c}
\Gamma_0^-, [[N, N, F, NF, FT, NB, FB, NFT, NFB, FTB, A: \neg_f A]] \quad \Gamma_1^- \\
\frac{\Gamma_0^+, [T, NT, TB, NTB: \neg_f A] \quad \Gamma_1^+, [T: A]}{\Gamma_0^+, \Gamma_1^+} \neg_f E: B
\end{array}$$

$$\begin{array}{c}
\Gamma_0^-, [[N, N, F, B, FT, NB, FB, NFT, NFB, FTB, A: \neg_f A]] \quad \Gamma_1^-, [[NF: A]] \\
\frac{\Gamma_0^+, [T, NT, TB, NTB: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: NF
\end{array}$$

$$\begin{array}{c}
\Gamma_1^-, [[FB: A]] \\
\frac{\Gamma_1^+}{\Gamma_1^+, [NT: \neg_f A]} \neg_f I: NT
\end{array}$$

$$\begin{array}{c}
\Gamma_0^-, [[N, N, F, B, NF, NB, FB, NFT, NFB, FTB, A: \neg_f A]] \quad \Gamma_1^-, [[NB: A]] \\
\frac{\Gamma_0^+, [T, NT, TB, NTB: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: FT
\end{array}$$

$$\begin{array}{c}
\Gamma_0^-, [[N, N, F, B, NF, FT, FB, NFT, NFB, FTB, A: \neg_f A]] \quad \Gamma_1^-, [[FT: A]] \\
\frac{\Gamma_0^+, [T, NT, TB, NTB: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: NB
\end{array}$$

$$\begin{array}{c}
\Gamma_0^-, [[N, N, F, B, NF, FT, NB, NFT, NFB, FTB, A: \neg_f A]] \quad \Gamma_1^- \\
\frac{\Gamma_0^+, [T, NT, TB, NTB: \neg_f A] \quad \Gamma_1^+, [NT: A]}{\Gamma_0^+, \Gamma_1^+} \neg_f E: FB
\end{array}$$

$$\frac{\Gamma_1^-, \Gamma_1^+, [\mathbf{TB}: A]}{\Gamma_1^+, [\mathbf{TB}: \neg_f A]} \neg_f \mathbf{I} : \mathbf{TB}$$

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{NFB}: A]]}{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+} \neg_f \mathbf{E} : \mathbf{NFT}$$

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{NFT}: A]]}{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+} \neg_f \mathbf{E} : \mathbf{NFB}$$

$$\frac{\Gamma_1^-, [[\mathbf{FTB}: A]] \quad \Gamma_1^+}{\Gamma_1^+, [\mathbf{NTB}: \neg_f A]} \neg_f \mathbf{I} : \mathbf{NTB}$$

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, \Gamma_1^+}{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+, [\mathbf{NTB}: A]} \neg_f \mathbf{E} : \mathbf{FTB}$$

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{A}: A]]}{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+} \neg_f \mathbf{E} : \mathbf{A}$$

The introduction and elimination rules for connective  $\wedge_t$  are given by

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_t B]] \quad \Gamma_1^-, [[\mathbf{N}, \mathbf{B}: B]] \quad \Gamma_2^-, [[\mathbf{N}: A, B]] \quad \Gamma_3^-, [[\mathbf{N}, \mathbf{B}: A, B]] \quad \Gamma_4^-, [[\mathbf{N}, \mathbf{B}: A]]}{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_t B] \quad \Gamma_1^+, [T, \mathbf{TB}: B] \quad \Gamma_2^+, [T: A, B] \quad \Gamma_3^+ \quad \Gamma_4^+, [T, \mathbf{TB}: A]} \wedge_t \mathbf{E} : \mathbf{N}$$

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_t B]] \quad \Gamma_1^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: B]] \quad \Gamma_2^-, [[\mathbf{N}, \mathbf{N}: A, B]] \quad \Gamma_3^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A, B]] \quad \Gamma_4^-, [[\mathbf{N}, \mathbf{NB}: A, B]] \quad \Gamma_5^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A]]}{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_t B] \quad \Gamma_1^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: B] \quad \Gamma_2^+, [T, \mathbf{NT}: A, B] \quad \Gamma_3^+ \quad \Gamma_4^+, [\mathbf{NT}, \mathbf{NTB}: A, B] \quad \Gamma_5^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A]} \wedge_t \mathbf{E} : \mathbf{N}$$

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_t B]] \quad \Gamma_1^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: B]] \quad \Gamma_2^-, [[\mathbf{N}, \mathbf{F}, \mathbf{FT}: A, B]] \quad \Gamma_3^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FB}: A, B]] \quad \Gamma_4^-, [[\mathbf{F}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A, B]] \quad \Gamma_5^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A]]}{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_t B] \quad \Gamma_1^+, [T, \mathbf{TB}: B] \quad \Gamma_2^+, [T: A, B] \quad \Gamma_3^+ \quad \Gamma_4^+ \quad \Gamma_5^+, [T, \mathbf{TB}: A]} \wedge_t \mathbf{E} : \mathbf{F}$$

$$\frac{\Gamma_1^-, \Gamma_2^-, \Gamma_3^-}{\Gamma_1^+, [T, \mathbf{TB}: B] \quad \Gamma_2^+, [T: A, B] \quad \Gamma_3^+, [T, \mathbf{TB}: A]} \wedge_t \mathbf{I} : \mathbf{T}$$

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_t B]] \quad \Gamma_1^-, [[\mathbf{B}: B]] \quad \Gamma_2^-, [[\mathbf{B}: A, B]] \quad \Gamma_3^-, [[\mathbf{B}: A]]}{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_t B] \quad \Gamma_1^+, [\mathbf{TB}: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [\mathbf{TB}: A]} \wedge_t \mathbf{E} : \mathbf{B}$$

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_t B]] \quad \Gamma_1^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: A, B]] \quad \Gamma_2^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A, B]] \quad \Gamma_3^-, [[\mathbf{N}, \mathbf{NF}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: A, B]] \quad \Gamma_4^-, [[\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B]]}{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_t B] \quad \Gamma_1^+, [T, \mathbf{NT}: A, B] \quad \Gamma_2^+ \quad \Gamma_3^+, [\mathbf{NT}, \mathbf{NTB}: A, B] \quad \Gamma_4^+} \wedge_t \mathbf{E} : \mathbf{NF}$$

$\Gamma_1^+, [T, \text{NT}, \text{TB}, \text{NTB}; B]$	$\Gamma_2^+, [T, \text{NT}; A, B]$	$\Gamma_3^+, [\text{NT}, \text{NTB}; A, B]$	$\Gamma_4^+, [T, \text{NT}, \text{TB}, \text{NTB}; A]$	$\Gamma_1^-, \dots, \Gamma_4^+, [\text{NT}; A \wedge_t B]$
$\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}; A \wedge_t B]]$	$\Gamma_1^-, [[\mathbf{FT}, \mathbf{FTB}; B]]$	$\Gamma_2^-, [[\mathbf{FT}; A, B]]$	$\Gamma_3^-, [[\mathbf{FT}, \mathbf{FTB}; A, B]]$	$\Gamma_4^-, [[\mathbf{FT}, \mathbf{FTB}; A]]$
$\Gamma_0^+, [T, \text{NT}, \text{TB}, \text{NTB}; A \wedge_t B]$	$\Gamma_1^+, [T, \text{TB}; B]$	$\Gamma_2^+, [T; A, B]$	$\Gamma_3^+$	$\Gamma_4^+, [T, \text{TB}; A]$
		$\Gamma_0^+, \dots, \Gamma_4^+$		$\wedge_{tE: \mathbf{FT}}$
$\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}; A \wedge_t B]]$	$\Gamma_1^-, [[\mathbf{B}, \mathbf{NB}; B]]$	$\Gamma_2^-, [[\mathbf{B}, \mathbf{NB}; A, B]]$	$\Gamma_3^-, [[\mathbf{NB}; A, B]]$	$\Gamma_4^-, [[\mathbf{B}, \mathbf{NB}; A]]$
$\Gamma_0^+, [T, \text{NT}, \text{TB}, \text{NTB}; A \wedge_t B]$	$\Gamma_1^+, [\mathbf{TB}, \text{NTB}; B]$	$\Gamma_2^+$	$\Gamma_3^+, [\text{NTB}; A, B]$	$\Gamma_4^+, [\mathbf{TB}, \text{NTB}; A]$
		$\Gamma_0^+, \dots, \Gamma_4^+$		$\wedge_{tE: \mathbf{NB}}$
$\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}; A \wedge_t B]]$	$\Gamma_1^-, [[\mathbf{B}, \mathbf{FB}, \mathbf{FTB}; B]]$	$\Gamma_2^-, [[\mathbf{B}, \mathbf{FB}; A, B]]$	$\Gamma_3^-, [[\mathbf{FB}, \mathbf{FTB}; A, B]]$	$\Gamma_4^-, [[\mathbf{B}, \mathbf{FB}, \mathbf{FTB}; A]]$
$\Gamma_0^+, [T, \text{NT}, \text{TB}, \text{NTB}; A \wedge_t B]$	$\Gamma_1^+, [\mathbf{TB}; B]$	$\Gamma_2^+$	$\Gamma_3^+$	$\Gamma_4^+, [\mathbf{TB}; A]$
		$\Gamma_0^+, \dots, \Gamma_4^+$		$\wedge_{tE: \mathbf{FB}}$
$\Gamma_1^-, \dots, \Gamma_4^+$	$\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-$	$\Gamma_1^+, \Gamma_2^+, \Gamma_3^+, \Gamma_4^+$	$\wedge_{tI: \mathbf{TB}}$	
$\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}; A \wedge_t B]]$	$\Gamma_1^-, [[\mathbf{FT}, \mathbf{NFT}, \mathbf{FTB}; A, B]]$	$\Gamma_2^-, [[\mathbf{FT}, \mathbf{NFT}; A, B]]$	$\Gamma_3^-, [[\mathbf{FT}, \mathbf{NFT}, \mathbf{FTB}; A; A, B]]$	$\Gamma_4^-, [[\mathbf{FT}, \mathbf{NFT}, \mathbf{FTB}; A; A]]$
$\Gamma_0^+, [T, \text{NT}, \text{TB}, \text{NTB}; A \wedge_t B]$	$\Gamma_1^+, [T, \text{NT}, \text{TB}, \text{NTB}; B]$	$\Gamma_2^+, [T, \text{NT}; A, B]$	$\Gamma_3^+, [\text{NT}, \text{NTB}; A, B]$	$\Gamma_4^+, [T, \text{NT}, \text{TB}, \text{NTB}; A]$
		$\Gamma_0^+, \dots, \Gamma_5^+$		$\wedge_{tE: \mathbf{NFT}}$
$\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}; A \wedge_t B]]$	$\Gamma_1^-, [[\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}; A, B]]$	$\Gamma_2^-, [[\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}; A, B]]$	$\Gamma_3^-, [[\mathbf{NB}, \mathbf{NFB}; A; A, B]]$	$\Gamma_4^-, [[\mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}; A; A, B]]$
$\Gamma_0^+, [T, \text{NT}, \text{TB}, \text{NTB}; A \wedge_t B]$	$\Gamma_1^+, [\mathbf{TB}, \text{NTB}; B]$	$\Gamma_2^+$	$\Gamma_3^+, [\text{NTB}; A, B]$	$\Gamma_4^+, [\mathbf{TB}, \text{NTB}; A]$
		$\Gamma_0^+, \dots, \Gamma_5^+$		$\wedge_{tE: \mathbf{NFB}}$
$\Gamma_1^-, \dots, \Gamma_4^+$	$\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-$	$\Gamma_1^+, \Gamma_2^+, \Gamma_3^+, \Gamma_4^+$	$\wedge_{tI: \mathbf{NTB}}$	
$\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}; A \wedge_t B]]$	$\Gamma_1^-, [[\mathbf{FTB}; B]]$	$\Gamma_2^-, [[\mathbf{FTB}; A, B]]$	$\Gamma_3^-, [[\mathbf{FTB}; A]]$	$\Gamma_4^-, [[\mathbf{FTB}; A]]$
$\Gamma_0^+, [T, \text{NT}, \text{TB}, \text{NTB}; A \wedge_t B]$	$\Gamma_1^+, [\mathbf{TB}; B]$	$\Gamma_2^+$	$\Gamma_3^+, [\mathbf{TB}; A]$	$\Gamma_4^+, [\mathbf{TB}, \text{NTB}; A]$
		$\Gamma_0^+, \dots, \Gamma_3^+$		$\wedge_{tE: \mathbf{FTB}}$
$\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}; A \wedge_t B]]$	$\Gamma_1^-, [[\mathbf{FTB}; A, B]]$	$\Gamma_2^-, [[\mathbf{A}; A, B]]$	$\Gamma_3^-, [[\mathbf{FTB}, \mathbf{A}; A, B]]$	$\Gamma_4^-, [[\mathbf{FTB}, \mathbf{A}; A]]$
$\Gamma_0^+, [T, \text{NT}, \text{TB}, \text{NTB}; A \wedge_t B]$	$\Gamma_1^+, [\mathbf{TB}, \text{NTB}; B]$	$\Gamma_2^+, [\text{NTB}; A, B]$	$\Gamma_3^+$	$\Gamma_4^+, [\mathbf{TB}, \text{NTB}; A]$
		$\Gamma_0^+, \dots, \Gamma_4^+$		$\wedge_{tE: \mathbf{A}}$

The introduction and elimination rules for connective  $\vee_t$  are given by

The introduction and elimination rules for connective  $\vee_t$  are given by

$$\frac{\Gamma_0^-, [[N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \vee_t B]] \quad \Gamma_1^-, [[N, NF: B]] \quad \Gamma_2^-, [[N: A, B]] \quad \Gamma_3^-, [[N, NF: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+} \quad \Gamma_4^-, [[N, N, F, NF: A]]$$

$$\frac{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B]}{\Gamma_0^+, \dots, \Gamma_4^+} \vee_t E: N$$

$$\frac{\Gamma_0^-, [[N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \vee_t B]] \quad \Gamma_1^-, [[N, NF: B]] \quad \Gamma_2^-, [[N: A, B]] \quad \Gamma_3^-, [[N, NF: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+} \quad \Gamma_4^-, [[N, N, F, NF: A]]$$

$$\frac{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B]}{\Gamma_0^+, \dots, \Gamma_3^+} \vee_t E: N$$



$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFT, NFB, FTB : A \vee_t B]] \quad \Gamma_1^-, [[NF, NFT, NFB, A : B]] \quad \Gamma_2^-, [[NFT, A : A, B]] \quad \Gamma_3^-, [[NFB, A : A, B]] \quad \Gamma_4^-, [[NF, NFT, NFB, A : A]]}{\Gamma_0^+, [T, NT, TB, NTB : A \vee_t B] \qquad \Gamma_1^+ \qquad \Gamma_2^+ \qquad \Gamma_3^+ \qquad \Gamma_4^+} \quad \forall_{t E : A}$$

The introduction and elimination rules for connective  $\wedge_f$  are given by

$$\frac{\Gamma_0^-, [[N, F, B, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, A : A \wedge_f B]] \quad \Gamma_1^-, [[N, F, B, \mathbf{FB} : B]] \quad \Gamma_2^-, [[N, F : A, B]] \quad \Gamma_3^-, [[N, B : A, B]] \quad \Gamma_4^-, [[N, F, B, \mathbf{FB} : A]]}{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB} : A \wedge_f B] \qquad \qquad \qquad \Gamma_1^+ \qquad \qquad \qquad \Gamma_2^+ \qquad \qquad \qquad \Gamma_3^+ \qquad \qquad \qquad \Gamma_4^+} \frac{}{\wedge_f E : \mathbf{N}} \\ \Gamma_0^+, \dots, \Gamma_4^+$$

$$\Gamma_0^+, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_f B]]] \quad \Gamma_1^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: B]]] \quad \Gamma_2^+, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}: A, B]]] \quad \Gamma_3^+, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A, B]]] \quad \Gamma_4^+, [[\mathbf{N}, \mathbf{NF}, \mathbf{NB}, \mathbf{NFB}: A, B]]] \quad \Gamma_5^+, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A]]] \\ \Gamma_6^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_f B]] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+ \quad \Gamma_5^+ \quad \Gamma_6^+ \quad \wedge_{f: E} \mathbf{N}$$

$$\frac{\Gamma_0^-, [[N, N, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[F, FB: B]] \quad \Gamma_2^-, [[F: A, B]] \quad \Gamma_3^-, [[F, FB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \qquad \qquad \Gamma_1^+ \qquad \qquad \Gamma_2^+ \qquad \qquad \Gamma_3^+} \frac{}{\Gamma_0^+, \dots, \Gamma_3^+} \wedge_f E: F$$

$$\frac{\Gamma_1^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: B]] \quad \Gamma_2^-, [[\mathbf{N}, \mathbf{F}, \mathbf{FT}: A, B]] \quad \Gamma_3^-, [[\mathbf{N}, \mathbf{B}: A, B]] \quad \Gamma_4^-, [[\mathbf{FT}, \mathbf{FTB}: A, B]] \quad \Gamma_5^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A]]}{\Gamma_1^+, [T, \mathbf{TB}: B] \qquad \qquad \Gamma_2^+, [T: A, B] \qquad \Gamma_3^+, [T, \mathbf{TB}: A, B] \qquad \Gamma_4^+, [T, \mathbf{TB}: A, B] \qquad \Gamma_5^+, [T, \mathbf{TB}: A]} \wedge_f I : T$$

$$\frac{\Gamma_0^-, [[N, N, F, NF, FT, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[B, FB: B]] \quad \Gamma_2^-, [[B: A, B]] \quad \Gamma_3^-, [[B, FB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \qquad \qquad \Gamma_1^+ \qquad \qquad \Gamma_2^+ \qquad \qquad \Gamma_3^+} \frac{}{\Gamma_0^+, \dots, \Gamma_3^+} \wedge_f E: B$$

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_f B]] \quad \Gamma_1^-, [[\mathbf{F}, \mathbf{NF}, \mathbf{FB}, \mathbf{NFB}: B]] \quad \Gamma_2^-, [[\mathbf{F}, \mathbf{NF}: A, B]] \quad \Gamma_3^-, [[\mathbf{NF}, \mathbf{NFB}: A, B]] \quad \Gamma_4^-, [[\mathbf{F}, \mathbf{NF}, \mathbf{FB}, \mathbf{NFB}: A]]}{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_f B] \qquad \qquad \qquad \Gamma_1^+ \qquad \qquad \qquad \Gamma_2^+ \qquad \qquad \qquad \Gamma_3^+ \qquad \qquad \qquad \Gamma_4^+} \wedge_{fE: \mathbf{NF}} \Gamma_0^+, \dots, \Gamma_4^+$$

$$\frac{\Gamma_1^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: A, B]]}{\Gamma_1^+, [T, \mathbf{NT}: A, B]} \quad \frac{\Gamma_2^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A, B]]}{\Gamma_2^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A, B]} \quad \frac{\Gamma_3^-, [[\mathbf{N}, \mathbf{NF}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: A, B]]}{\Gamma_3^+, [\mathbf{NT}, \mathbf{NTB}: A, B]} \quad \frac{\Gamma_4^-, [[\mathbf{FT}, \mathbf{NFT}, \mathbf{FTB}, \mathbf{A}: A, B]]}{\Gamma_4^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A, B]} \quad \wedge_f \mathbf{NT}$$

$$\Gamma_1^+, \dots, \Gamma_4^+, [\mathbf{NT}: A \wedge_f B]$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[F, FT, FB, FTB: B]] \quad \Gamma_2^-, [[F, FT: A, B]] \quad \Gamma_3^-, [[FT, FTB: A, B]] \quad \Gamma_4^-, [[F, FT, FB, FTB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+} \quad \frac{}{\wedge_f E: FT}$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[B, NB, FB, NFB: B]] \quad \Gamma_2^-, [[B, NB: A, B]] \quad \Gamma_3^-, [[NB, NFB: A, B]] \quad \Gamma_4^-, [[B, NB, FB, NFB: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+}{\Gamma_0^+, \dots, \Gamma_4^+} \wedge_{fE: NB}}$$

$$\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_f B]] \quad \Gamma_1^-, [[\mathbf{FB}: B]] \quad \Gamma_2^-, [[\mathbf{FB}: A]]}{\frac{\Gamma_0^+, [T, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_f B]}{\Gamma_0^+, \Gamma_1^+, \Gamma_2^+} \quad \frac{\Gamma_1^+}{\Gamma_1^+} \quad \frac{\Gamma_2^+}{\Gamma_2^+}} \quad \wedge_f^{E: \mathbf{FB}}$$

$$\frac{\Gamma_1^-, [[B, \mathbf{FB}, \mathbf{FTB}: B]] \quad \Gamma_2^-, [[B: A, B]] \quad \Gamma_3^-, [[\mathbf{FTB}: A, B]] \quad \Gamma_4^-, [[B, \mathbf{FB}, \mathbf{FTB}: A]]}{\Gamma_1^+, [\mathbf{TB}: B] \qquad \Gamma_2^+, [\mathbf{TB}: A, B] \qquad \Gamma_3^+, [\mathbf{TB}: A, B] \qquad \Gamma_4^+, [\mathbf{TB}: A]} \wedge_{f^{\mathbf{I}}: \mathbf{TB}} \Gamma_1^+, \dots, \Gamma_4^+, [\mathbf{TB}: A \wedge_f B]$$

$$\Gamma_0^+, [[\text{N}, \text{F}, \text{B}, \text{NF}, \text{FT}, \text{NB}, \text{FB}, \text{NFB}, \text{FTB}, \text{A}: A \wedge B]] \quad \Gamma_1^+, [[\text{F}, \text{NF}, \text{FT}, \text{FB}, \text{NFT}, \text{NFB}, \text{FTB}, \text{A}: B]] \quad \Gamma_2^+, [[\text{F}, \text{NF}, \text{FT}, \text{NFT}, \text{A}: B]] \quad \Gamma_3^+, [[\text{NF}, \text{NFT}, \text{NFB}, \text{A}: A, B]] \quad \Gamma_4^+, [[\text{FT}, \text{NFT}, \text{FTB}, \text{A}: B]] \quad \Gamma_5^+, [[\text{F}, \text{NF}, \text{FT}, \text{FB}, \text{NFT}, \text{NFB}, \text{FTB}, \text{A}: A]] \\ \Gamma_0^-, [\text{T}, \text{NT}, \text{TB}, \text{NTB}: A \wedge B] \quad \Gamma_1^-, \quad \Gamma_2^-, \quad \Gamma_3^-, \quad \Gamma_4^-, \quad \Gamma_5^- \quad \wedge_B \text{E: NFT}$$

$$\begin{array}{c}
\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFT, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[FB, NFB: B]] \quad \Gamma_2^-, [[NFB: A, B]] \quad \Gamma_3^-, [[FB, NFB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+} \wedge_f E: NFB \\
\\
\frac{\Gamma_1^-, [[B, NB, FB, NFB, FTB, A: B]] \quad \Gamma_2^-, [[B, NB: A, B]] \quad \Gamma_3^-, [[NB, NFB, A: A, B]] \quad \Gamma_4^-, [[FTB, A: A, B]] \quad \Gamma_5^-, [[B, NB, FB, NFB, FTB, A: A]]}{\Gamma_1^+, [TB, NTB: B] \quad \Gamma_2^+, [TB, NTB: A, B] \quad \Gamma_3^+, [NTB: A, B] \quad \Gamma_4^+, [TB, NTB: A, B] \quad \Gamma_5^+, [TB, NTB: A]} \wedge_f I: NTB \\
\\
\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[FB, FTB: B]] \quad \Gamma_2^-, [[FTB: A, B]] \quad \Gamma_3^-, [[FB, FTB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+} \wedge_f E: FTB \\
\\
\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[FB, NFB, FTB, A: B]] \quad \Gamma_2^-, [[NFB, A: A, B]] \quad \Gamma_3^-, [[FTB, A: A, B]] \quad \Gamma_4^-, [[FB, NFB, FTB, A: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+} \wedge_f E: A
\end{array}$$

The introduction and elimination rules for connective  $\vee_f$  are given by

$$\begin{array}{c}
\frac{\Gamma_0^-, [[N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[N, N: B]] \quad \Gamma_2^-, [[N, N: A, B]] \quad \Gamma_3^-, [[N: A, B]] \quad \Gamma_4^-, [[N, N: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [T, NT: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [T: A, B] \quad \Gamma_4^+, [T, NT: A]} \vee_f E: N \\
\\
\frac{\Gamma_0^-, [[N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[N: B]] \quad \Gamma_2^-, [[N: A, B]] \quad \Gamma_3^-, [[N: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [NT: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [NT: A]} \vee_f E: N \\
\\
\frac{\Gamma_0^-, [[N, N, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[N, N, F, NF, FT, NFT: B]] \quad \Gamma_2^-, [[N, N, F, NF: A, B]] \quad \Gamma_3^-, [[N, F, FT: A, B]] \quad \Gamma_4^-, [[F, NF, FT, NFT: A, B]] \quad \Gamma_5^-, [[N, N, F, NF, FT, NFT: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [T, NT: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [T: A, B] \quad \Gamma_4^+, [T, TB: A] \quad \Gamma_5^+, [T, NT: A]} \vee_f E: F \\
\\
\frac{\Gamma_1^- \quad \Gamma_2^- \quad \Gamma_3^-}{\Gamma_1^+, [T, NT: B] \quad \Gamma_2^+, [T: A, B] \quad \Gamma_3^+, [T, NT: A]} \vee_f I: T \\
\\
\frac{\Gamma_0^-, [[N, N, F, NF, FT, NB, FB, NFT, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[N, N, B, NB: B]] \quad \Gamma_2^-, [[N, N, B, NB: A, B]] \quad \Gamma_3^-, [[N, B: A, B]] \quad \Gamma_4^-, [[B, NB: A, B]] \quad \Gamma_5^-, [[N, N, B, NB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [T, NT, TB, NTB: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [T, TB: A, B] \quad \Gamma_4^+, [TB, NTB: A, B] \quad \Gamma_5^+, [T, NT, TB, NTB: A]} \vee_f E: B \\
\\
\frac{\Gamma_0^-, [[N, N, F, B, FT, NB, FB, NFT, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[N, NF, NFT: B]] \quad \Gamma_2^-, [[N, NF: A, B]] \quad \Gamma_3^-, [[NF, NFT: A, B]] \quad \Gamma_4^-, [[N, NF, NFT: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [NT: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [NT: A] \quad \Gamma_4^+, [NT: A]} \vee_f E: NF \\
\\
\frac{\Gamma_1^- \quad \Gamma_2^-}{\Gamma_1^+, [NT: B] \quad \Gamma_2^+, [NT: A]} \vee_f I: NT
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma_0^-, [[N, N, F, B, NF, NB, FB, NFT, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[FT, NFT: B]] \quad \Gamma_2^-, [[FT: A, B]] \quad \Gamma_3^-, [[FT, NFT: A, B]] \quad \Gamma_4^-, [[FT, NFT: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [T, NT: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [T: A, B] \quad \Gamma_4^+, [T, NT: A]} \vee_f E: FT \\
\\
\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, FB, NFT, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[N, NB: B]] \quad \Gamma_2^-, [[N, NB: A, B]] \quad \Gamma_3^-, [[NB: A, B]] \quad \Gamma_4^-, [[N, NB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [NT, NTB: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [NTB: A, B] \quad \Gamma_4^+, [NT, NTB: A]} \vee_f E: NB
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, NFT, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[N, N, F, B, NF, NB, FB, NFB: A, B]] \quad \Gamma_2^-, [[N, F, B, FT, FB, FTB: A, B]] \quad \Gamma_3^-, [[F, NF, FT, FB, NFT, NFB, FTB, A: A, B]] \quad \Gamma_4^-, [[B, NB, FB, NFB, FTB, A: A, B]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_2^+, [T, TB: A, B] \quad \Gamma_3^+, [TB, NTB: A, B] \quad \Gamma_4^+, [TB, NTB: A, B]} \quad \frac{}{\Gamma_0^+, \dots, \Gamma_4^+} \quad \frac{}{\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-} \\
\frac{\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-}{\Gamma_1^+, [T, NT, TB, NTB: B] \quad \Gamma_2^+, [T, TB: A, B] \quad \Gamma_3^+, [TB, NTB: A, B] \quad \Gamma_4^+, [T, NT, TB, NTB: A]} \quad \frac{}{\Gamma_1^+, \dots, \Gamma_4^+} \quad \frac{}{\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-} \\
\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[NFT: B]] \quad \Gamma_2^-, [[NFT: A, B]] \quad \Gamma_3^-, [[NFT: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [NT: B] \quad \Gamma_2^+, [NT: A]} \quad \frac{}{\Gamma_0^+, \dots, \Gamma_3^+} \quad \frac{}{\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-} \\
\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[N, NF, NB, NFT, NFB: A, B]] \quad \Gamma_2^-, [[N, NF, NB, NFB: A, B]] \quad \Gamma_3^-, [[NF, NFT, NFB, A: A, B]] \quad \Gamma_4^-, [[NB, NFB, A: A, B]] \quad \Gamma_5^-, [[N, NF, NB, NFT, NFB, A: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [NT, NTB: B] \quad \Gamma_2^+, [NTB: A] \quad \Gamma_3^+, [NT: B] \quad \Gamma_4^+, [NTB: A] \quad \Gamma_5^+, [NT, NTB: A]} \quad \frac{}{\Gamma_0^+, \dots, \Gamma_5^+} \quad \frac{}{\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-, \Gamma_5^-} \\
\frac{\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-, \Gamma_5^-}{\Gamma_1^+, [NT, NTB: B] \quad \Gamma_2^+, [NTB: A, B] \quad \Gamma_3^+, [NT, NTB: A]} \quad \frac{}{\Gamma_1^+, \Gamma_2^+, \Gamma_3^+, [NTB: A \vee_f B]} \quad \frac{}{\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-, \Gamma_5^-} \\
\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFT, NFB, FTB: A \vee_f B]] \quad \Gamma_1^-, [[FT, NFT, FTB, A: B]] \quad \Gamma_2^-, [[FT, FTB: A, B]] \quad \Gamma_3^-, [[FT, NFT, FTB, A: A, B]] \quad \Gamma_4^-, [[FTB, A: A, B]] \quad \Gamma_5^-, [[FT, NFT, FTB, A: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [T, NT, TB, NTB: B] \quad \Gamma_2^+, [T, TB: A, B] \quad \Gamma_3^+, [TB, NTB: A, B] \quad \Gamma_4^+, [TB, NTB: A, B] \quad \Gamma_5^+, [T, NT, TB, NTB: A]} \quad \frac{}{\Gamma_0^+, \dots, \Gamma_5^+} \quad \frac{}{\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-, \Gamma_5^-} \\
\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFT, NFB, FTB: A \vee_f B]] \quad \Gamma_1^-, [[NFT, A: B]] \quad \Gamma_2^-, [[NFT, A: A, B]] \quad \Gamma_3^-, [[A: A, B]] \quad \Gamma_4^-, [[NFT, A: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [NT, NTB: B] \quad \Gamma_2^+, [NTB: A, B] \quad \Gamma_3^+, [NTB: A, B] \quad \Gamma_4^+, [NT, NTB: A]} \quad \frac{}{\Gamma_0^+, \dots, \Gamma_4^+} \quad \frac{}{\Gamma_1^-, \Gamma_2^-, \Gamma_3^-, \Gamma_4^-, \Gamma_5^-}
\end{array}$$

**Definition 22.** A *natural deduction derivation* is defined inductively as follows:

1. Let  $A$  be any formula. Then

$$\frac{[V^-: A]}{[V^+: A]}$$

is a derivation of  $A$  from the assumption  $[V^-: A]$  (an *initial derivation*).

2. If  $D_k$  are derivations of  $\Gamma_k^+, \Delta_k^+$  from the assumptions  $\Gamma_k^-, \hat{\Delta}_k^-$ , and

$$\frac{\left\langle \begin{array}{c} \Gamma_k^-, [\Delta_k^-] \\ \Gamma_k^+, \Delta_k^+ \end{array} \right\rangle_{k \in K}}{\Pi^+}$$

is an instance of a deduction rule with  $\hat{\Delta}_k^-$  a subsequent of  $\Delta_k^-$ , and all eigenvariable conditions are satisfied, then

$$\frac{\langle D_k \rangle_{k \in K}}{\Pi^+}$$

is a derivation of  $\Pi^+$  from the assumptions  $\bigcup_{k \in K} \Gamma_k^-$ . The formulas in  $\hat{\Delta}_k^-$  which do not occur in  $\bigcup_{k \in K} \Gamma_k^-$  are said to be *discharged* at this inference.

**Theorem 23.** A partial sequent  $\Gamma^+$  can be derived from the assumptions  $\Gamma^-$  in the natural deduction system for Shramko-Wansing logic iff, for every interpretation  $\mathfrak{I}$ , either some formula in  $\Gamma_v^-$  ( $v \in V^-$ ) evaluates to the truth value  $v$ , or there is a  $w \in V^+$  and a formula in  $\Gamma_w^+$  that evaluates to  $w$ .

*Proof.* See Theorems 4.7 and 5.4 of Baaz et al. [3] or Theorems 4.2.8 and 4.3.4 of Zach [15].  $\square$

**Corollary 24.**  $\Gamma \models A$  iff there is a natural deduction derivation of  $[V^+ : A]$  from  $[V^- : \Gamma]$ .

## 6 Resolution and clause formation rules for Shramko-Wansing logic

The many-valued resolution calculus of Baaz and Fermüller [1] applies to Shramko-Wansing logic. We present the framework here, as well as a language preserving clause translation system for Shramko-Wansing logic.

**Definition 25** (Signed formula). A *signed formula* is an expression of the form  $A^v$ , where  $A$  is a formula and  $v \in V$ . If  $A$  is a propositional variable,  $A^v$  is a *signed atom*.

**Definition 26** (Signed clause). A (signed) *clause*  $C = \{A_1^{v_1}, \dots, A_n^{v_n}\}$  is a finite set of signed atoms (proper clause). The empty clause is denoted by  $\square$ .

An *extended clause* is a finite set of signed formulas.

**Definition 27** (Semantics of clause sets). Let  $\mathfrak{I}$  be an interpretation.  $\mathfrak{I}$  satisfies a clause  $C$  iff there is some signed formula  $A^v \in C$ , so that  $\text{val}_{\mathfrak{I}}(A) = v$ .  $\mathfrak{I}$  satisfies a clause set  $\mathcal{C}$  iff it satisfies every clause in  $\mathcal{C}$ .  $\mathcal{C}$  is called *satisfiable* iff some structure satisfies it, and *unsatisfiable* otherwise.

The clause formation rules for connective  $\neg_t$  are given by

$$\begin{array}{ll}
 \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^N\}\}}{\mathcal{C} \cup \{C \cup \{A^N\}\}} \quad \neg_t:N & \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^N\}\}}{\mathcal{C} \cup \{C \cup \{A^T\}\}} \quad \neg_t:N \\
 \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^F\}\}}{\mathcal{C} \cup \{C \cup \{A^B\}\}} \quad \neg_t:F & \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^T\}\}}{\mathcal{C} \cup \{C \cup \{A^N\}\}} \quad \neg_t:T \\
 \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^B\}\}}{\mathcal{C} \cup \{C \cup \{A^F\}\}} \quad \neg_t:B & \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{NF}\}\}}{\mathcal{C} \cup \{C \cup \{A^{TB}\}\}} \quad \neg_t:NF \\
 \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{NT}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NT}\}\}} \quad \neg_t:NT & \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{FT}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NB}\}\}} \quad \neg_t:FT \\
 \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{NB}\}\}}{\mathcal{C} \cup \{C \cup \{A^{FT}\}\}} \quad \neg_t:NB & \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{FB}\}\}}{\mathcal{C} \cup \{C \cup \{A^{FB}\}\}} \quad \neg_t:FB
 \end{array}$$

$$\begin{array}{ll}
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{TB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NF}}\}\}} \neg_t : \mathbf{TB} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{NFT}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NTB}}\}\}} \neg_t : \mathbf{NFT} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{NFB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{FTB}}\}\}} \neg_t : \mathbf{NFB} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{NTB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NFT}}\}\}} \neg_t : \mathbf{NTB} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{FTB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NFB}}\}\}} \neg_t : \mathbf{FTB} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{A}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{A}}\}\}} \neg_t : \mathbf{A}
\end{array}$$

The clause formation rules for connective  $\neg_f$  are given by

$$\begin{array}{ll}
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{N}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{N}}\}\}} \neg_f : \mathbf{N} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{N}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{F}}\}\}} \neg_f : \mathbf{N} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{F}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{N}}\}\}} \neg_f : \mathbf{F} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{T}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{B}}\}\}} \neg_f : \mathbf{T} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{B}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{T}}\}\}} \neg_f : \mathbf{B} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{NF}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NF}}\}\}} \neg_f : \mathbf{NF} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{NT}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{FB}}\}\}} \neg_f : \mathbf{NT} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{FT}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NB}}\}\}} \neg_f : \mathbf{FT} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{NB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{FT}}\}\}} \neg_f : \mathbf{NB} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{FB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NT}}\}\}} \neg_f : \mathbf{FB} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{TB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{TB}}\}\}} \neg_f : \mathbf{TB} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{NFT}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NFB}}\}\}} \neg_f : \mathbf{NFT} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{NFB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NFT}}\}\}} \neg_f : \mathbf{NFB} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{NTB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{FTB}}\}\}} \neg_f : \mathbf{NTB} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{FTB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NTB}}\}\}} \neg_f : \mathbf{FTB} & \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{\mathbf{A}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{A}}\}\}} \neg_f : \mathbf{A}
\end{array}$$

The clause formation rules for connective  $\wedge_t$  are given by

$$\begin{array}{c}
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\mathbf{N}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\mathbf{N}}, B^{\mathbf{T}}, B^{\mathbf{B}}, B^{\mathbf{TB}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{T}}, B^{\mathbf{N}}, B^{\mathbf{T}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{B}}, B^{\mathbf{N}}, B^{\mathbf{B}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{T}}, A^{\mathbf{B}}, A^{\mathbf{TB}}\}\}} \wedge_t : \mathbf{N} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\mathbf{F}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\mathbf{N}}, B^{\mathbf{F}}, B^{\mathbf{B}}, B^{\mathbf{TB}}, B^{\mathbf{FT}}, B^{\mathbf{FB}}, B^{\mathbf{NFB}}, B^{\mathbf{NTB}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{F}}, A^{\mathbf{T}}, A^{\mathbf{NT}}, B^{\mathbf{N}}, B^{\mathbf{N}}, B^{\mathbf{F}}, B^{\mathbf{NT}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{N}}, A^{\mathbf{B}}, A^{\mathbf{NB}}, B^{\mathbf{N}}, B^{\mathbf{N}}, B^{\mathbf{B}}, B^{\mathbf{NFB}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{N}}, A^{\mathbf{NT}}, A^{\mathbf{NB}}, A^{\mathbf{NTB}}, B^{\mathbf{N}}, B^{\mathbf{NT}}, B^{\mathbf{NFB}}, B^{\mathbf{NTB}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{N}}, A^{\mathbf{T}}, A^{\mathbf{B}}, A^{\mathbf{NT}}, A^{\mathbf{NB}}, A^{\mathbf{TB}}, A^{\mathbf{NTB}}\}\}} \wedge_t : \mathbf{F} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\mathbf{T}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\mathbf{T}}, B^{\mathbf{TB}}\}, C \cup \{A^{\mathbf{T}}, B^{\mathbf{T}}\}, C \cup \{A^{\mathbf{T}}, A^{\mathbf{TB}}\}\}} \wedge_t : \mathbf{T}
\end{array}$$

$$\begin{aligned}
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^B\}\}}{\mathcal{C} \cup \{C \cup \{B^B, B^{\text{TB}}\}, C \cup \{A^B, B^B\}, C \cup \{A^B, A^{\text{TB}}\}\}} \wedge_t: \mathbf{B} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\text{NT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^T, B^{\text{NT}}, B^{\text{TB}}, B^{\text{NTB}}\}, C \cup \{A^T, A^{\text{NT}}, B^T, B^{\text{NT}}\}, C \cup \{A^{\text{NT}}, A^{\text{NTB}}, B^{\text{NT}}, B^{\text{NTB}}\}, C \cup \{A^T, A^{\text{NT}}, A^{\text{NTB}}, B^{\text{NT}}, B^{\text{NTB}}, B^{\text{NT}}, B^{\text{NT}}, B^{\text{NT}}, B^{\text{NT}}, B^{\text{NT}}, B^{\text{NT}}, B^{\text{NT}}\}\}} \wedge_t: \text{NT} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\text{FT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^T, B^{\text{FT}}, B^{\text{TB}}, B^{\text{FTB}}\}, C \cup \{A^T, A^{\text{FT}}, B^T, B^{\text{FT}}\}, C \cup \{A^{\text{FT}}, A^{\text{FTB}}, B^{\text{FT}}, B^{\text{FTB}}\}, C \cup \{A^T, A^{\text{FT}}, A^{\text{TB}}, A^{\text{FTB}}\}\}} \wedge_t: \text{FT} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\text{NB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^B, B^{\text{NB}}, B^{\text{TB}}, B^{\text{NTB}}\}, C \cup \{A^B, A^{\text{NB}}, B^B, B^{\text{NB}}\}, C \cup \{A^{\text{NB}}, A^{\text{NTB}}, B^{\text{NB}}, B^{\text{NTB}}\}, C \cup \{A^B, A^{\text{NB}}, A^{\text{TB}}, A^{\text{NTB}}\}\}} \wedge_t: \text{NB} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\text{FB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^B, B^{\text{FB}}, B^{\text{TB}}, B^{\text{FTB}}\}, C \cup \{A^B, A^{\text{FB}}, B^B, B^{\text{FB}}\}, C \cup \{A^{\text{FB}}, A^{\text{FTB}}, B^{\text{FB}}, B^{\text{FTB}}\}, C \cup \{A^B, A^{\text{FB}}, A^{\text{TB}}, A^{\text{FTB}}\}\}} \wedge_t: \text{FB} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\text{TB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{TB}}\}, C \cup \{A^{\text{TB}}\}\}} \wedge_t: \text{TB} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\text{NFT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^T, B^{\text{NFT}}, B^{\text{TB}}, B^{\text{NTB}}, B^{\text{FTB}}, B^{\text{NTB}}, B^{\text{FTB}}, B^{\text{NFT}}, B^{\text{TB}}, B^{\text{A}}\}, C \cup \{A^T, A^{\text{NFT}}, A^{\text{FT}}, A^{\text{NFT}}, B^T, B^{\text{NFT}}, B^{\text{FT}}, B^{\text{NFT}}\}, C \cup \{A^{\text{NFT}}, A^{\text{NTB}}, A^{\text{A}}, B^{\text{NFT}}, B^{\text{NFT}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^{\text{FT}}, A^{\text{NFT}}, A^{\text{FTB}}, A^{\text{A}}, B^{\text{FT}}, B^{\text{NFT}}, B^{\text{FTB}}, B^{\text{A}}\}, C \cup \{A^T, A^{\text{NFT}}, A^{\text{FT}}, A^{\text{TB}}, A^{\text{NFT}}, A^{\text{NTB}}, A^{\text{FTB}}, A^{\text{A}}\}\}} \wedge_t: \text{NFT} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\text{NFT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^B, B^{\text{NB}}, B^{\text{FB}}, B^{\text{NFB}}, B^{\text{NTB}}, B^{\text{FTB}}, B^{\text{A}}\}, C \cup \{A^B, A^{\text{NB}}, A^{\text{FB}}, A^{\text{NFB}}, B^B, B^{\text{NB}}, B^{\text{FB}}, B^{\text{NFB}}\}, C \cup \{A^{\text{NB}}, A^{\text{NFB}}, A^{\text{NTB}}, A^{\text{A}}, B^{\text{NB}}, B^{\text{NFB}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^{\text{FB}}, A^{\text{NFB}}, A^{\text{FTB}}, A^{\text{A}}, B^{\text{FB}}, B^{\text{NFB}}, B^{\text{FTB}}, B^{\text{A}}\}, C \cup \{A^B, A^{\text{NB}}, A^{\text{FB}}, A^{\text{TB}}, A^{\text{NFB}}, A^{\text{NTB}}, A^{\text{FTB}}, A^{\text{A}}\}\}} \wedge_t: \text{NFTB} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\text{NTB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{TB}}, B^{\text{NTB}}\}, C \cup \{A^{\text{NTB}}, B^{\text{NTB}}\}, C \cup \{A^{\text{TB}}, A^{\text{NTB}}\}\}} \wedge_t: \text{NTB} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\text{FTB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{TB}}, B^{\text{FTB}}\}, C \cup \{A^{\text{FTB}}, B^{\text{FTB}}\}, C \cup \{A^{\text{TB}}, A^{\text{FTB}}\}\}} \wedge_t: \text{FTB} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^{\text{A}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{TB}}, B^{\text{NTB}}, B^{\text{FTB}}, B^{\text{A}}\}, C \cup \{A^{\text{NTB}}, A^{\text{A}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^{\text{FTB}}, A^{\text{A}}, B^{\text{FTB}}, B^{\text{A}}\}, C \cup \{A^{\text{TB}}, A^{\text{NTB}}, A^{\text{FTB}}, A^{\text{A}}\}\}} \wedge_t: \text{A} \\
\end{aligned}$$

The clause formation rules for connective  $\vee_t$  are given by

$$\begin{aligned}
& \frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\mathbf{N}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\mathbf{N}}, B^{\mathbf{N}}, B^F, B^{\text{NF}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{N}}, B^{\mathbf{N}}, B^{\mathbf{N}}\}, C \cup \{A^{\mathbf{N}}, A^F, B^{\mathbf{N}}, B^F\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{N}}, A^F, A^{\text{NF}}\}\}} \vee_t: \mathbf{N} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\mathbf{N}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\mathbf{N}}, B^{\text{NF}}\}, C \cup \{A^{\mathbf{N}}, B^{\mathbf{N}}\}, C \cup \{A^{\mathbf{N}}, A^{\text{NF}}\}\}} \vee_t: \mathbf{N} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\mathbf{F}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\mathbf{F}}, B^{\text{NF}}\}, C \cup \{A^{\mathbf{F}}, B^{\mathbf{F}}\}, C \cup \{A^{\mathbf{F}}, A^{\text{NF}}\}\}} \vee_t: \mathbf{F} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\mathbf{T}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\mathbf{N}}, B^{\mathbf{F}}, B^{\mathbf{T}}, B^{\mathbf{NF}}, B^{\mathbf{NT}}, B^{\mathbf{FT}}, B^{\mathbf{NFT}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{N}}, A^{\mathbf{F}}, A^{\mathbf{N}^{\mathbf{T}}}, B^{\mathbf{N}}, B^{\mathbf{N}}, B^{\mathbf{T}}, B^{\mathbf{N}^{\mathbf{T}}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{F}}, A^{\mathbf{T}}, A^{\mathbf{N}^{\mathbf{T}}}, B^{\mathbf{N}}, B^{\mathbf{F}}, B^{\mathbf{T}}, B^{\mathbf{F}^{\mathbf{T}}}\}, C \cup \{A^{\mathbf{T}}, A^{\mathbf{N}^{\mathbf{T}}}, A^{\mathbf{F}^{\mathbf{T}}}, A^{\mathbf{N}^{\mathbf{FT}}}, B^{\mathbf{T}}, B^{\mathbf{N}^{\mathbf{T}}}, B^{\mathbf{F}^{\mathbf{T}}}, B^{\mathbf{N}^{\mathbf{FT}}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{N}}, A^{\mathbf{F}}, A^{\mathbf{T}}, A^{\mathbf{N}^{\mathbf{F}}}, A^{\mathbf{N}^{\mathbf{T}}}, A^{\mathbf{F}^{\mathbf{T}}}, A^{\mathbf{N}^{\mathbf{FT}}}\}\}} \vee_t: \mathbf{T} \\
& \frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\mathbf{B}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\mathbf{N}}, B^{\mathbf{F}}, B^{\mathbf{B}}, B^{\mathbf{NF}}, B^{\mathbf{NB}}, B^{\mathbf{FB}}, B^{\mathbf{NFB}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{N}}, A^{\mathbf{B}}, A^{\mathbf{NB}}, B^{\mathbf{N}}, B^{\mathbf{N}}, B^{\mathbf{B}}, B^{\mathbf{N}^{\mathbf{B}}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{F}}, A^{\mathbf{B}}, A^{\mathbf{FB}}, B^{\mathbf{N}}, B^{\mathbf{F}}, B^{\mathbf{B}}, B^{\mathbf{FB}}\}, C \cup \{A^{\mathbf{B}}, A^{\mathbf{N}^{\mathbf{B}}}, A^{\mathbf{FB}}, A^{\mathbf{NFB}}, B^{\mathbf{B}}, B^{\mathbf{N}^{\mathbf{B}}}, B^{\mathbf{F}}\}, C \cup \{A^{\mathbf{N}}, A^{\mathbf{N}}, A^{\mathbf{F}}, A^{\mathbf{B}}, A^{\mathbf{N}^{\mathbf{F}}}, A^{\mathbf{B}}\}\}} \vee_t: \mathbf{B}
\end{aligned}$$

$$\begin{array}{c}
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\text{NF}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{NF}}\}, C \cup \{A^{\text{NF}}\}\}} \vee_t:\text{NF} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\text{NT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^N, B^{\text{NF}}, B^{\text{NT}}, B^{\text{NFT}}\}, C \cup \{A^N, A^{\text{NT}}, B^N, B^{\text{NT}}\}, C \cup \{A^{\text{NT}}, A^{\text{NFT}}, B^{\text{NT}}, B^{\text{NFT}}\}, C \cup \{A^N, A^{\text{NF}}, A^{\text{NT}}, A^{\text{NFT}}\}\}} \vee_t:\text{NT} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\text{FT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^F, B^{\text{NF}}, B^{\text{FT}}, B^{\text{NFT}}\}, C \cup \{A^F, A^{\text{FT}}, B^F, B^{\text{FT}}\}, C \cup \{A^{\text{FT}}, A^{\text{NFT}}, B^{\text{FT}}, B^{\text{NFT}}\}, C \cup \{A^F, A^{\text{NF}}, A^{\text{FT}}, A^{\text{NFT}}\}\}} \vee_t:\text{FT} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\text{NB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^N, B^{\text{NF}}, B^{\text{NB}}, B^{\text{NFB}}\}, C \cup \{A^N, A^{\text{NB}}, B^N, B^{\text{NB}}\}, C \cup \{A^{\text{NB}}, A^{\text{NFB}}, B^{\text{NB}}, B^{\text{NFB}}\}, C \cup \{A^N, A^{\text{NF}}, A^{\text{NB}}, A^{\text{NFB}}\}\}} \vee_t:\text{NB} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\text{FB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^F, B^{\text{NF}}, B^{\text{FB}}, B^{\text{NFB}}\}, C \cup \{A^F, A^{\text{FB}}, B^F, B^{\text{FB}}\}, C \cup \{A^{\text{FB}}, A^{\text{NFB}}, B^{\text{FB}}, B^{\text{NFB}}\}, C \cup \{A^F, A^{\text{NF}}, A^{\text{FB}}, A^{\text{NFB}}\}\}} \vee_t:\text{FB} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\text{NFT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{NF}}, B^{\text{NFT}}\}, C \cup \{A^{\text{NFT}}, B^{\text{NFT}}\}, C \cup \{A^{\text{NF}}, A^{\text{NFT}}\}\}} \vee_t:\text{NFT} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\text{NFB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{NF}}, B^{\text{NFB}}\}, C \cup \{A^{\text{NFB}}, B^{\text{NFB}}\}, C \cup \{A^{\text{NF}}, A^{\text{NFB}}\}\}} \vee_t:\text{NFB} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\text{NTB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^N, B^{\text{NF}}, B^{\text{NT}}, B^{\text{NB}}, B^{\text{NFT}}, B^{\text{NFB}}, B^{\text{NTB}}, B^A\}, C \cup \{A^N, A^{\text{NT}}, A^{\text{NB}}, A^{\text{NFT}}, B^N, B^{\text{NT}}, B^{\text{NB}}, B^{\text{NFT}}, B^{\text{NTB}}, B^A\}, C \cup \{A^{\text{NT}}, A^{\text{NFT}}, A^{\text{NB}}, A^{\text{NFB}}, B^{\text{NT}}, B^{\text{NB}}, B^{\text{NFB}}, B^{\text{NTB}}, B^A\}, C \cup \{A^N, A^{\text{NF}}, A^{\text{NT}}, A^{\text{NB}}, A^{\text{NFB}}, A^{\text{NTB}}, A^A\}, C \cup \{A^N, A^{\text{NF}}, A^{\text{NT}}, A^{\text{NB}}, A^{\text{NFB}}, A^{\text{NTB}}, A^A\}\}} \vee_t:\text{NTB} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\text{FTB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^F, B^{\text{NF}}, B^{\text{FT}}, B^{\text{FB}}, B^{\text{NFT}}, B^{\text{NFB}}, B^{\text{FTB}}, B^A\}, C \cup \{A^F, A^{\text{FT}}, A^{\text{FB}}, A^{\text{FTB}}, B^F, B^{\text{FT}}, B^{\text{FB}}, B^{\text{NFTB}}, B^A\}, C \cup \{A^{\text{FT}}, A^{\text{NFT}}, A^{\text{FTB}}, A^{\text{FB}}, B^{\text{FT}}, B^{\text{FB}}, B^{\text{NFTB}}, B^A\}, C \cup \{A^{\text{FT}}, A^{\text{NFT}}, A^{\text{FTB}}, A^{\text{FB}}, B^{\text{FT}}, B^{\text{FB}}, B^{\text{NFTB}}, B^A\}, C \cup \{A^F, A^{\text{NF}}, A^{\text{FT}}, A^{\text{FB}}, A^{\text{NFT}}, A^{\text{FTB}}, A^A\}\}} \vee_t:\text{FTB} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_t B)^{\text{A}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{NF}}, B^{\text{NFT}}, B^{\text{NFB}}, B^A\}, C \cup \{A^{\text{NFT}}, A^{\text{A}}, B^{\text{NFT}}, B^A\}, C \cup \{A^{\text{NFB}}, A^{\text{A}}, B^{\text{NFB}}, B^A\}, C \cup \{A^{\text{NF}}, A^{\text{NFT}}, A^{\text{NFB}}, A^{\text{A}}\}\}} \vee_t:\text{A} \\
\end{array}$$

The clause formation rules for connective  $\wedge_f$  are given by

$$\begin{array}{c}
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{N}}\}\}}{\mathcal{C} \cup \{C \cup \{B^N, B^F, B^B, B^{\text{FB}}\}, C \cup \{A^N, A^F, B^N, B^F\}, C \cup \{A^{\text{N}}, A^B, B^N, B^B\}, C \cup \{A^N, A^F, A^B, A^{\text{FB}}\}\}} \wedge_f:\text{N} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{F}}\}\}}{\mathcal{C} \cup \{C \cup \{B^F, B^{\text{FB}}\}, C \cup \{A^F, B^F\}, C \cup \{A^F, A^{\text{FB}}\}\}} \wedge_f:\text{F} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{B}}\}\}}{\mathcal{C} \cup \{C \cup \{B^B, B^{\text{FB}}\}, C \cup \{A^B, B^B\}, C \cup \{A^B, A^{\text{FB}}\}\}} \wedge_f:\text{B} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{NF}}\}\}}{\mathcal{C} \cup \{C \cup \{B^F, B^{\text{NF}}, B^{\text{FB}}, B^{\text{NFB}}\}, C \cup \{A^F, A^{\text{NF}}, B^F, B^{\text{NF}}\}, C \cup \{A^{\text{NF}}, A^{\text{NFB}}, B^{\text{NF}}, B^{\text{NFB}}\}, C \cup \{A^F, A^{\text{NF}}, A^{\text{FB}}, A^{\text{NFB}}\}\}} \wedge_f:\text{NF} \\
\\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{NT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^N, A^N, A^F, A^T, A^{\text{NF}}, A^{\text{NT}}, A^{\text{FT}}, A^{\text{NFT}}, A^{\text{FTB}}, A^{\text{NFTB}}, B^N, B^A, B^F, B^{\text{FT}}, B^{\text{NFT}}, B^{\text{FTB}}, B^{\text{NFTB}}\}, C \cup \{A^N, A^N, A^F, A^T, A^{\text{FT}}, A^{\text{NFT}}, B^N, B^F, B^{\text{FT}}, B^{\text{FTB}}\}, C \cup \{A^N, A^T, A^{\text{FT}}, A^{\text{NFT}}, B^N, B^F, B^{\text{FT}}, B^{\text{NFT}}, B^{\text{FTB}}, B^{\text{NFTB}}\}, C \cup \{A^T, A^{\text{FT}}, A^{\text{NFT}}, A^{\text{FTB}}, A^{\text{NFTB}}, B^T, B^{\text{FT}}, B^{\text{NFT}}, B^{\text{FTB}}, B^{\text{NFTB}}\}, C \cup \{A^N, A^F, A^T, A^B, A^{\text{FT}}, A^{\text{NFT}}, A^{\text{FTB}}, A^{\text{NFTB}}, B^A, B^T, B^{\text{FT}}, B^{\text{NFT}}, B^{\text{FTB}}, B^{\text{NFTB}}\}\}} \wedge_f:\text{NT}
\end{array}$$

$$\begin{array}{c}
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{FT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^F, B^{\text{FT}}, B^{\text{FB}}, B^{\text{FTB}}\}, C \cup \{A^F, A^{\text{FT}}, B^F, B^{\text{FT}}\}, C \cup \{A^{\text{FT}}, A^{\text{FTB}}, B^{\text{FT}}, B^{\text{FTB}}\}, C \cup \{A^F, A^{\text{FT}}, A^{\text{FB}}, A^{\text{FTB}}\}\}} \wedge_f:\text{FT} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{NB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^B, B^{\text{NB}}, B^{\text{FB}}, B^{\text{NFB}}\}, C \cup \{A^B, A^{\text{NB}}, B^B, B^{\text{NB}}\}, C \cup \{A^{\text{NB}}, A^{\text{NFB}}, B^{\text{NB}}, B^{\text{NFB}}\}, C \cup \{A^B, A^{\text{NB}}, A^{\text{FB}}, A^{\text{NFB}}\}\}} \wedge_f:\text{NB} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{FB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{FB}}\}, C \cup \{A^{\text{FB}}\}\}} \wedge_f:\text{FB} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{TB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^B, B^{\text{FB}}, B^{\text{TB}}, B^{\text{FTB}}\}, C \cup \{A^B, A^{\text{TB}}, B^B, B^{\text{TB}}\}, C \cup \{A^{\text{TB}}, A^{\text{FTB}}, B^{\text{TB}}, B^{\text{FTB}}\}, C \cup \{A^B, A^{\text{FB}}, A^{\text{TB}}, A^{\text{FTB}}\}\}} \wedge_f:\text{TB} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{NFT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^F, B^{\text{NF}}, B^{\text{FT}}, B^{\text{FB}}, B^{\text{NFT}}, B^{\text{NFB}}, B^{\text{FTB}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^F, A^{\text{NF}}, A^{\text{FT}}, A^{\text{NFT}}, B^F, B^{\text{NF}}, B^{\text{FT}}, B^{\text{NFT}}\}, C \cup \{A^{\text{NF}}, A^{\text{NFT}}, A^{\text{NF}}, A^{\text{FT}}, A^{\text{NFT}}, B^{\text{NF}}, B^{\text{FT}}, B^{\text{NFT}}, B^{\text{A}}\}, C \cup \{A^{\text{FT}}, A^{\text{NFT}}, A^{\text{FT}}, A^{\text{FB}}, A^{\text{FT}}, B^{\text{NFT}}, B^{\text{FTB}}, B^{\text{A}}\}, C \cup \{A^F, A^{\text{NF}}, A^{\text{FT}}, A^{\text{FB}}, A^{\text{NFT}}, A^{\text{FTB}}, A^{\text{A}}\}\}} \wedge_f:\text{NFT} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{NFB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{FB}}, B^{\text{NFB}}\}, C \cup \{A^{\text{NFB}}, B^{\text{NFB}}\}, C \cup \{A^{\text{FB}}, A^{\text{NFB}}\}\}} \wedge_f:\text{NFB} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{NTB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^B, B^{\text{NB}}, B^{\text{FB}}, B^{\text{TB}}, B^{\text{NFB}}, B^{\text{FTB}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^B, A^{\text{NB}}, A^{\text{TB}}, A^{\text{NTB}}, B^B, B^{\text{NB}}, B^{\text{TB}}, B^{\text{NTB}}\}, C \cup \{A^{\text{NB}}, A^{\text{NFB}}, A^{\text{NTB}}, A^{\text{A}}, B^{\text{NB}}, B^{\text{NFB}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^{\text{TB}}, A^{\text{NFTB}}, A^{\text{FTB}}, A^{\text{A}}, B^{\text{TB}}, B^{\text{NTB}}, B^{\text{FTB}}, B^{\text{A}}\}, C \cup \{A^B, A^{\text{NB}}, A^{\text{FB}}, A^{\text{TB}}, A^{\text{NFB}}, A^{\text{FTB}}, A^{\text{A}}\}\}} \wedge_f:\text{NTB} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{FTB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{FB}}, B^{\text{FTB}}\}, C \cup \{A^{\text{FTB}}, B^{\text{FTB}}\}, C \cup \{A^{\text{FB}}, A^{\text{FTB}}\}\}} \wedge_f:\text{FTB} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_f B)^{\text{A}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{FB}}, B^{\text{NFB}}, B^{\text{FTB}}, B^{\text{A}}\}, C \cup \{A^{\text{NFB}}, A^{\text{A}}, B^{\text{NFB}}, B^{\text{A}}\}, C \cup \{A^{\text{FTB}}, A^{\text{A}}, B^{\text{FTB}}, B^{\text{A}}\}, C \cup \{A^{\text{FB}}, A^{\text{NFB}}, A^{\text{FTB}}, A^{\text{A}}\}\}} \wedge_f:\text{A} \\
\end{array}$$

The clause formation rules for connective  $\vee_f$  are given by

$$\begin{array}{c}
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{N}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{N}}, B^{\text{N}}, B^{\text{T}}, B^{\text{NT}}\}, C \cup \{A^{\text{N}}, A^{\text{N}}, B^{\text{N}}, B^{\text{N}}\}, C \cup \{A^{\text{N}}, A^{\text{T}}, B^{\text{N}}, B^{\text{T}}\}, C \cup \{A^{\text{N}}, A^{\text{N}}, A^{\text{T}}, A^{\text{NT}}\}\}} \vee_f:\text{N} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{N}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{N}}, B^{\text{NT}}\}, C \cup \{A^{\text{N}}, B^{\text{N}}\}, C \cup \{A^{\text{N}}, A^{\text{NT}}\}\}} \vee_f:\text{N} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{F}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{N}}, B^{\text{F}}, B^{\text{T}}, B^{\text{NT}}, B^{\text{FB}}, B^{\text{NFT}}, B^{\text{FT}}, B^{\text{NFT}}\}, C \cup \{A^{\text{N}}, A^{\text{N}}, A^{\text{F}}, A^{\text{NF}}, B^{\text{N}}, B^{\text{N}}, B^{\text{F}}, B^{\text{NF}}\}, C \cup \{A^{\text{N}}, A^{\text{F}}, A^{\text{T}}, A^{\text{FT}}, B^{\text{N}}, B^{\text{F}}, B^{\text{FT}}\}, C \cup \{A^{\text{F}}, A^{\text{NF}}, A^{\text{FT}}, A^{\text{NFT}}, B^{\text{F}}, B^{\text{NF}}, B^{\text{FT}}, B^{\text{NFT}}\}, C \cup \{A^{\text{N}}, A^{\text{F}}, A^{\text{T}}, A^{\text{NF}}, A^{\text{NT}}, A^{\text{FT}}, A^{\text{NFT}}\}\}} \vee_f:\text{F} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{T}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{T}}, B^{\text{NT}}\}, C \cup \{A^{\text{T}}, B^{\text{T}}\}, C \cup \{A^{\text{T}}, A^{\text{NT}}\}\}} \vee_f:\text{T} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{NT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{N}}, B^{\text{N}}, B^{\text{B}}, B^{\text{NT}}, B^{\text{NB}}, B^{\text{TB}}, B^{\text{NTB}}\}, C \cup \{A^{\text{N}}, A^{\text{N}}, A^{\text{B}}, A^{\text{NB}}, B^{\text{N}}, B^{\text{N}}, B^{\text{B}}, B^{\text{NB}}\}, C \cup \{A^{\text{N}}, A^{\text{T}}, A^{\text{B}}, A^{\text{TB}}, B^{\text{N}}, B^{\text{T}}, B^{\text{B}}, B^{\text{TB}}\}, C \cup \{A^{\text{B}}, A^{\text{NB}}, A^{\text{TB}}, A^{\text{NTB}}, B^{\text{B}}, B^{\text{NB}}, B^{\text{TB}}, B^{\text{NTB}}\}, C \cup \{A^{\text{N}}, A^{\text{N}}, A^{\text{T}}, A^{\text{B}}, A^{\text{NB}}, A^{\text{TB}}, A^{\text{NTB}}\}\}} \vee_f:\text{NT} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{NF}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{N}}, B^{\text{NF}}, B^{\text{NT}}, B^{\text{NFT}}\}, C \cup \{A^{\text{N}}, A^{\text{NF}}, B^{\text{N}}, B^{\text{NF}}\}, C \cup \{A^{\text{NF}}, A^{\text{NFT}}, B^{\text{NF}}, B^{\text{NFT}}\}, C \cup \{A^{\text{N}}, A^{\text{NF}}, A^{\text{NT}}, A^{\text{NFT}}\}\}} \vee_f:\text{NF} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{NT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{NT}}\}, C \cup \{A^{\text{NT}}\}\}} \vee_f:\text{NT} \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{FT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{T}}, B^{\text{NT}}, B^{\text{FT}}, B^{\text{NFT}}\}, C \cup \{A^{\text{T}}, A^{\text{FT}}, B^{\text{T}}, B^{\text{FT}}\}, C \cup \{A^{\text{FT}}, A^{\text{NFT}}, B^{\text{FT}}, B^{\text{NFT}}\}, C \cup \{A^{\text{T}}, A^{\text{NT}}, A^{\text{FT}}, A^{\text{NFT}}\}\}} \vee_f:\text{FT}$$

$$\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{NB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^N, B^{\text{NT}}, B^{\text{NB}}, B^{\text{NTB}}\}, C \cup \{A^N, A^{\text{NB}}, B^N, B^{\text{NB}}\}, C \cup \{A^{\text{NB}}, A^{\text{NTB}}, B^{\text{NB}}, B^{\text{NTB}}\}, C \cup \{A^N, A^{\text{NT}}, A^{\text{NB}}, A^{\text{NTB}}\}\}} \vee_f : \text{NB}$$

$$\mathcal{C} \cup \{C \cup \{A^N, A^{\text{NT}}, A^{\text{NB}}, A^{\text{NTB}}\}, C \cup \{A^N, A^{\text{NB}}, A^{\text{NT}}, A^{\text{NTB}}\}, C \cup \{A^{\text{NB}}, A^{\text{NT}}, A^{\text{NTB}}, B^{\text{NTB}}\}, C \cup \{A^{\text{NT}}, A^{\text{NB}}, A^{\text{NTB}}\}\} \cup \{C \cup \{(A \vee_f B)^{\text{TB}}\}\}$$

$$\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{TB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^T, B^{\text{NT}}, B^{\text{TB}}, B^{\text{NTB}}\}, C \cup \{A^T, A^{\text{TB}}, B^T, B^{\text{TB}}\}, C \cup \{A^{\text{TB}}, A^{\text{NTB}}, B^{\text{TB}}, B^{\text{NTB}}\}, C \cup \{A^T, A^{\text{NT}}, A^{\text{TB}}, A^{\text{NTB}}\}\}} \vee_f : \text{TB}$$

$$\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{NFT}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{NT}}, B^{\text{NFT}}\}, C \cup \{A^{\text{NFT}}, B^{\text{NFT}}\}, C \cup \{A^{\text{NT}}, A^{\text{NFT}}\}\}} \vee_f : \text{NFT}$$

$$\mathcal{C} \cup \{C \cup \{B^N, B^{\text{NF}}, B^{\text{NT}}, B^{\text{NB}}, B^{\text{NFT}}, B^{\text{NFB}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^N, A^{\text{NF}}, A^{\text{NB}}, A^{\text{NFB}}, B^N, B^{\text{NF}}, B^{\text{NB}}, B^{\text{NFT}}\}, C \cup \{A^{\text{NF}}, A^{\text{NFT}}, A^{\text{NFB}}, A^{\text{NTB}}, B^{\text{NFT}}, B^{\text{NFB}}, B^{\text{A}}\}, C \cup \{A^{\text{NFB}}, A^{\text{NTB}}, A^{\text{A}}, B^{\text{NB}}, B^{\text{NFB}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^N, A^{\text{NF}}, A^{\text{NT}}, A^{\text{NB}}, A^{\text{NFB}}, A^{\text{NTB}}, A^{\text{A}}\}\} \cup \{C \cup \{(A \vee_f B)^{\text{NFB}}\}\}$$

$$\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{NTB}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{NT}}, B^{\text{NTB}}\}, C \cup \{A^{\text{NTB}}, B^{\text{NTB}}\}, C \cup \{A^{\text{NT}}, A^{\text{NTB}}\}\}} \vee_f : \text{NTB}$$

$$\mathcal{C} \cup \{C \cup \{B^T, B^{\text{NT}}, B^{\text{TB}}, B^{\text{NFT}}, B^{\text{NFB}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^T, A^{\text{NT}}, A^{\text{TB}}, A^{\text{NFT}}, B^T, B^{\text{NT}}, B^{\text{TB}}, B^{\text{NFT}}\}, C \cup \{A^{\text{NT}}, A^{\text{NFT}}, A^{\text{TB}}, B^{\text{NT}}, B^{\text{NFT}}, B^{\text{A}}\}, C \cup \{A^{\text{TB}}, A^{\text{NT}}, A^{\text{NFT}}, A^{\text{TB}}, B^{\text{NT}}, B^{\text{NFT}}, B^{\text{A}}\}, C \cup \{A^T, A^{\text{NT}}, A^{\text{TB}}, A^{\text{NFT}}, A^{\text{TB}}, A^{\text{NT}}, A^{\text{A}}\}\} \cup \{C \cup \{(A \vee_f B)^{\text{FTB}}\}\}$$

$$\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{\text{A}}\}\}}{\mathcal{C} \cup \{C \cup \{B^{\text{NT}}, B^{\text{NFT}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^{\text{NFT}}, A^{\text{A}}, B^{\text{NFT}}, B^{\text{A}}\}, C \cup \{A^{\text{NTB}}, A^{\text{A}}, B^{\text{NTB}}, B^{\text{A}}\}, C \cup \{A^{\text{NT}}, A^{\text{NFT}}, A^{\text{NTB}}, A^{\text{A}}\}\}} \vee_f : \text{A}$$

**Theorem 28.** Let  $T(\mathcal{C})$  be the result of exhaustively applying the translation rules to a clause set  $\mathcal{C}$ . Then  $T(\mathcal{C})$  is a set of proper clauses, i.e.,  $T(\mathcal{C})$  contains only signed atoms (all connectives are eliminated). Furthermore,  $T(\mathcal{C})$  is satisfiable iff  $\mathcal{C}$  is.

**Proposition 29.** Let  $A$  be a sentence and  $\Delta$  be a set of sentences. Then

1.  $\models A$  iff  $\{\{A^v \mid v \in V^-\}\}$  is unsatisfiable.
2.  $\Delta \models A$  iff

$$\{\{B^w \mid w \in V^+\} \mid B \in \Delta\} \cup \{\{A^v \mid v \in V^-\}\}$$

is unsatisfiable.

**Definition 30.** A clause  $R$  is a *resolvent* of clauses  $C_1, C_2$  if  $R = (C_1 \setminus \{A^{v_1}\}) \cup (C_2 \setminus \{A^{v_2}\})$  where  $v_1 \neq v_2$

**Definition 31.** A *resolution refutation* of a clause set  $\mathcal{C}$  is a sequence of clauses  $C_1, \dots, C_n$  so that for every  $i$ ,  $C_i \in \mathcal{C}$  or  $C_i$  is a resolvent of  $C_j, C_k$  with  $j, k < i$ , and  $C_n = \emptyset$ .

**Theorem 32.** A clause set  $\mathcal{C}$  is unsatisfiable iff it has a resolution refutation.

*Proof.* See Theorems 3.14 and 3.19 of Baaz and Fermüller [1] or Theorems 2.5.5 and 2.5.8 of Zach [15].  $\square$

**Corollary 33.**  $\Delta \models A$  iff

$$T(\{\{B^w \mid w \in V^+\} \mid B \in \Delta\} \cup \{\{A^v \mid v \in V^-\}\})$$

has a resolution refutation.

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## A shramko-wansing.lgc – specification of Shramko-Wansing logic

```

logic "Shramko-Wansing".
truth_values { e , n , f , t , b , nf , nt , ft , nb , fb , tb ,
               nft , nfb , ntb , ftb , a }.
designated_truth_values { t , nt , tb , ntb }.
operator(negt/1, mapping {
    (e ) : e ,
    (n ) : t ,
    (f ) : b ,
    (t ) : n ,
    (b ) : f ,
    (nf ) : tb ,
    (nt ) : nt ,
    (nb ) : ft ,
    (ft ) : nb ,
    (fb ) : fb ,
    (tb ) : nf ,
    (nft) : ntb ,
    (nfb) : ftb ,
    (ntb) : nft ,
    (ftb) : nfb ,
    (a ) : a
}).
operator(negf/1, mapping {
    (e ) : e ,
    (n ) : f ,
    (f ) : n ,
    (t ) : b ,
    (b ) : t ,
    (nf ) : nf ,
    (nt ) : fb ,
    (nb ) : ft ,
    (ft ) : nb ,
    (fb ) : nt ,
    (tb ) : tb ,
    (nft) : nfb ,
    (nfb) : nft ,
    (ntb) : ftb ,
}
).
```

```

        (ftb) : ntb,
        (a ) : a
    }
).

ordering(truth, "nf < { f < { fb < {ftb < tb, b < tb}, ft < {ftb, t
    < tb}, e < {b,t}}}, nfb < { fb, a < {ftb, ntb}, nb < {b, ntb
    }}, nft < { ft, a, nt < {t, ntb < tb}}, n < {e, nb, nt}}}").
ordering(falsity, "nt < { n < { nf < {f < fb, nfb < fb}, e < {f, b
    < fb}, nb < {nfb, b}}}, nft < { nf, ft < {f, ftb}, a < {nfb,
    ftb}}, t < { e, ft, tb < {b, ftb < fb}}, ntb < {nb, a, tb}}}").
operator(andt /2, inf(truth)).
operator(ort /2, sup(truth)).
operator(andf /2, inf(falsity)).
operator(orf /2, sup(falsity)).

```

## B shramko-wansing.cfg – L<sup>A</sup>T<sub>E</sub>X translations

```

texName(n, "\\boldsymbol{N}").
texName(f, "\\boldsymbol{F}").
texName(t, "\\boldsymbol{T}").
texName(b, "\\boldsymbol{B}").
texName(e, "\\mathbf{N}").
texName(a, "\\mathbf{A}").
texName(nf, "\\mathbf{NF}").
texName(nt, "\\mathbf{NT}").
texName(nb, "\\mathbf{NB}").
texName(ft, "\\mathbf{FT}").
texName(fb, "\\mathbf{FB}").
texName(tb, "\\mathbf{TB}").
texName(nft, "\\mathbf{NFT}").
texName(nfb, "\\mathbf{NFB}").
texName(ntb, "\\mathbf{NTB}").
texName(ftb, "\\mathbf{FTB}").

texName(andt, "\\land_t").
texName(ort, "\\lor_t").
texName(negt, "\\lnot_t").
texInfix(andt).
texInfix(ort).
texPrefix(negt).
texName(andf, "\\land_f").
texName(orf, "\\lor_f").
texName(negf, "\\lnot_f").
texInfix(andf).
texInfix(orf).
texPrefix(negf).

texName(forall, "\\forall_t").
texName(exists, "\\exists_t").

texExtra("ShortName", "\\mathbf{SIXTEEN}_{\{2,\\mathit{des}\}}").
texExtra("Semantics", "The connectives $\\land_t$ and $\\lor_t$ of
$\\ShortName$ are defined as the inf and sup of the $\\le_t$ ordering given in \\cref{fig}. The connectives $\\land_f$ and $\\lor_f$ correspond to inf and sup of $\\le_f$. For simplicity, we leave out the $\\le_i$ ordering and other operators and

```

```
inversions defined in \\cite{ml}.\\begin{figure}\\centering\\
includegraphics[width=.5\\textwidth]{shramko-wansing.png}\\\\
caption{The $\\le_t$ and $\\le_f$ orderings}\\label{fig}\\end{\\
figure}).\\
texExtra("Preamble","\\usepackage{graphicx,amsbsy}").\\
texExtra("Link","https://logic.at/multlog/shramko-wansing.pdf").
```