

# Analytic Proof Systems for Shramko-Wansing logic **SIXTEEN**<sub>2,des</sub>

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June 17, 2021

## Abstract

We give sequent calculus, analytic tableaux, natural deduction, and clause translation systems for resolution for Shramko-Wansing logic **SIXTEEN**<sub>2,des</sub>.

## 1 Introduction

In this paper we present calculi for Shramko-Wansing logic **SIXTEEN**<sub>2,des</sub> [12]. **SIXTEEN**<sub>2,des</sub> has sixteen truth values **N**, **N**, **F**, **T**, **B**, **NF**, **NT**, **FT**, **NB**, **FB**, **TB**, **NFT**, **NFB**, **NTB**, **FTB**, **A** (with **T**, **NT**, **TB**, **NTB** designated), and connectives  $\neg_t$ ,  $\neg_f$ ,  $\wedge_t$ ,  $\vee_t$ ,  $\wedge_f$ ,  $\vee_f$ . Its syntax and semantics is detailed in section 2.

We first present a 16-sided sequent calculus in section 3. The fundamental idea for many-sided sequent calculi for finite-valued logics goes back to Schröter [11], Rousseau [8], Takahashi [14]. We follow the method given by Baaz, Fermüller, and Zach [4] and Zach [15] for constructing inference rules. This guarantees that our system automatically has soundness and completeness theorems, cut-elimination theorem and Maehara lemma (interpolation). For proofs of these results see [4, 15].

Signed tableau systems for finite-valued logics were proposed by Surma [13] and Carnielli [6], and generalized by Hähnle [7]. In section 4, we present a signed tableau system for Shramko-Wansing logic.

Many-valued natural deduction systems for finite-valued logics have been investigated by Baaz, Fermüller, and Zach [3] and Zach [15]. We give the introduction and elimination rules for the natural deduction system for **SIXTEEN**<sub>2,des</sub> in section 5.

In addition to Hähnle’s work on tableaux-based theorem proving for finite-valued logic, Baaz and Fermüller [1] have studied resolution calculi for clauses

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\*Cite as: Multlog (2021), “Analytic proof systems for Shramko-Wansing logic **SIXTEEN**<sub>2,des</sub>.” [PDF generated by MULTLOG, v. 1.16-pre3, <https://logic.at/multlog>]. Available at <https://logic.at/multlog/shramko-wansing.pdf>.

See appendices A and B for the specification of Shramko-Wansing logic.

of signed literals. In order for these calculi to be used to prove that formulas of Shramko-Wansing logic are valid or follow from some others, it is necessary to produce sets of signed clauses. In section 6, we present a translation calculus that yields a set of clauses from a set of formulas.

The rules we provide are optimal in each case, and use the algorithms developed by Salzer [9, 10].

## 2 Syntax and semantics

**Definition 1.** The *propositional language*  $\mathcal{L}$  for Shramko-Wansing logic consists of

1. propositional variables:  $x_0, x_1, x_2, \dots$
2. propositional connectives, arity given in parenthesis:  $\neg_t$  (1),  $\neg_f$  (1),  $\wedge_t$  (2),  $\vee_t$  (2),  $\wedge_f$  (2), and  $\vee_f$  (2)
3. auxiliary symbols: “(”, “)” and “,”

*Formulas* are defined inductively:

1. Every propositional variable is a formula.
2. If  $A$  is a formula, so is  $\neg_t A$ .
3. If  $A$  is a formula, so is  $\neg_f A$ .
4. If  $A$  and  $B$  are formulas, so is  $(A \wedge_t B)$ .
5. If  $A$  and  $B$  are formulas, so is  $(A \vee_t B)$ .
6. If  $A$  and  $B$  are formulas, so is  $(A \wedge_f B)$ .
7. If  $A$  and  $B$  are formulas, so is  $(A \vee_f B)$ .

As a notational convention, lowercase letters will be used to denote variables, possibly indexed. Uppercase letters  $A, B, C, \dots$  will stand for formulas, greek letters  $\Gamma, \Delta, \Lambda, \dots$  for sequences and sets of formulas. The symbol  $\square$  stands for general propositional connectives.

The connectives  $\wedge_t$  and  $\vee_t$  of **SIXTEEN**<sub>2,des</sub> are defined as the inf and sup of the  $\leq_t$  ordering given in fig. 1. The connectives  $\wedge_f$  and  $\vee_f$  correspond to inf and sup of  $\leq_f$ . For simplicity, we leave out the  $\leq_i$  ordering and other operators and inversions defined in [12].

**Definition 2.** The *matrix* for Shramko-Wansing logic is given by:

1. the set of *truth values*  $V = \{\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}\}$ ,
2. the set  $V^+ = \{\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}\} \subseteq V$  of *designated truth values*,

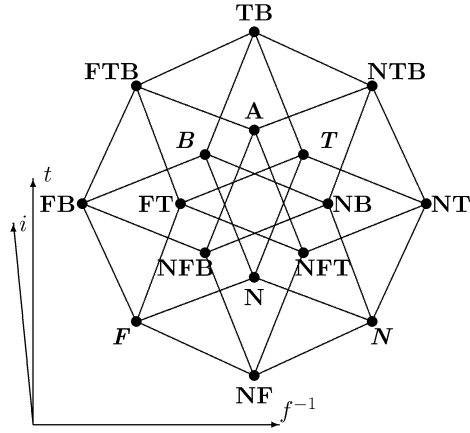


Figure 1: The  $\leq_t$  and  $\leq_f$  orderings

3. the truth functions for connectives  $\neg_t$ ,  $\neg_f$ ,  $\wedge_t$ ,  $\vee_t$ ,  $\wedge_f$  and  $\vee_f$ , as given below;

The set of *undesignated values* is  $V^- = V \setminus V^+ = \{N, N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A\}$ .

The truth functions for connectives  $\neg_t$ ,  $\neg_f$ ,  $\wedge_t$ ,  $\vee_t$ ,  $\wedge_f$  and  $\vee_f$  are defined by

$\neg_t$		$\neg_f$	
N	N	N	N
N	T	N	F
F	B	F	N
T	N	T	B
B	F	B	T
NF	TB	NF	NF
NT	NT	NT	FB
FT	NB	FT	NB
NB	FT	NB	FT
FB	FB	FB	NT
TB	NF	TB	TB
NFT	NTB	NFT	NFT
NFB	FTB	NFB	NFT
NTB	NFT	NTB	FTB
FTB	NTB	FTB	NTB
A	A	A	A

$\widetilde{\wedge}_t$	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
N	N	N	F	N	N	NF	N	F	N	F	N	NF	NF	N	F	NF
N	N	N	NF	N	N	NF	N	NF	N	NF	N	NF	NF	N	NF	NF
F	F	NF	F	F	F	NF	NF	F	NF	F	F	NF	NF	NF	F	NF
T	N	N	F	T	N	NF	NT	FT	N	F	T	NFT	NF	NT	FT	NFT
B	N	N	F	N	B	NF	N	F	NB	FB	B	NF	NFB	NB	FB	NFB
NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF	NF
NT	N	N	NF	NT	N	NF	NT	NFT	N	NF	NT	NFT	NF	NT	NFT	NFT
FT	F	NF	F	FT	F	NF	NFT	FT	NF	F	FT	NFT	NF	NFT	FT	NFT
NB	N	N	NF	N	NB	NF	N	NF	NB	NFB	NB	NF	NFB	NB	NFB	NFB
FB	F	NF	F	F	FB	NF	NF	F	NFB	FB	FB	NF	NFB	NFB	FB	NFB
TB	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
NFT	NF	NF	NF	NFT	NF	NF	NFT	NFT	NF	NF	NFT	NFT	NF	NFT	NFT	NFT
NFB	NF	NF	NF	NF	NFB	NF	NF	NF	NFB	NFB	NFB	NF	NFB	NFB	NFB	NFB
NTB	N	N	NF	NT	NB	NF	NT	NFT	NB	NFB	NTB	NFT	NFB	NTB	A	A
FTB	F	NF	F	FT	FB	NF	NFT	FT	NFB	FB	FTB	NFT	NFB	A	FTB	A
A	NF	NF	NF	NFT	NFB	NF	NFT	NFT	NFB	NFB	A	NFT	NFB	A	A	A
$\widetilde{\vee}_t$	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
N	N	N	N	T	B	N	T	T	B	B	TB	T	B	TB	TB	TB
N	N	N	N	T	B	N	NT	T	NB	B	TB	NT	NB	NTB	TB	NTB
F	N	N	F	T	B	F	T	FT	B	FB	TB	FT	FB	TB	FTB	FTB
T	T	T	T	T	TB	T	T	T	TB	TB	TB	T	TB	TB	TB	TB
B	B	B	B	TB	B	B	TB	TB	B	B	TB	TB	B	TB	TB	TB
NF	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
NT	T	NT	T	T	TB	NT	NT	T	NTB	TB	TB	NT	NTB	NTB	TB	NTB
FT	T	T	FT	T	TB	FT	T	FT	TB	FTB	TB	FT	FTB	TB	FTB	FTB
NB	B	NB	B	TB	B	NB	NTB	TB	NB	B	TB	NTB	NB	NTB	TB	NTB
FB	B	B	FB	TB	B	FB	TB	FTB	B	FB	TB	FTB	FB	TB	FTB	FTB
TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB	TB
NFT	T	NT	FT	T	TB	NFT	NT	FT	NTB	FTB	TB	NFT	A	NTB	FTB	A
NFB	B	NB	FB	TB	B	NFB	NTB	FTB	NB	FB	TB	A	NFB	NTB	FTB	A
NTB	TB	NTB	TB	TB	TB	NTB	NTB	TB	NTB	TB	TB	NTB	NTB	NTB	TB	NTB
FTB	TB	TB	FTB	TB	TB	FTB	TB	FTB	TB	FTB	TB	FTB	FTB	TB	FTB	FTB
A	TB	NTB	FTB	TB	TB	A	NTB	FTB	NTB	FTB	TB	A	A	NTB	FTB	A

$\widetilde{\wedge}_f$	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
N	N	N	N	T	N	N	NT	T	N	N	T	NT	N	NT	T	NT
N	N	N	N	NT	N	N	NT	NT	N	N	NT	NT	N	NT	NT	NT
F	N	N	F	T	N	NF	NT	FT	N	F	T	NFT	NF	NT	FT	NFT
T	T	NT	T	T	T	NT	NT	T	NT	T	T	NT	NT	NT	T	NT
B	N	N	N	T	B	N	NT	T	NB	B	TB	NT	NB	NTB	TB	NTB
NF	N	N	NF	NT	N	NF	NT	NFT	N	NF	NT	NFT	NF	NT	NFT	NFT
NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT	NT
FT	T	NT	FT	T	T	NFT	NT	FT	NT	FT	T	NFT	NFT	NT	FT	NFT
NB	N	N	N	NT	NB	N	NT	NT	NB	NB	NTB	NT	NB	NTB	NTB	NTB
FB	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
TB	T	NT	T	T	TB	NT	NT	T	NTB	TB	TB	NT	NTB	NTB	TB	NTB
NFT	NT	NT	NFT	NT	NT	NFT	NT	NFT	NT	NFT	NT	NFT	NFT	NT	NFT	NFT
NFB	N	N	NF	NT	NB	NF	NT	NFT	NB	NFB	NTB	NFT	NFB	NTB	A	A
NTB	NT	NT	NT	NT	NTB	NT	NT	NT	NTB	NTB	NTB	NT	NTB	NTB	NTB	NTB
FTB	T	NT	FT	T	TB	NFT	NT	FT	NTB	FTB	TB	NFT	A	NTB	FTB	A
A	NT	NT	NFT	NT	NTB	NFT	NT	NFT	NTB	A	NTB	NFT	A	NTB	A	A

$\widetilde{\vee}_f$	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
N	N	N	F	N	B	F	N	F	B	FB	B	F	FB	B	FB	FB
N	N	N	F	N	B	NF	N	F	NB	FB	B	NF	NFB	NB	FB	NFB
F	F	F	F	F	FB	F	F	F	FB	FB	FB	F	FB	FB	FB	FB
T	N	N	F	T	B	F	T	FT	B	FB	TB	FT	FB	TB	FTB	FTB
B	B	B	FB	B	B	FB	B	FB	B	FB	B	FB	FB	B	FB	FB
NF	F	NF	F	F	FB	NF	NF	F	NFB	FB	FB	NF	NFB	NFB	FB	NFB
NT	N	N	F	T	B	NF	NT	FT	NB	FB	TB	NFT	NFB	NTB	FTB	A
FT	F	F	F	FT	FB	F	FT	FT	FB	FB	FTB	FT	FB	FTB	FTB	FTB
NB	B	NB	FB	B	B	NFB	NB	FB	NB	FB	B	NFB	NFB	NB	FB	NFB
FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB	FB
TB	B	B	FB	TB	B	FB	TB	FTB	B	FB	TB	FTB	FB	TB	FTB	FTB
NFT	F	NF	F	FT	FB	NF	NFT	FT	NFB	FB	FTB	NFT	NFB	A	FTB	A
NFB	FB	NFB	FB	FB	FB	NFB	NFB	FB	NFB	FB	FB	NFB	NFB	NFB	FB	NFB
NTB	B	NB	FB	TB	B	NFB	NTB	FTB	NB	FB	TB	A	NFB	NTB	FTB	A
FTB	FB	FB	FB	FTB	FB	FB	FTB	FTB	FB	FB	FTB	FTB	FB	FTB	FTB	FTB
A	FB	NFB	FB	FTB	FB	NFB	A	FTB	NFB	FB	FTB	A	NFB	A	FTB	A

**Definition 3.** Let  $A$  be a formula and  $x_0, x_1, \dots, x_k$  the variables occurring in  $A$ . Then an *interpretation*  $\mathcal{J}$  of  $A$  is an assignment of truth values to the variables.

**Definition 4.** Given an interpretation  $\mathcal{J}$ , we define the *valuation*  $\text{val}_{\mathcal{J}}$  for formulas  $A$  to truth values as follows:

1. If  $A$  is atomic, then  $\text{val}_{\mathcal{J}}(A)$  simply is the interpretation of  $A$ .
2. If  $A = \square(A_1, \dots, A_n)$ , where  $A_1, \dots, A_n$  are formulas, and  $\widetilde{\square}$  is the associated truth function to  $\square$ , then  $\text{val}_{\mathcal{J}}(A) = \widetilde{\square}(\text{val}_{\mathcal{J}}(A_1), \dots, \text{val}_{\mathcal{J}}(A_n))$ .

**Definition 5.** An interpretation  $\mathcal{J}$  *satisfies* a formula  $A$ , in symbols:  $\mathcal{J} \models A$ , iff  $\text{val}_{\mathcal{J}}(A) \in V^+$ .

**Definition 6.**  $\Delta$  entails  $A$  iff  $\mathfrak{J} \models A$  for every interpretation  $\mathfrak{J}$  such that  $\mathfrak{J} \models B$  for all  $B \in \Delta$ .  $A$  is a *tautology* iff it is satisfied by every interpretation  $\mathfrak{J}$ .

### 3 Sequent calculus for Shramko-Wansing logic

**Definition 7** (Syntax of Sequents). A *sequent*  $\Gamma$  is a 16-tuple

$$\Gamma_{\mathbf{N}} \mid \Gamma_{\mathbf{N}} \mid \Gamma_{\mathbf{F}} \mid \Gamma_{\mathbf{T}} \mid \Gamma_{\mathbf{B}} \mid \Gamma_{\mathbf{NF}} \mid \Gamma_{\mathbf{NT}} \mid \Gamma_{\mathbf{FT}} \mid \Gamma_{\mathbf{NB}} \mid \Gamma_{\mathbf{FB}} \mid \Gamma_{\mathbf{TB}} \mid \Gamma_{\mathbf{NFT}} \mid \Gamma_{\mathbf{NFB}} \mid \Gamma_{\mathbf{NTB}} \mid \Gamma_{\mathbf{FTB}} \mid \Gamma_{\mathbf{A}}$$

of finite sequences  $\Gamma_v$  of formulas, where  $v \in V$ . The  $\Gamma_v$  are called the *components* of  $\Gamma$ .

For a sequence of formulas  $\Delta$ , and  $W \subseteq V$ , let  $[W: \Delta]$  denote the sequent whose component  $\Gamma_v$  is  $\Delta$  if  $v \in W$  and empty otherwise. For  $\{\{w_1, \dots, w_k\}: \Delta\}$  we also write  $[w_1, \dots, w_k: \Delta]$ . If  $\Gamma$  and  $\Gamma'$  are sequents, then  $\Gamma, \Gamma'$  denotes the component-wise union, i.e., the  $v$ -component of  $\Gamma, \Gamma'$  is  $\Gamma_v, \Gamma'_v$ .

**Definition 8.** Let  $\mathfrak{J}$  be an interpretation.  $\mathfrak{J}$  *satisfies* a sequent  $\Gamma$  iff there is a  $v \in V$  so that for some formula  $A \in \Gamma_v$ ,  $\text{val}_{\mathfrak{J}}(A) = v$ .  $\mathfrak{J}$  is called a *model* of  $\Gamma$ , in symbols  $\mathfrak{J} \models \Gamma$ .

$\Gamma$  is called *satisfiable* iff there is an interpretation  $\mathfrak{J}$  so that  $\mathfrak{J} \models \Gamma$  and *valid* iff for every interpretation  $\mathfrak{J}$ ,  $\mathfrak{J} \models \Gamma$ .

**Proposition 9.**  $\Delta \models A$  iff the sequent  $[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \Delta], [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A]$  is valid.

**Definition 10.** The *sequent calculus* for Shramko-Wansing logic is given by:

1. axiom schemas of the form  $[V: A]$ ,
2. weakening rules for every truth value  $v$ :

$$\frac{\Gamma}{\Gamma, [v: A]} \text{w:v}$$

3. exchange rules for every truth value  $v$ :

$$\frac{\Gamma, [v: A, B], \Delta}{\Gamma, [v: B, A], \Delta} \text{x:v}$$

4. contraction rules for every truth value  $v$ :

$$\frac{\Gamma, [v: A, A]}{\Gamma, [v: A]} \text{c:v}$$

5. cut rules for every two truth values  $v \neq w$ :

$$\frac{\Gamma, [v: A] \quad \Delta, [w: A]}{\Gamma, \Delta} \text{cut:vw}$$

6. an introduction rule  $\Box: v$  for every connective  $\Box$  and every truth value  $v$ , as specified below.

(2)–(5) are called *structural rules*. (6) are called *logical rules*.

The introduction rules for connective  $\neg_t$  are given by

$$\begin{array}{c}
\frac{\Gamma, [\mathbf{N}: A]}{\Gamma, [\mathbf{N}: \neg_t A]} \neg_t: \mathbf{N} \quad \frac{\Gamma, [\mathbf{T}: A]}{\Gamma, [\mathbf{N}: \neg_t A]} \neg_t: \mathbf{N} \quad \frac{\Gamma, [\mathbf{B}: A]}{\Gamma, [\mathbf{F}: \neg_t A]} \neg_t: \mathbf{F} \\
\frac{\Gamma, [\mathbf{N}: A]}{\Gamma, [\mathbf{T}: \neg_t A]} \neg_t: \mathbf{T} \quad \frac{\Gamma, [\mathbf{F}: A]}{\Gamma, [\mathbf{B}: \neg_t A]} \neg_t: \mathbf{B} \quad \frac{\Gamma, [\mathbf{TB}: A]}{\Gamma, [\mathbf{NF}: \neg_t A]} \neg_t: \mathbf{NF} \\
\frac{\Gamma, [\mathbf{NT}: A]}{\Gamma, [\mathbf{NT}: \neg_t A]} \neg_t: \mathbf{NT} \quad \frac{\Gamma, [\mathbf{NB}: A]}{\Gamma, [\mathbf{FT}: \neg_t A]} \neg_t: \mathbf{FT} \quad \frac{\Gamma, [\mathbf{FT}: A]}{\Gamma, [\mathbf{NB}: \neg_t A]} \neg_t: \mathbf{NB} \\
\frac{\Gamma, [\mathbf{FB}: A]}{\Gamma, [\mathbf{FB}: \neg_t A]} \neg_t: \mathbf{FB} \quad \frac{\Gamma, [\mathbf{NF}: A]}{\Gamma, [\mathbf{TB}: \neg_t A]} \neg_t: \mathbf{TB} \quad \frac{\Gamma, [\mathbf{NTB}: A]}{\Gamma, [\mathbf{NFT}: \neg_t A]} \neg_t: \mathbf{NFT} \\
\frac{\Gamma}{\Gamma, [\mathbf{NFB}: \neg_t A]} \neg_t: \mathbf{NFB} \quad \frac{\Gamma, [\mathbf{NFT}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{NTB}: \neg_t A]} \neg_t: \mathbf{NTB} \\
\frac{\Gamma, [\mathbf{NFB}: A]}{\Gamma, [\mathbf{FTB}: \neg_t A]} \neg_t: \mathbf{FTB} \quad \frac{\Gamma, [\mathbf{A}: A]}{\Gamma, [\mathbf{A}: \neg_t A]} \neg_t: \mathbf{A}
\end{array}$$

The introduction rules for connective  $\neg_f$  are given by

$$\begin{array}{c}
\frac{\Gamma, [\mathbf{N}: A]}{\Gamma, [\mathbf{N}: \neg_f A]} \neg_f: \mathbf{N} \quad \frac{\Gamma, [\mathbf{F}: A]}{\Gamma, [\mathbf{N}: \neg_f A]} \neg_f: \mathbf{N} \quad \frac{\Gamma, [\mathbf{N}: A]}{\Gamma, [\mathbf{F}: \neg_f A]} \neg_f: \mathbf{F} \\
\frac{\Gamma, [\mathbf{B}: A]}{\Gamma, [\mathbf{T}: \neg_f A]} \neg_f: \mathbf{T} \quad \frac{\Gamma, [\mathbf{T}: A]}{\Gamma, [\mathbf{B}: \neg_f A]} \neg_f: \mathbf{B} \quad \frac{\Gamma, [\mathbf{NF}: A]}{\Gamma, [\mathbf{NF}: \neg_f A]} \neg_f: \mathbf{NF} \\
\frac{\Gamma, [\mathbf{FB}: A]}{\Gamma, [\mathbf{NT}: \neg_f A]} \neg_f: \mathbf{NT} \quad \frac{\Gamma, [\mathbf{NB}: A]}{\Gamma, [\mathbf{FT}: \neg_f A]} \neg_f: \mathbf{FT} \quad \frac{\Gamma, [\mathbf{FT}: A]}{\Gamma, [\mathbf{NB}: \neg_f A]} \neg_f: \mathbf{NB} \\
\frac{\Gamma, [\mathbf{NT}: A]}{\Gamma, [\mathbf{FB}: \neg_f A]} \neg_f: \mathbf{FB} \quad \frac{\Gamma, [\mathbf{TB}: A]}{\Gamma, [\mathbf{TB}: \neg_f A]} \neg_f: \mathbf{TB} \\
\frac{\Gamma, [\mathbf{NFT}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NFT}: \neg_f A]} \neg_f: \mathbf{NFT} \quad \frac{\Gamma}{\Gamma, [\mathbf{NFB}: \neg_f A]} \neg_f: \mathbf{NFB} \\
\frac{\Gamma, [\mathbf{FTB}: A]}{\Gamma, [\mathbf{NTB}: \neg_f A]} \neg_f: \mathbf{NTB} \quad \frac{\Gamma, [\mathbf{NTB}: A]}{\Gamma, [\mathbf{FTB}: \neg_f A]} \neg_f: \mathbf{FTB} \quad \frac{\Gamma, [\mathbf{A}: A]}{\Gamma, [\mathbf{A}: \neg_f A]} \neg_f: \mathbf{A}
\end{array}$$

The introduction rules for connective  $\wedge_t$  are given by

$$\frac{\Gamma, [\mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{TB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{T}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{B}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{TB}: A]}{\Gamma, [\mathbf{N}: A \wedge_t B]} \wedge_t: \mathbf{N}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{NT}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{NT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NT}, \mathbf{NB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{NT}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: A]}{\Gamma, [\mathbf{N}: A \wedge_t B]} \wedge_t: \mathbf{N}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FB}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{F}: A \wedge_t B]} \wedge_t: \mathbf{F}$$

$$\frac{\Gamma, [\mathbf{T}, \mathbf{TB}: B] \quad \Gamma, [\mathbf{T}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{TB}: A]}{\Gamma, [\mathbf{T}: A \wedge_t B]} \wedge_t: \mathbf{T}$$

$$\frac{\Gamma, [\mathbf{B}, \mathbf{TB}: B] \quad \Gamma, [\mathbf{B}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{TB}: A]}{\Gamma, [\mathbf{B}: A \wedge_t B]} \wedge_t: \mathbf{B}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B]}{\Gamma, [\mathbf{NF}: A \wedge_t B]} \wedge_t: \mathbf{NF}$$

$$\frac{\Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}: A, B] \quad \Gamma, [\mathbf{NT}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A]}{\Gamma, [\mathbf{NT}: A \wedge_t B]} \wedge_t: \mathbf{NT}$$

$$\frac{\Gamma, [\mathbf{T}, \mathbf{FT}, \mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{T}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{FT}, \mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{FT}: A \wedge_t B]} \wedge_t: \mathbf{FT}$$

$$\frac{\Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: A]}{\Gamma, [\mathbf{NB}: A \wedge_t B]} \wedge_t: \mathbf{NB}$$

$$\frac{\Gamma, [\mathbf{B}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{B}, \mathbf{FB}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{FB}: A \wedge_t B]} \wedge_t: \mathbf{FB}$$

$$\frac{\Gamma, [\mathbf{TB}: B] \quad \Gamma, [\mathbf{TB}: A]}{\Gamma, [\mathbf{TB}: A \wedge_t B]} \wedge_t: \mathbf{TB}$$

$$\frac{\Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{TB}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{NT}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{NFT}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{TB}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{NFT}: A \wedge_t B]} \wedge_t: \mathbf{NFT}$$

$$\frac{\Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{NFB}: A \wedge_t B]} \wedge_t: \mathbf{NFB}$$

$$\frac{\Gamma, [\mathbf{TB}, \mathbf{NTB}: B] \quad \Gamma, [\mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{TB}, \mathbf{NTB}: A]}{\Gamma, [\mathbf{NTB}: A \wedge_t B]} \wedge_t: \mathbf{NTB}$$

$$\frac{\Gamma, [\mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{FTB}: A \wedge_t B]} \wedge_t: \mathbf{FTB}$$

$$\frac{\Gamma, [\mathbf{TB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{TB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{A}: A \wedge_t B]} \wedge_t: \mathbf{A}$$

The introduction rules for connective  $\vee_t$  are given by

$$\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}: A]}{\Gamma, [\mathbf{N}: A \vee_t B]} \vee_t: \mathbf{N}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{NF}: B] \quad \Gamma, [\mathbf{N}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}: A]}{\Gamma, [\mathbf{N}: A \vee_t B]} \vee_t: \mathbf{N}$$



$$\frac{\Gamma, [\mathbf{F}, \mathbf{NF}: B] \quad \Gamma, [\mathbf{F}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}: A]}{\Gamma, [\mathbf{F}: A \vee_t B]} \vee_t: \mathbf{F}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{NT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A]}{\Gamma, [\mathbf{T}: A \vee_t B]} \vee_t: \mathbf{T}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FB}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{B}: A \vee_t B]} \vee_t: \mathbf{B}$$

$$\frac{\Gamma, [\mathbf{NF}: B] \quad \Gamma, [\mathbf{NF}: A]}{\Gamma, [\mathbf{NF}: A \vee_t B]} \vee_t: \mathbf{NF}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NFT}: B] \quad \Gamma, [\mathbf{N}, \mathbf{NT}: A, B] \quad \Gamma, [\mathbf{NT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NFT}: A]}{\Gamma, [\mathbf{NT}: A \vee_t B]} \vee_t: \mathbf{NT}$$

$$\frac{\Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: B] \quad \Gamma, [\mathbf{F}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: A]}{\Gamma, [\mathbf{FT}: A \vee_t B]} \vee_t: \mathbf{FT}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NB}: A \vee_t B]} \vee_t: \mathbf{NB}$$

$$\frac{\Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{F}, \mathbf{FB}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{FB}: A \vee_t B]} \vee_t: \mathbf{FB}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{NT}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{TB}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B]}{\Gamma, [\mathbf{TB}: A \vee_t B]} \vee_t: \mathbf{TB}$$

$$\frac{\Gamma, [\mathbf{NF}, \mathbf{NFT}: B] \quad \Gamma, [\mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{NF}, \mathbf{NFT}: A]}{\Gamma, [\mathbf{NFT}: A \vee_t B]} \vee_t: \mathbf{NFT}$$

$$\frac{\Gamma, [\mathbf{NF}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{NF}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NFB}: A \vee_t B]} \vee_t: \mathbf{NFB}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NT}, \mathbf{NB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{NT}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{NTB}: A \vee_t B]} \vee_t: \mathbf{NTB}$$

$$\frac{\Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{NFT}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{FTB}: A \vee_t B]} \vee_t: \mathbf{FTB}$$

$$\frac{\Gamma, [\mathbf{NF}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: B] \quad \Gamma, [\mathbf{NFT}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{NFB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{NF}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: A]}{\Gamma, [\mathbf{A}: A \vee_t B]} \vee_t: \mathbf{A}$$

The introduction rules for connective  $\wedge_f$  are given by

$$\frac{\Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{F}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{B}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FB}: A]}{\Gamma, [\mathbf{N}: A \wedge_f B]} \wedge_f: \mathbf{N}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{N}: A \wedge_f B]} \wedge_f: \mathbf{N}$$

$$\frac{\Gamma, [\mathbf{F}, \mathbf{FB}: B] \quad \Gamma, [\mathbf{F}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{FB}: A]}{\Gamma, [\mathbf{F}: A \wedge_f B]} \wedge_f: \mathbf{F}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{TB}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{FT}, \mathbf{TB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{T}: A \wedge_f B]} \wedge_f: \mathbf{T}$$

$$\frac{\Gamma, [\mathbf{B}, \mathbf{FB}: B] \quad \Gamma, [\mathbf{B}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{FB}: A]}{\Gamma, [\mathbf{B}: A \wedge_f B]} \wedge_f: \mathbf{B}$$

$$\frac{\Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}: A, B] \quad \Gamma, [\mathbf{NF}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NF}: A \wedge_f B]} \wedge_f: \mathbf{NF}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{B}, \mathbf{NT}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NF}, \mathbf{NT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}, \mathbf{FT}, \mathbf{TB}, \mathbf{NFT}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B]}{\Gamma, [\mathbf{NT}: A \wedge_f B]} \wedge_f: \mathbf{NT}$$

$$\frac{\Gamma, [\mathbf{F}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{F}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{FT}: A \wedge_f B]} \wedge_f: \mathbf{FT}$$

$$\frac{\Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NB}: A \wedge_f B]} \wedge_f: \mathbf{NB}$$

$$\frac{\Gamma, [\mathbf{FB}: B] \quad \Gamma, [\mathbf{FB}: A]}{\Gamma, [\mathbf{FB}: A \wedge_f B]} \wedge_f: \mathbf{FB}$$

$$\frac{\Gamma, [\mathbf{B}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{B}, \mathbf{TB}: A, B] \quad \Gamma, [\mathbf{TB}, \mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{FB}, \mathbf{TB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{TB}: A \wedge_f B]} \wedge_f: \mathbf{TB}$$

$$\frac{\Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{NF}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FT}, \mathbf{NFT}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{NFT}: A \wedge_f B]} \wedge_f: \mathbf{NFT}$$

$$\frac{\Gamma, [\mathbf{FB}, \mathbf{NFB}: B] \quad \Gamma, [\mathbf{NFB}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{NFB}: A]}{\Gamma, [\mathbf{NFB}: A \wedge_f B]} \wedge_f: \mathbf{NFB}$$

$$\frac{\Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{TB}, \mathbf{NTB}: A, B] \quad \Gamma, [\mathbf{NB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{TB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{B}, \mathbf{NB}, \mathbf{FB}, \mathbf{TB}, \mathbf{NFB}, \mathbf{NTB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{NTB}: A \wedge_f B]} \wedge_f: \mathbf{NTB}$$

$$\frac{\Gamma, [\mathbf{FB}, \mathbf{FTB}: B] \quad \Gamma, [\mathbf{FTB}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{FTB}: A]}{\Gamma, [\mathbf{FTB}: A \wedge_f B]} \wedge_f: \mathbf{FTB}$$

$$\frac{\Gamma, [\mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{NFB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FTB}, \mathbf{A}: A, B] \quad \Gamma, [\mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A]}{\Gamma, [\mathbf{A}: A \wedge_f B]} \wedge_f: \mathbf{A}$$

The introduction rules for connective  $\vee_f$  are given by

$$\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{NT}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{T}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{T}, \mathbf{NT}: A]}{\Gamma, [\mathbf{N}: A \vee_f B]} \vee_f: \mathbf{N}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{NT}: B] \quad \Gamma, [\mathbf{N}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{NT}: A]}{\Gamma, [\mathbf{N}: A \vee_f B]} \vee_f: \mathbf{N}$$

$$\frac{\Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{FT}: A, B] \quad \Gamma, [\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: A, B] \quad \Gamma, [\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{T}, \mathbf{NF}, \mathbf{NT}, \mathbf{FT}, \mathbf{NFT}: A]}{\Gamma, [\mathbf{F}: A \vee_f B]} \vee_f: \mathbf{F}$$

$$\frac{\Gamma, [\mathbf{T}, \mathbf{NT}: B] \quad \Gamma, [\mathbf{T}: A, B] \quad \Gamma, [\mathbf{T}, \mathbf{NT}: A]}{\Gamma, [\mathbf{T}: A \vee_f B]} \vee_f: \mathbf{T}$$

$$\begin{array}{c}
\frac{\Gamma, [N, N, T, B, NT, NB, TB, NTB: B] \quad \Gamma, [N, N, B, NB: A, B] \quad \Gamma, [N, T, B, TB: A, B] \quad \Gamma, [B, NB, TB, NTB: A, B] \quad \Gamma, [N, N, T, B, NT, NB, TB, NTB: A]}{\Gamma, [B: A \vee_f B]} \vee_f: B \\
\\
\frac{\Gamma, [N, NF, NT, NFT: B] \quad \Gamma, [N, NF: A, B] \quad \Gamma, [NF, NFT: A, B] \quad \Gamma, [N, NF, NT, NFT: A]}{\Gamma, [NF: A \vee_f B]} \vee_f: NF \\
\\
\frac{\Gamma, [NT: B] \quad \Gamma, [NT: A]}{\Gamma, [NT: A \vee_f B]} \vee_f: NT \\
\\
\frac{\Gamma, [T, NT, FT, NFT: B] \quad \Gamma, [T, FT: A, B] \quad \Gamma, [FT, NFT: A, B] \quad \Gamma, [T, NT, FT, NFT: A]}{\Gamma, [FT: A \vee_f B]} \vee_f: FT \\
\\
\frac{\Gamma, [N, NT, NB, NTB: B] \quad \Gamma, [N, NB: A, B] \quad \Gamma, [NB, NTB: A, B] \quad \Gamma, [N, NT, NB, NTB: A]}{\Gamma, [NB: A \vee_f B]} \vee_f: NB \\
\\
\frac{\Gamma, [N, N, F, B, NF, NB, FB, NFB: A, B] \quad \Gamma, [N, F, T, B, FT, FB, TB, FTB: A, B] \quad \Gamma, [F, NF, FT, FB, NFT, NFB, FTB, A: A, B] \quad \Gamma, [B, NB, FB, TB, NFB, NTB, FTB, A: A, B]}{\Gamma, [FB: A \vee_f B]} \vee_f: FB \\
\\
\frac{\Gamma, [T, NT, TB, NTB: B] \quad \Gamma, [T, TB: A, B] \quad \Gamma, [TB, NTB: A, B] \quad \Gamma, [T, NT, TB, NTB: A]}{\Gamma, [TB: A \vee_f B]} \vee_f: TB \\
\\
\frac{\Gamma, [NT, NFT: B] \quad \Gamma, [NFT: A, B] \quad \Gamma, [NT, NFT: A]}{\Gamma, [NFT: A \vee_f B]} \vee_f: NFT \\
\\
\frac{\Gamma, [N, NF, NT, NB, NFT, NFB, NTB, A: B] \quad \Gamma, [N, NF, NB, NFB: A, B] \quad \Gamma, [NF, NFT, NFB, A: A, B] \quad \Gamma, [NB, NFB, NTB, A: A, B] \quad \Gamma, [N, NF, NT, NB, NFT, NFB, NTB, A: A]}{\Gamma, [NFB: A \vee_f B]} \vee_f: NFB \\
\\
\frac{\Gamma, [NT, NTB: B] \quad \Gamma, [NTB: A, B] \quad \Gamma, [NT, NTB: A]}{\Gamma, [NTB: A \vee_f B]} \vee_f: NTB \\
\\
\frac{\Gamma, [T, NT, FT, TB, NFT, NTB, FTB, A: B] \quad \Gamma, [T, FT, TB, FTB: A, B] \quad \Gamma, [FT, NFT, FTB, A: A, B] \quad \Gamma, [TB, NTB, FTB, A: A, B] \quad \Gamma, [T, NT, FT, TB, NFT, NTB, FTB, A: A]}{\Gamma, [FTB: A \vee_f B]} \vee_f: FTB \\
\\
\frac{\Gamma, [NT, NFT, NTB, A: B] \quad \Gamma, [NFT, A: A, B] \quad \Gamma, [NTB, A: A, B] \quad \Gamma, [NT, NFT, NTB, A: A]}{\Gamma, [A: A \vee_f B]} \vee_f: A
\end{array}$$

**Definition 11.** An upward tree of sequents is called a *proof* in the sequent calculus iff every leaf is an axiom, and all other sequents in it are obtained from the ones standing immediately above it by applying one of the rules. The sequent at the root of  $P$  is called its *end-sequent*. A sequent  $\Gamma$  is called *provable* iff it is the end-sequent of some proof.

**Theorem 12** (Soundness and Completeness). *A sequent is provable if and only if it is valid.*

*Proof.* See Theorems 3.1 and 3.2 of Baaz et al. [4] or Theorems 3.3.8 and 3.3.10 of Zach [15].  $\square$

**Corollary 13.** *In Shramko-Wansing logic,  $\models A$  iff  $[T, NT, TB, NTB: A]$  has a sequent proof, and  $\Delta \models A$  iff  $[N, N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: \Delta], [T, NT, TB, NTB: A]$  has a proof.*

**Theorem 14** (Cut-elimination). *The cut rule is eliminable in the sequent calculus for Shramko-Wansing logic.*

*Proof.* See Theorem 4.1 of Baaz et al. [4] or Theorem 3.5.3 of Zach [15].  $\square$

**Theorem 15** (Maehara lemma). *The Maehara lemma holds for the sequent calculus for Shramko-Wansing logic.*

*Proof.* See Theorem 3.8.1 of Zach [15].  $\square$

## 4 Tableaux for Shramko-Wansing logic

Although the method of Surma [13] and Carnielli [6] for obtaining signed analytic tableaux systems applies to Shramko-Wansing logic, it has a drawback. As Hähnle [7] pointed out, to show that a formula is valid, it is required to provide as many closed tableaux as there are non-designated values. This is usually not desirable; the generalized approach by Hähnle [7] solves this problem. Below we give a tableau system for Shramko-Wansing logic using the sets of signs  $V \setminus \{v\}$ , i.e., the tableau system exactly dual to that of Carnielli (in the sense of [2]).

**Definition 16.** A *signed formula* is an expression of the form  $v: A$  where  $v \in V$  and  $A$  is a formula.

**Definition 17.** A *tableau* for a set of signed formulas  $\Delta$  is a downward rooted tree of signed formulas where each one is either an element of  $\Delta$  or results from a signed formula in the branch above it by a branch expansion rule. A tableau is *closed* if every branch contains, for some formula  $A$ , the signed formulas  $v: A$  for all  $v \in V$ , or a signed formula  $v: A$  with a branch expansion rule that explicitly closes the branch ( $\otimes$ ).

The branch expansion rules for connective  $\neg_t$  are given by

$$\begin{array}{c} \frac{\mathbf{N}: \neg_t A}{\mathbf{N}: A} \quad \frac{\mathbf{N}: \neg_t A}{\mathbf{T}: A} \quad \frac{\mathbf{F}: \neg_t A}{\mathbf{B}: A} \quad \frac{\mathbf{T}: \neg_t A}{\mathbf{N}: A} \quad \frac{\mathbf{B}: \neg_t A}{\mathbf{F}: A} \quad \frac{\mathbf{NF}: \neg_t A}{\mathbf{TB}: A} \\ \\ \frac{\mathbf{NT}: \neg_t A}{\mathbf{NT}: A} \quad \frac{\mathbf{FT}: \neg_t A}{\mathbf{NB}: A} \quad \frac{\mathbf{NB}: \neg_t A}{\mathbf{FT}: A} \quad \frac{\mathbf{FB}: \neg_t A}{\mathbf{FB}: A} \quad \frac{\mathbf{TB}: \neg_t A}{\mathbf{NF}: A} \quad \frac{\mathbf{NFT}: \neg_t A}{\mathbf{NTB}: A} \\ \\ \frac{\mathbf{NTB}: \neg_t A}{\mathbf{NFT}: A} \quad \frac{\mathbf{FTB}: \neg_t A}{\mathbf{NFB}: A} \quad \frac{\mathbf{A}: \neg_t A}{\mathbf{A}: A} \\ \mathbf{FTB}: A \end{array}$$

The branch expansion rules for connective  $\neg_f$  are given by

$$\begin{array}{c} \frac{\mathbf{N}: \neg_f A}{\mathbf{N}: A} \quad \frac{\mathbf{N}: \neg_f A}{\mathbf{F}: A} \quad \frac{\mathbf{F}: \neg_f A}{\mathbf{N}: A} \quad \frac{\mathbf{T}: \neg_f A}{\mathbf{B}: A} \quad \frac{\mathbf{B}: \neg_f A}{\mathbf{T}: A} \quad \frac{\mathbf{NF}: \neg_f A}{\mathbf{NF}: A} \\ \\ \frac{\mathbf{NT}: \neg_f A}{\mathbf{FB}: A} \quad \frac{\mathbf{FT}: \neg_f A}{\mathbf{NB}: A} \quad \frac{\mathbf{NB}: \neg_f A}{\mathbf{FT}: A} \quad \frac{\mathbf{FB}: \neg_f A}{\mathbf{NT}: A} \quad \frac{\mathbf{TB}: \neg_f A}{\mathbf{TB}: A} \end{array}$$

$$\begin{array}{c}
\frac{\mathbf{NFT}: \neg_f A}{\mathbf{NFT}: A} \quad \frac{\mathbf{NTB}: \neg_f A}{\mathbf{FTB}: A} \quad \frac{\mathbf{FTB}: \neg_f A}{\mathbf{NTB}: A} \quad \frac{\mathbf{A}: \neg_f A}{\mathbf{A}: A} \\
\mathbf{NFB}: A
\end{array}$$

The branch expansion rules for connective  $\wedge_t$  are given by

$$\frac{\mathbf{N}: A \wedge_t B}{\begin{array}{cccc} \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A \\ \mathbf{T}: B & \mathbf{T}: A & \mathbf{B}: A & \mathbf{T}: A \\ \mathbf{B}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{B}: A \\ \mathbf{TB}: B & \mathbf{T}: B & \mathbf{B}: B & \mathbf{TB}: A \end{array}}$$

$$\frac{\mathbf{N}: A \wedge_t B}{\begin{array}{ccccc} \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A \\ \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{NT}: A & \mathbf{N}: A \\ \mathbf{T}: B & \mathbf{T}: A & \mathbf{B}: A & \mathbf{NB}: A & \mathbf{T}: A \\ \mathbf{B}: B & \mathbf{NT}: A & \mathbf{NB}: A & \mathbf{NTB}: A & \mathbf{B}: A \\ \mathbf{NT}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{NT}: A \\ \mathbf{NB}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{NT}: B & \mathbf{NB}: A \\ \mathbf{TB}: B & \mathbf{T}: B & \mathbf{B}: B & \mathbf{NB}: B & \mathbf{TB}: A \\ \mathbf{NTB}: B & \mathbf{NT}: B & \mathbf{NB}: B & \mathbf{NTB}: B & \mathbf{NTB}: A \end{array}}$$

$$\frac{\mathbf{F}: A \wedge_t B}{\begin{array}{ccccc} \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{F}: A & \mathbf{N}: A \\ \mathbf{F}: B & \mathbf{F}: A & \mathbf{F}: A & \mathbf{FT}: A & \mathbf{F}: A \\ \mathbf{T}: B & \mathbf{T}: A & \mathbf{B}: A & \mathbf{FB}: A & \mathbf{T}: A \\ \mathbf{B}: B & \mathbf{FT}: A & \mathbf{FB}: A & \mathbf{FTB}: A & \mathbf{B}: A \\ \mathbf{FT}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{F}: B & \mathbf{FT}: A \\ \mathbf{FB}: B & \mathbf{F}: B & \mathbf{F}: B & \mathbf{FT}: B & \mathbf{FB}: A \\ \mathbf{TB}: B & \mathbf{T}: B & \mathbf{B}: B & \mathbf{FB}: B & \mathbf{TB}: A \\ \mathbf{FTB}: B & \mathbf{FT}: B & \mathbf{FB}: B & \mathbf{FTB}: B & \mathbf{FTB}: A \end{array}}$$

$$\frac{\mathbf{T}: A \wedge_t B}{\begin{array}{ccc} \mathbf{T}: B & \mathbf{T}: A & \mathbf{T}: A \\ \mathbf{TB}: B & \mathbf{T}: B & \mathbf{TB}: A \end{array}} \quad \frac{\mathbf{B}: A \wedge_t B}{\begin{array}{ccc} \mathbf{B}: B & \mathbf{B}: A & \mathbf{B}: A \\ \mathbf{TB}: B & \mathbf{B}: B & \mathbf{TB}: A \end{array}}$$

**NF:  $A \wedge_t B$**

<b>N: A</b>	<b>N: A</b>	<b>N: A</b>	<b>F: A</b>
<b>N: A</b>	<b>N: A</b>	<b>NF: A</b>	<b>NF: A</b>
<b>F: A</b>	<b>F: A</b>	<b>NT: A</b>	<b>FT: A</b>
<b>T: A</b>	<b>B: A</b>	<b>NB: A</b>	<b>FB: A</b>
<b>NF: A</b>	<b>NF: A</b>	<b>NFT: A</b>	<b>NFT: A</b>
<b>NT: A</b>	<b>NB: A</b>	<b>NFB: A</b>	<b>NFB: A</b>
<b>FT: A</b>	<b>FB: A</b>	<b>NTB: A</b>	<b>FTB: A</b>
<b>NFT: A</b>	<b>NFB: A</b>	<b>A: A</b>	<b>A: A</b>
<b>N: B</b>	<b>N: B</b>	<b>N: B</b>	<b>F: B</b>
<b>N: B</b>	<b>N: B</b>	<b>NF: B</b>	<b>NF: B</b>
<b>F: B</b>	<b>F: B</b>	<b>NT: B</b>	<b>FT: B</b>
<b>T: B</b>	<b>B: B</b>	<b>NB: B</b>	<b>FB: B</b>
<b>NF: B</b>	<b>NF: B</b>	<b>NFT: B</b>	<b>NFT: B</b>
<b>NT: B</b>	<b>NB: B</b>	<b>NFB: B</b>	<b>NFB: B</b>
<b>FT: B</b>	<b>FB: B</b>	<b>NTB: B</b>	<b>FTB: B</b>
<b>NFT: B</b>	<b>NFB: B</b>	<b>A: B</b>	<b>A: B</b>

**NT:  $A \wedge_t B$**

<b>T: B</b>	<b>T: A</b>	<b>NT: A</b>	<b>T: A</b>
<b>NT: B</b>	<b>NT: A</b>	<b>NTB: A</b>	<b>NT: A</b>
<b>TB: B</b>	<b>T: B</b>	<b>NT: B</b>	<b>TB: A</b>
<b>NTB: B</b>	<b>NT: B</b>	<b>NTB: B</b>	<b>NTB: A</b>

**FT:  $A \wedge_t B$**

<b>T: B</b>	<b>T: A</b>	<b>FT: A</b>	<b>T: A</b>
<b>FT: B</b>	<b>FT: A</b>	<b>FTB: A</b>	<b>FT: A</b>
<b>TB: B</b>	<b>T: B</b>	<b>FT: B</b>	<b>TB: A</b>
<b>FTB: B</b>	<b>FT: B</b>	<b>FTB: B</b>	<b>FTB: A</b>

**NB:  $A \wedge_t B$**

<b>B: B</b>	<b>B: A</b>	<b>NB: A</b>	<b>B: A</b>
<b>NB: B</b>	<b>NB: A</b>	<b>NTB: A</b>	<b>NB: A</b>
<b>TB: B</b>	<b>B: B</b>	<b>NB: B</b>	<b>TB: A</b>
<b>NTB: B</b>	<b>NB: B</b>	<b>NTB: B</b>	<b>NTB: A</b>

**FB:  $A \wedge_t B$**

<b>B: B</b>	<b>B: A</b>	<b>FB: A</b>	<b>B: A</b>
<b>FB: B</b>	<b>FB: A</b>	<b>FTB: A</b>	<b>FB: A</b>
<b>TB: B</b>	<b>B: B</b>	<b>FB: B</b>	<b>TB: A</b>
<b>FTB: B</b>	<b>FB: B</b>	<b>FTB: B</b>	<b>FTB: A</b>

**TB:  $A \wedge_t B$**

<b>TB: B</b>	<b>TB: A</b>
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$\mathbf{NFT}: A \wedge_t B$

$T: B$	$T: A$	$\mathbf{NT}: A$	$\mathbf{FT}: A$	$T: A$
$\mathbf{NT}: B$	$\mathbf{NT}: A$	$\mathbf{NFT}: A$	$\mathbf{NFT}: A$	$\mathbf{NT}: A$
$\mathbf{FT}: B$	$\mathbf{FT}: A$	$\mathbf{NTB}: A$	$\mathbf{FTB}: A$	$\mathbf{FT}: A$
$\mathbf{TB}: B$	$\mathbf{NFT}: A$	$A: A$	$A: A$	$\mathbf{TB}: A$
$\mathbf{NFT}: B$	$T: B$	$\mathbf{NT}: B$	$\mathbf{FT}: B$	$\mathbf{NFT}: A$
$\mathbf{NTB}: B$	$\mathbf{NT}: B$	$\mathbf{NFT}: B$	$\mathbf{NFT}: B$	$\mathbf{NTB}: A$
$\mathbf{FTB}: B$	$\mathbf{FT}: B$	$\mathbf{NTB}: B$	$\mathbf{FTB}: B$	$\mathbf{FTB}: A$
$A: B$	$\mathbf{NFT}: B$	$A: B$	$A: B$	$A: A$

$\mathbf{NFB}: A \wedge_t B$

$B: B$	$B: A$	$\mathbf{NB}: A$	$\mathbf{FB}: A$	$B: A$
$\mathbf{NB}: B$	$\mathbf{NB}: A$	$\mathbf{NFB}: A$	$\mathbf{NFB}: A$	$\mathbf{NB}: A$
$\mathbf{FB}: B$	$\mathbf{FB}: A$	$\mathbf{NTB}: A$	$\mathbf{FTB}: A$	$\mathbf{FB}: A$
$\mathbf{TB}: B$	$\mathbf{NFB}: A$	$A: A$	$A: A$	$\mathbf{TB}: A$
$\mathbf{NFB}: B$	$B: B$	$\mathbf{NB}: B$	$\mathbf{FB}: B$	$\mathbf{NFB}: A$
$\mathbf{NTB}: B$	$\mathbf{NB}: B$	$\mathbf{NFB}: B$	$\mathbf{NFB}: B$	$\mathbf{NTB}: A$
$\mathbf{FTB}: B$	$\mathbf{FB}: B$	$\mathbf{NTB}: B$	$\mathbf{FTB}: B$	$\mathbf{FTB}: A$
$A: B$	$\mathbf{NFB}: B$	$A: B$	$A: B$	$A: A$

$\mathbf{NTB}: A \wedge_t B$

$\mathbf{TB}: B$	$\mathbf{NTB}: A$	$\mathbf{TB}: A$	$\mathbf{TB}: B$	$\mathbf{FTB}: A$	$\mathbf{TB}: A$
$\mathbf{NTB}: B$	$\mathbf{NTB}: B$	$\mathbf{NTB}: A$	$\mathbf{FTB}: B$	$\mathbf{FTB}: B$	$\mathbf{FTB}: A$

$\mathbf{FTB}: A \wedge_t B$

$A: A \wedge_t B$

$\mathbf{TB}: B$	$\mathbf{NTB}: A$	$\mathbf{FTB}: A$	$\mathbf{TB}: A$
$\mathbf{NTB}: B$	$A: A$	$A: A$	$\mathbf{NTB}: A$
$\mathbf{FTB}: B$	$\mathbf{NTB}: B$	$\mathbf{FTB}: B$	$\mathbf{FTB}: A$
$A: B$	$A: B$	$A: B$	$A: A$

The branch expansion rules for connective  $\vee_t$  are given by

$\mathbf{N}: A \vee_t B$

$\mathbf{N}: B$	$\mathbf{N}: A$	$\mathbf{N}: A$	$\mathbf{N}: A$	$\mathbf{N}: B$	$\mathbf{N}: A$	$\mathbf{N}: A$
$\mathbf{N}: B$	$\mathbf{N}: A$	$\mathbf{F}: A$	$\mathbf{N}: A$	$\mathbf{NF}: B$	$\mathbf{N}: B$	$\mathbf{NF}: A$
$\mathbf{F}: B$	$\mathbf{N}: B$	$\mathbf{N}: B$	$\mathbf{F}: A$			
$\mathbf{NF}: B$	$\mathbf{N}: B$	$\mathbf{F}: B$	$\mathbf{NF}: A$			

$\mathbf{N}: A \vee_t B$

$\mathbf{F}: A \vee_t B$

$\mathbf{F}: B$	$\mathbf{F}: A$	$\mathbf{F}: A$
$\mathbf{NF}: B$	$\mathbf{F}: B$	$\mathbf{NF}: A$

$$T: A \vee_t B$$

<b>N</b> : <i>B</i>	<b>N</b> : <i>A</i>	<b>N</b> : <i>A</i>	<b>T</b> : <i>A</i>	<b>N</b> : <i>A</i>
<b>N</b> : <i>B</i>	<b>N</b> : <i>A</i>	<b>F</b> : <i>A</i>	<b>NT</b> : <i>A</i>	<b>N</b> : <i>A</i>
<b>F</b> : <i>B</i>	<b>T</b> : <i>A</i>	<b>T</b> : <i>A</i>	<b>FT</b> : <i>A</i>	<b>F</b> : <i>A</i>
<b>T</b> : <i>B</i>	<b>NT</b> : <i>A</i>	<b>FT</b> : <i>A</i>	<b>NFT</b> : <i>A</i>	<b>T</b> : <i>A</i>
<b>NF</b> : <i>B</i>	<b>N</b> : <i>B</i>	<b>N</b> : <i>B</i>	<b>T</b> : <i>B</i>	<b>NF</b> : <i>A</i>
<b>NT</b> : <i>B</i>	<b>N</b> : <i>B</i>	<b>F</b> : <i>B</i>	<b>NT</b> : <i>B</i>	<b>NT</b> : <i>A</i>
<b>FT</b> : <i>B</i>	<b>T</b> : <i>B</i>	<b>T</b> : <i>B</i>	<b>FT</b> : <i>B</i>	<b>FT</b> : <i>A</i>
<b>NFT</b> : <i>B</i>	<b>NT</b> : <i>B</i>	<b>FT</b> : <i>B</i>	<b>NFT</b> : <i>B</i>	<b>NFT</b> : <i>A</i>

$B: A \vee_t B$					$NF: A \vee_t B$	
<b>N</b> : <i>B</i>	<b>N</b> : <i>A</i>	<b>N</b> : <i>A</i>	<b>B</b> : <i>A</i>	<b>N</b> : <i>A</i>	<b>NF</b> : <i>B</i>	<b>NF</b> : <i>A</i>
<b>N</b> : <i>B</i>	<b>N</b> : <i>A</i>	<b>F</b> : <i>A</i>	<b>NB</b> : <i>A</i>	<b>N</b> : <i>A</i>		
<b>F</b> : <i>B</i>	<b>B</b> : <i>A</i>	<b>B</b> : <i>A</i>	<b>FB</b> : <i>A</i>	<b>F</b> : <i>A</i>		
<b>B</b> : <i>B</i>	<b>NB</b> : <i>A</i>	<b>FB</b> : <i>A</i>	<b>NFB</b> : <i>A</i>	<b>B</b> : <i>A</i>		
<b>NF</b> : <i>B</i>	<b>N</b> : <i>B</i>	<b>N</b> : <i>B</i>	<b>B</b> : <i>B</i>	<b>NF</b> : <i>A</i>		
<b>NB</b> : <i>B</i>	<b>N</b> : <i>B</i>	<b>F</b> : <i>B</i>	<b>NB</b> : <i>B</i>	<b>NB</b> : <i>A</i>		
<b>FB</b> : <i>B</i>	<b>B</b> : <i>B</i>	<b>B</b> : <i>B</i>	<b>FB</b> : <i>B</i>	<b>FB</b> : <i>A</i>		
<b>NFB</b> : <i>B</i>	<b>NB</b> : <i>B</i>	<b>FB</b> : <i>B</i>	<b>NFB</b> : <i>B</i>	<b>NFB</b> : <i>A</i>		

$$NT: A \vee_t B$$

<b>N</b> : <i>B</i>	<b>N</b> : <i>A</i>	<b>NT</b> : <i>A</i>	<b>N</b> : <i>A</i>
<b>NF</b> : <i>B</i>	<b>NT</b> : <i>A</i>	<b>NFT</b> : <i>A</i>	<b>NF</b> : <i>A</i>
<b>NT</b> : <i>B</i>	<b>N</b> : <i>B</i>	<b>NT</b> : <i>B</i>	<b>NT</b> : <i>A</i>
<b>NFT</b> : <i>B</i>	<b>NT</b> : <i>B</i>	<b>NFT</b> : <i>B</i>	<b>NFT</b> : <i>A</i>

$$FT: A \vee_t B$$

<b>F</b> : <i>B</i>	<b>F</b> : <i>A</i>	<b>FT</b> : <i>A</i>	<b>F</b> : <i>A</i>
<b>NF</b> : <i>B</i>	<b>FT</b> : <i>A</i>	<b>NFT</b> : <i>A</i>	<b>NF</b> : <i>A</i>
<b>FT</b> : <i>B</i>	<b>F</b> : <i>B</i>	<b>FT</b> : <i>B</i>	<b>FT</b> : <i>A</i>
<b>NFT</b> : <i>B</i>	<b>FT</b> : <i>B</i>	<b>NFT</b> : <i>B</i>	<b>NFT</b> : <i>A</i>

$$NB: A \vee_t B$$

<b>N</b> : <i>B</i>	<b>N</b> : <i>A</i>	<b>NB</b> : <i>A</i>	<b>N</b> : <i>A</i>
<b>NF</b> : <i>B</i>	<b>NB</b> : <i>A</i>	<b>NFB</b> : <i>A</i>	<b>NF</b> : <i>A</i>
<b>NB</b> : <i>B</i>	<b>N</b> : <i>B</i>	<b>NB</b> : <i>B</i>	<b>NB</b> : <i>A</i>
<b>NFB</b> : <i>B</i>	<b>NB</b> : <i>B</i>	<b>NFB</b> : <i>B</i>	<b>NFB</b> : <i>A</i>

$$FB: A \vee_t B$$

<b>F</b> : <i>B</i>	<b>F</b> : <i>A</i>	<b>FB</b> : <i>A</i>	<b>F</b> : <i>A</i>
<b>NF</b> : <i>B</i>	<b>FB</b> : <i>A</i>	<b>NFB</b> : <i>A</i>	<b>NF</b> : <i>A</i>
<b>FB</b> : <i>B</i>	<b>F</b> : <i>B</i>	<b>FB</b> : <i>B</i>	<b>FB</b> : <i>A</i>
<b>NFB</b> : <i>B</i>	<b>FB</b> : <i>B</i>	<b>NFB</b> : <i>B</i>	<b>NFB</b> : <i>A</i>



**TB:  $A \vee_t B$**

<b>N: A</b>	<b>N: A</b>	<b>T: A</b>	<b>B: A</b>
<b>N: A</b>	<b>F: A</b>	<b>NT: A</b>	<b>NB: A</b>
<b>T: A</b>	<b>T: A</b>	<b>FT: A</b>	<b>FB: A</b>
<b>B: A</b>	<b>B: A</b>	<b>TB: A</b>	<b>TB: A</b>
<b>NT: A</b>	<b>FT: A</b>	<b>NFT: A</b>	<b>NFB: A</b>
<b>NB: A</b>	<b>FB: A</b>	<b>NTB: A</b>	<b>NTB: A</b>
<b>TB: A</b>	<b>TB: A</b>	<b>FTB: A</b>	<b>FTB: A</b>
<b>NTB: A</b>	<b>FTB: A</b>	<b>A: A</b>	<b>A: A</b>
<b>N: B</b>	<b>N: B</b>	<b>T: B</b>	<b>B: B</b>
<b>N: B</b>	<b>F: B</b>	<b>NT: B</b>	<b>NB: B</b>
<b>T: B</b>	<b>T: B</b>	<b>FT: B</b>	<b>FB: B</b>
<b>B: B</b>	<b>B: B</b>	<b>TB: B</b>	<b>TB: B</b>
<b>NT: B</b>	<b>FT: B</b>	<b>NFT: B</b>	<b>NFB: B</b>
<b>NB: B</b>	<b>FB: B</b>	<b>NTB: B</b>	<b>NTB: B</b>
<b>TB: B</b>	<b>TB: B</b>	<b>FTB: B</b>	<b>FTB: B</b>
<b>NTB: B</b>	<b>FTB: B</b>	<b>A: B</b>	<b>A: B</b>

**NFT:  $A \vee_t B$**

<b>NF: B</b>	<b>NFT: A</b>	<b>NF: A</b>
<b>NFT: B</b>	<b>NFT: B</b>	<b>NFT: A</b>

**NFB:  $A \vee_t B$**

<b>NF: B</b>	<b>NFB: A</b>	<b>NF: A</b>
<b>NFB: B</b>	<b>NFB: B</b>	<b>NFB: A</b>

**NTB:  $A \vee_t B$**

<b>N: B</b>	<b>N: A</b>	<b>NT: A</b>	<b>NB: A</b>	<b>N: A</b>
<b>NF: B</b>	<b>NT: A</b>	<b>NFT: A</b>	<b>NFB: A</b>	<b>NF: A</b>
<b>NT: B</b>	<b>NB: A</b>	<b>NTB: A</b>	<b>NTB: A</b>	<b>NT: A</b>
<b>NB: B</b>	<b>NTB: A</b>	<b>A: A</b>	<b>A: A</b>	<b>NB: A</b>
<b>NFT: B</b>	<b>N: B</b>	<b>NT: B</b>	<b>NB: B</b>	<b>NFT: A</b>
<b>NFB: B</b>	<b>NT: B</b>	<b>NFT: B</b>	<b>NFB: B</b>	<b>NFB: A</b>
<b>NTB: B</b>	<b>NB: B</b>	<b>NTB: B</b>	<b>NTB: B</b>	<b>NTB: A</b>
<b>A: B</b>	<b>NTB: B</b>	<b>A: B</b>	<b>A: B</b>	<b>A: A</b>

**FTB:  $A \vee_t B$**

<b>F: B</b>	<b>F: A</b>	<b>FT: A</b>	<b>FB: A</b>	<b>F: A</b>
<b>NF: B</b>	<b>FT: A</b>	<b>NFT: A</b>	<b>NFB: A</b>	<b>NF: A</b>
<b>FT: B</b>	<b>FB: A</b>	<b>FTB: A</b>	<b>FTB: A</b>	<b>FT: A</b>
<b>FB: B</b>	<b>FTB: A</b>	<b>A: A</b>	<b>A: A</b>	<b>FB: A</b>
<b>NFT: B</b>	<b>F: B</b>	<b>FT: B</b>	<b>FB: B</b>	<b>NFT: A</b>
<b>NFB: B</b>	<b>FT: B</b>	<b>NFT: B</b>	<b>NFB: B</b>	<b>NFB: A</b>
<b>FTB: B</b>	<b>FB: B</b>	<b>FTB: B</b>	<b>FTB: B</b>	<b>FTB: A</b>
<b>A: B</b>	<b>FTB: B</b>	<b>A: B</b>	<b>A: B</b>	<b>A: A</b>

$$\frac{\mathbf{A}: A \vee_t B}{\begin{array}{cccc} \mathbf{NF}: B & \mathbf{NFT}: A & \mathbf{NFB}: A & \mathbf{NF}: A \\ \mathbf{NFT}: B & \mathbf{A}: A & \mathbf{A}: A & \mathbf{NFT}: A \\ \mathbf{NFB}: B & \mathbf{NFT}: B & \mathbf{NFB}: B & \mathbf{NFB}: A \\ \mathbf{A}: B & \mathbf{A}: B & \mathbf{A}: B & \mathbf{A}: A \end{array}}$$

The branch expansion rules for connective  $\wedge_f$  are given by

$$\frac{\mathbf{N}: A \wedge_f B}{\begin{array}{cccc} \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A \\ \mathbf{F}: B & \mathbf{F}: A & \mathbf{B}: A & \mathbf{F}: A \\ \mathbf{B}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{B}: A \\ \mathbf{FB}: B & \mathbf{F}: B & \mathbf{B}: B & \mathbf{FB}: A \end{array}}$$

$$\frac{\mathbf{N}: A \wedge_f B}{\begin{array}{ccccc} \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A & \mathbf{N}: A \\ \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{NF}: A & \mathbf{N}: A \\ \mathbf{F}: B & \mathbf{F}: A & \mathbf{B}: A & \mathbf{NB}: A & \mathbf{F}: A \\ \mathbf{B}: B & \mathbf{NF}: A & \mathbf{NB}: A & \mathbf{NFB}: A & \mathbf{B}: A \\ \mathbf{NF}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{NF}: A \\ \mathbf{NB}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{NF}: B & \mathbf{NB}: A \\ \mathbf{FB}: B & \mathbf{F}: B & \mathbf{B}: B & \mathbf{NB}: B & \mathbf{FB}: A \\ \mathbf{NFB}: B & \mathbf{NF}: B & \mathbf{NB}: B & \mathbf{NFB}: B & \mathbf{NFB}: A \end{array}}$$

$$\frac{\mathbf{F}: A \wedge_f B}{\begin{array}{ccc} \mathbf{F}: B & \mathbf{F}: A & \mathbf{F}: A \\ \mathbf{FB}: B & \mathbf{F}: B & \mathbf{FB}: A \end{array}}$$

$$\frac{\mathbf{T}: A \wedge_f B}{\begin{array}{ccccc} \mathbf{N}: B & \mathbf{N}: A & \mathbf{N}: A & \mathbf{T}: A & \mathbf{N}: A \\ \mathbf{F}: B & \mathbf{F}: A & \mathbf{T}: A & \mathbf{FT}: A & \mathbf{F}: A \\ \mathbf{T}: B & \mathbf{T}: A & \mathbf{B}: A & \mathbf{TB}: A & \mathbf{T}: A \\ \mathbf{B}: B & \mathbf{FT}: A & \mathbf{TB}: A & \mathbf{FTB}: A & \mathbf{B}: A \\ \mathbf{FT}: B & \mathbf{N}: B & \mathbf{N}: B & \mathbf{T}: B & \mathbf{FT}: A \\ \mathbf{FB}: B & \mathbf{F}: B & \mathbf{T}: B & \mathbf{FT}: B & \mathbf{FB}: A \\ \mathbf{TB}: B & \mathbf{T}: B & \mathbf{B}: B & \mathbf{TB}: B & \mathbf{TB}: A \\ \mathbf{FTB}: B & \mathbf{FT}: B & \mathbf{TB}: B & \mathbf{FTB}: B & \mathbf{FTB}: A \end{array}}$$

$$\frac{\mathbf{B}: A \wedge_f B}{\begin{array}{ccc} \mathbf{B}: B & \mathbf{B}: A & \mathbf{B}: A \\ \mathbf{FB}: B & \mathbf{B}: B & \mathbf{FB}: A \end{array}} \quad \frac{\mathbf{NF}: A \wedge_f B}{\begin{array}{cccc} \mathbf{F}: B & \mathbf{F}: A & \mathbf{NF}: A & \mathbf{F}: A \\ \mathbf{NF}: B & \mathbf{NF}: A & \mathbf{NFB}: A & \mathbf{NF}: A \\ \mathbf{FB}: B & \mathbf{F}: B & \mathbf{NF}: B & \mathbf{FB}: A \\ \mathbf{NFB}: B & \mathbf{NF}: B & \mathbf{NFB}: B & \mathbf{NFB}: A \end{array}}$$

$$\text{NT: } A \wedge_f B$$

<b>N: A</b>	<b>N: A</b>	<b>N: A</b>	<b>T: A</b>
<b>N: A</b>	<b>N: A</b>	<b>NF: A</b>	<b>NT: A</b>
<b>F: A</b>	<b>T: A</b>	<b>NT: A</b>	<b>FT: A</b>
<b>T: A</b>	<b>B: A</b>	<b>NB: A</b>	<b>TB: A</b>
<b>NF: A</b>	<b>NT: A</b>	<b>NFT: A</b>	<b>NFT: A</b>
<b>NT: A</b>	<b>NB: A</b>	<b>NFB: A</b>	<b>NTB: A</b>
<b>FT: A</b>	<b>TB: A</b>	<b>NTB: A</b>	<b>FTB: A</b>
<b>NFT: A</b>	<b>NTB: A</b>	<b>A: A</b>	<b>A: A</b>
<b>N: B</b>	<b>N: B</b>	<b>N: B</b>	<b>T: B</b>
<b>N: B</b>	<b>N: B</b>	<b>NF: B</b>	<b>NT: B</b>
<b>F: B</b>	<b>T: B</b>	<b>NT: B</b>	<b>FT: B</b>
<b>T: B</b>	<b>B: B</b>	<b>NB: B</b>	<b>TB: B</b>
<b>NF: B</b>	<b>NT: B</b>	<b>NFT: B</b>	<b>NFT: B</b>
<b>NT: B</b>	<b>NB: B</b>	<b>NFB: B</b>	<b>NTB: B</b>
<b>FT: B</b>	<b>TB: B</b>	<b>NTB: B</b>	<b>FTB: B</b>
<b>NFT: B</b>	<b>NTB: B</b>	<b>A: B</b>	<b>A: B</b>

$$\text{FT: } A \wedge_f B$$

<b>F: B</b>	<b>F: A</b>	<b>FT: A</b>	<b>F: A</b>
<b>FT: B</b>	<b>FT: A</b>	<b>FTB: A</b>	<b>FT: A</b>
<b>FB: B</b>	<b>F: B</b>	<b>FT: B</b>	<b>FB: A</b>
<b>FTB: B</b>	<b>FT: B</b>	<b>FTB: B</b>	<b>FTB: A</b>

$$\text{NB: } A \wedge_f B$$

<b>B: B</b>	<b>B: A</b>	<b>NB: A</b>	<b>B: A</b>
<b>NB: B</b>	<b>NB: A</b>	<b>NFB: A</b>	<b>NB: A</b>
<b>FB: B</b>	<b>B: B</b>	<b>NB: B</b>	<b>FB: A</b>
<b>NFB: B</b>	<b>NB: B</b>	<b>NFB: B</b>	<b>NFB: A</b>

$$\text{FB: } A \wedge_f B$$

<b>FB: B</b>	<b>FB: A</b>
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$$\text{TB: } A \wedge_f B$$

<b>B: B</b>	<b>B: A</b>	<b>TB: A</b>	<b>B: A</b>
<b>FB: B</b>	<b>TB: A</b>	<b>FTB: A</b>	<b>FB: A</b>
<b>TB: B</b>	<b>B: B</b>	<b>TB: B</b>	<b>TB: A</b>
<b>FTB: B</b>	<b>TB: B</b>	<b>FTB: B</b>	<b>FTB: A</b>

$$\text{NFT: } A \wedge_f B$$

<b>F: B</b>	<b>F: A</b>	<b>NF: A</b>	<b>FT: A</b>	<b>F: A</b>
<b>NF: B</b>	<b>NF: A</b>	<b>NFT: A</b>	<b>NFT: A</b>	<b>NF: A</b>
<b>FT: B</b>	<b>FT: A</b>	<b>NFB: A</b>	<b>FTB: A</b>	<b>FT: A</b>
<b>FB: B</b>	<b>NFT: A</b>	<b>A: A</b>	<b>A: A</b>	<b>FB: A</b>
<b>NFT: B</b>	<b>F: B</b>	<b>NF: B</b>	<b>FT: B</b>	<b>NFT: A</b>
<b>NFB: B</b>	<b>NF: B</b>	<b>NFT: B</b>	<b>NFT: B</b>	<b>NFB: A</b>
<b>FTB: B</b>	<b>FT: B</b>	<b>NFB: B</b>	<b>FTB: B</b>	<b>FTB: A</b>
<b>A: B</b>	<b>NFT: B</b>	<b>A: B</b>	<b>A: B</b>	<b>A: A</b>

$$\frac{\text{NFB: } A \wedge_f B}{\begin{array}{ccc} \text{FB: } B & \text{NFB: } A & \text{FB: } A \\ \text{NFB: } B & \text{NFB: } B & \text{NFB: } A \end{array}}$$

$$\frac{\text{NTB: } A \wedge_f B}{\begin{array}{ccccc} \text{B: } B & \text{B: } A & \text{NB: } A & \text{TB: } A & \text{B: } A \\ \text{NB: } B & \text{NB: } A & \text{NFB: } A & \text{NTB: } A & \text{NB: } A \\ \text{FB: } B & \text{TB: } A & \text{NTB: } A & \text{FTB: } A & \text{FB: } A \\ \text{TB: } B & \text{NTB: } A & \text{A: } A & \text{A: } A & \text{TB: } A \\ \text{NFB: } B & \text{B: } B & \text{NB: } B & \text{TB: } B & \text{NFB: } A \\ \text{NTB: } B & \text{NB: } B & \text{NFB: } B & \text{NTB: } B & \text{NTB: } A \\ \text{FTB: } B & \text{TB: } B & \text{NTB: } B & \text{FTB: } B & \text{FTB: } A \\ \text{A: } B & \text{NTB: } B & \text{A: } B & \text{A: } B & \text{A: } A \end{array}}$$

$$\frac{\text{FTB: } A \wedge_f B}{\begin{array}{ccc} \text{FB: } B & \text{FTB: } A & \text{FB: } A \\ \text{FTB: } B & \text{FTB: } B & \text{FTB: } A \end{array}}$$

$$\frac{\text{A: } A \wedge_f B}{\begin{array}{cccc} \text{FB: } B & \text{NFB: } A & \text{FTB: } A & \text{FB: } A \\ \text{NFB: } B & \text{A: } A & \text{A: } A & \text{NFB: } A \\ \text{FTB: } B & \text{NFB: } B & \text{FTB: } B & \text{FTB: } A \\ \text{A: } B & \text{A: } B & \text{A: } B & \text{A: } A \end{array}}$$

The branch expansion rules for connective  $\vee_f$  are given by

$$\frac{\text{N: } A \vee_f B}{\begin{array}{cccc} \text{N: } B & \text{N: } A & \text{N: } A & \text{N: } A \\ \text{N: } B & \text{N: } A & \text{T: } A & \text{N: } A \\ \text{T: } B & \text{N: } B & \text{N: } B & \text{T: } A \\ \text{NT: } B & \text{N: } B & \text{T: } B & \text{NT: } A \end{array}} \quad \frac{\text{N: } A \vee_f B}{\begin{array}{ccc} \text{N: } B & \text{N: } A & \text{N: } A \\ \text{NT: } B & \text{N: } B & \text{NT: } A \end{array}}$$

$$\frac{\text{F: } A \vee_f B}{\begin{array}{ccccc} \text{N: } B & \text{N: } A & \text{N: } A & \text{F: } A & \text{N: } A \\ \text{N: } B & \text{N: } A & \text{F: } A & \text{NF: } A & \text{N: } A \\ \text{F: } B & \text{F: } A & \text{T: } A & \text{FT: } A & \text{F: } A \\ \text{T: } B & \text{NF: } A & \text{FT: } A & \text{NFT: } A & \text{T: } A \\ \text{NF: } B & \text{N: } B & \text{N: } B & \text{F: } B & \text{NF: } A \\ \text{NT: } B & \text{N: } B & \text{F: } B & \text{NF: } B & \text{NT: } A \\ \text{FT: } B & \text{F: } B & \text{T: } B & \text{FT: } B & \text{FT: } A \\ \text{NFT: } B & \text{NF: } B & \text{FT: } B & \text{NFT: } B & \text{NFT: } A \end{array}}$$

$$\frac{\text{T: } A \vee_f B}{\begin{array}{ccc} \text{T: } B & \text{T: } A & \text{T: } A \\ \text{NT: } B & \text{T: } B & \text{NT: } A \end{array}}$$

$B: A \vee_f B$					
$N: B$	$N: A$	$N: A$	$B: A$	$N: A$	
$N: B$	$N: A$	$T: A$	$NB: A$	$N: A$	
$T: B$	$B: A$	$B: A$	$TB: A$	$T: A$	
$B: B$	$NB: A$	$TB: A$	$NTB: A$	$B: A$	
$NT: B$	$N: B$	$N: B$	$B: B$	$NT: A$	
$NB: B$	$N: B$	$T: B$	$NB: B$	$NB: A$	
$TB: B$	$B: B$	$B: B$	$TB: B$	$TB: A$	
$NTB: B$	$NB: B$	$TB: B$	$NTB: B$	$NTB: A$	
$NF: A \vee_f B$			$NT: A \vee_f B$		
$N: B$	$N: A$	$NF: A$	$N: A$	$NT: B$	$NT: A$
$NF: B$	$NF: A$	$NFT: A$	$NF: A$		
$NT: B$	$N: B$	$NF: B$	$NT: A$		
$NFT: B$	$NF: B$	$NFT: B$	$NFT: A$		
$FT: A \vee_f B$					
$T: B$	$T: A$	$FT: A$	$T: A$		
$NT: B$	$FT: A$	$NFT: A$	$NT: A$		
$FT: B$	$T: B$	$FT: B$	$FT: A$		
$NFT: B$	$FT: B$	$NFT: B$	$NFT: A$		
$NB: A \vee_f B$					
$N: B$	$N: A$	$NB: A$	$N: A$		
$NT: B$	$NB: A$	$NTB: A$	$NT: A$		
$NB: B$	$N: B$	$NB: B$	$NB: A$		
$NTB: B$	$NB: B$	$NTB: B$	$NTB: A$		
$FB: A \vee_f B$					
$N: A$	$N: A$	$F: A$	$B: A$		
$N: A$	$F: A$	$NF: A$	$NB: A$		
$F: A$	$T: A$	$FT: A$	$FB: A$		
$B: A$	$B: A$	$FB: A$	$TB: A$		
$NF: A$	$FT: A$	$NFT: A$	$NFB: A$		
$NB: A$	$FB: A$	$NFB: A$	$NTB: A$		
$FB: A$	$TB: A$	$FTB: A$	$FTB: A$		
$NFB: A$	$FTB: A$	$A: A$	$A: A$		
$N: B$	$N: B$	$F: B$	$B: B$		
$N: B$	$F: B$	$NF: B$	$NB: B$		
$F: B$	$T: B$	$FT: B$	$FB: B$		
$B: B$	$B: B$	$FB: B$	$TB: B$		
$NF: B$	$FT: B$	$NFT: B$	$NFB: B$		
$NB: B$	$FB: B$	$NFB: B$	$NTB: B$		
$FB: B$	$TB: B$	$FTB: B$	$FTB: B$		
$NFB: B$	$FTB: B$	$A: B$	$A: B$		

$$\frac{\text{TB}: A \vee_f B}{\begin{array}{cccc} \text{T}: B & \text{T}: A & \text{TB}: A & \text{T}: A \\ \text{NT}: B & \text{TB}: A & \text{NTB}: A & \text{NT}: A \\ \text{TB}: B & \text{T}: B & \text{TB}: B & \text{TB}: A \\ \text{NTB}: B & \text{TB}: B & \text{NTB}: B & \text{NTB}: A \end{array}}$$

$$\frac{\text{NFT}: A \vee_f B}{\begin{array}{ccc} \text{NT}: B & \text{NFT}: A & \text{NT}: A \\ \text{NFT}: B & \text{NFT}: B & \text{NFT}: A \end{array}}$$

$$\frac{\text{NFB}: A \vee_f B}{\begin{array}{ccccc} \text{N}: B & \text{N}: A & \text{NF}: A & \text{NB}: A & \text{N}: A \\ \text{NF}: B & \text{NF}: A & \text{NFT}: A & \text{NFB}: A & \text{NF}: A \\ \text{NT}: B & \text{NB}: A & \text{NFB}: A & \text{NTB}: A & \text{NT}: A \\ \text{NB}: B & \text{NFB}: A & \text{A}: A & \text{A}: A & \text{NB}: A \\ \text{NFT}: B & \text{N}: B & \text{NF}: B & \text{NB}: B & \text{NFT}: A \\ \text{NFB}: B & \text{NF}: B & \text{NFT}: B & \text{NFB}: B & \text{NFB}: A \\ \text{NTB}: B & \text{NB}: B & \text{NFB}: B & \text{NTB}: B & \text{NTB}: A \\ \text{A}: B & \text{NFB}: B & \text{A}: B & \text{A}: B & \text{A}: A \end{array}}$$

$$\frac{\text{NTB}: A \vee_f B}{\begin{array}{ccc} \text{NT}: B & \text{NTB}: A & \text{NT}: A \\ \text{NTB}: B & \text{NTB}: B & \text{NTB}: A \end{array}}$$

$$\frac{\text{FTB}: A \vee_f B}{\begin{array}{ccccc} \text{T}: B & \text{T}: A & \text{FT}: A & \text{TB}: A & \text{T}: A \\ \text{NT}: B & \text{FT}: A & \text{NFT}: A & \text{NTB}: A & \text{NT}: A \\ \text{FT}: B & \text{TB}: A & \text{FTB}: A & \text{FTB}: A & \text{FT}: A \\ \text{TB}: B & \text{FTB}: A & \text{A}: A & \text{A}: A & \text{TB}: A \\ \text{NFT}: B & \text{T}: B & \text{FT}: B & \text{TB}: B & \text{NFT}: A \\ \text{NTB}: B & \text{FT}: B & \text{NFT}: B & \text{NTB}: B & \text{NTB}: A \\ \text{FTB}: B & \text{TB}: B & \text{FTB}: B & \text{FTB}: B & \text{FTB}: A \\ \text{A}: B & \text{FTB}: B & \text{A}: B & \text{A}: B & \text{A}: A \end{array}}$$

$$\frac{\text{A}: A \vee_f B}{\begin{array}{cccc} \text{NT}: B & \text{NFT}: A & \text{NTB}: A & \text{NT}: A \\ \text{NFT}: B & \text{A}: A & \text{A}: A & \text{NFT}: A \\ \text{NTB}: B & \text{NFT}: B & \text{NTB}: B & \text{NTB}: A \\ \text{A}: B & \text{A}: B & \text{A}: B & \text{A}: A \end{array}}$$

**Definition 18.** An interpretation  $\mathfrak{I}$  *satisfies* a signed formula  $v: A$  iff  $\text{val}_{\mathfrak{I}}(A) \neq v$ . A set of signed formulas is *satisfiable* if some interpretation  $\mathfrak{I}$  satisfies all signed formulas in it.

**Theorem 19.** *A set of signed formulas is unsatisfiable iff it has a closed tableau.*

*Proof.* Apply Theorems 4.14 and 4.21 of Hähnle [7]; interpreting  $v: A$  as  $S \ A$  where  $S = V \setminus \{v\}$ .  $\square$

**Corollary 20.** *In Shramko-Wansing logic,  $\models A$  iff  $\{v: A \mid v \in V^+\}$  has a closed tableau.  $\Delta \models A$  iff  $\{v: B \mid v \in V^-, B \in \Delta\} \cup \{v: A \mid v \in V^+\}$  has a closed tableau.*

## 5 Natural deduction for Shramko-Wansing logic

Let  $\Gamma$  be a (set) sequent,  $V^+ \subseteq V$  the set of *designated truth values*. The set of non-designated truth values is then  $V^- = V \setminus V^+$ . We divide the sequent  $\Gamma$  into its designated part  $\Gamma^+$  and its non-designated part  $\Gamma^-$  in the obvious way:

$$\begin{aligned}\Gamma^+ &:= \langle \Gamma_v \mid v \in V^+ \rangle \\ \Gamma^- &:= \langle \Gamma_v \mid v \in V^- \rangle\end{aligned}$$

**Definition 21.** The *natural deduction rules* for Shramko-Wansing logic are given by:

1. A weakening rule for all  $v \in V^+$ :

$$\frac{\Gamma^+}{\Gamma^+, [v: A]} \text{ w:v}$$

2. For every connective  $\square$  and every truth value  $v$  an introduction rule  $\square I: v$  (if  $v \in V^+$ ) or an elimination rule  $\square E: v$  (if  $v \in V^-$ ).

The introduction and elimination rules for connective  $\neg_t$  are given by

$$\begin{array}{c} \frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{N}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{N}} \\ \\ \frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{N}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+, [\mathbf{T}: A]}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{N}} \\ \\ \frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{B}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{F}} \\ \\ \frac{\Gamma_1^-, [[\mathbf{N}: A]]}{\frac{\Gamma_1^+}{\Gamma_1^+, [\mathbf{T}: \neg_t A]} \neg_t I: \mathbf{T}} \\ \\ \frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{F}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{B}} \end{array}$$

$$\begin{array}{c}
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+, [\mathbf{TB}: A]}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{NF}} \\
\frac{\Gamma_1^-}{\frac{\Gamma_1^+, [\mathbf{NT}: A]}{\Gamma_1^+, [\mathbf{NT}: \neg_t A]} \neg_t I: \mathbf{NT}} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{NB}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{FT}} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{FT}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{NB}} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{FB}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{FB}} \\
\frac{\Gamma_1^-, [[\mathbf{NF}: A]]}{\frac{\Gamma_1^+}{\Gamma_1^+, [\mathbf{TB}: \neg_t A]} \neg_t I: \mathbf{TB}} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+, [\mathbf{NTB}: A]}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{NFT}} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{FTB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{NFB}} \\
\frac{\Gamma_1^-, [[\mathbf{NFT}, \mathbf{FTB}: A]]}{\frac{\Gamma_1^+}{\Gamma_1^+, [\mathbf{NTB}: \neg_t A]} \neg_t I: \mathbf{NTB}} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{NFB}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{FTB}} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}: \neg_t A]] \quad \Gamma_1^-, [[\mathbf{A}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_t A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_t E: \mathbf{A}}
\end{array}$$



The introduction and elimination rules for connective  $\neg_f$  are given by

$$\begin{array}{c}
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{N}: A]]}{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+} \neg_f E: \mathbf{N} \\
\Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{F}: A]] \\
\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: \mathbf{N} \\
\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{N}: A]] \\
\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: \mathbf{F} \\
\frac{\Gamma_1^-, [[\mathbf{B}: A]] \quad \Gamma_1^+}{\Gamma_1^+, [\mathbf{T}: \neg_f A]} \neg_f I: \mathbf{T} \\
\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [\mathbf{T}: A] \\
\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: \mathbf{B} \\
\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{NF}: A]] \\
\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: \mathbf{NF} \\
\frac{\Gamma_1^-, [[\mathbf{FB}: A]] \quad \Gamma_1^+}{\Gamma_1^+, [\mathbf{NT}: \neg_f A]} \neg_f I: \mathbf{NT} \\
\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{NB}: A]] \\
\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: \mathbf{FT} \\
\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{FT}: A]] \\
\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: \mathbf{NB} \\
\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [\mathbf{NT}: A] \\
\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A] \quad \Gamma_1^+}{\Gamma_0^+, \Gamma_1^+} \neg_f E: \mathbf{FB}
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma_1^-}{\Gamma_1^+, [\mathbf{TB}: A]} \neg_f \mathbf{I}: \mathbf{TB} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{NFT}, \mathbf{NFB}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A]}{\Gamma_0^+, \Gamma_1^+} \quad \Gamma_1^+} \neg_f \mathbf{E}: \mathbf{NFT} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{FTB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A]}{\Gamma_0^+, \Gamma_1^+} \quad \Gamma_1^+} \neg_f \mathbf{E}: \mathbf{NFB} \\
\frac{\Gamma_1^-, [[\mathbf{FTB}: A]]}{\Gamma_1^+, [\mathbf{NTB}: \neg_f A]} \neg_f \mathbf{I}: \mathbf{NTB} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: \neg_f A]] \quad \Gamma_1^-}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A]}{\Gamma_0^+, \Gamma_1^+} \quad \Gamma_1^+, [\mathbf{NTB}: A]} \neg_f \mathbf{E}: \mathbf{FTB} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}: \neg_f A]] \quad \Gamma_1^-, [[\mathbf{A}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: \neg_f A]}{\Gamma_0^+, \Gamma_1^+} \quad \Gamma_1^+} \neg_f \mathbf{E}: \mathbf{A}
\end{array}$$

The introduction and elimination rules for connective  $\wedge_t$  are given by

$$\begin{array}{c}
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_t B]] \quad \Gamma_1^-, [[\mathbf{N}, \mathbf{B}: B]] \quad \Gamma_2^-, [[\mathbf{N}: A, B]] \quad \Gamma_3^-, [[\mathbf{N}, \mathbf{B}: A, B]] \quad \Gamma_4^-, [[\mathbf{N}, \mathbf{B}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_t B]}{\Gamma_0^+, \dots, \Gamma_4^+} \quad \Gamma_1^+, [\mathbf{T}, \mathbf{TB}: B] \quad \Gamma_2^+, [\mathbf{T}: A, B] \quad \Gamma_3^+ \quad \Gamma_4^+, [\mathbf{T}, \mathbf{TB}: A]} \wedge_t \mathbf{E}: \mathbf{N} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_t B]] \quad \Gamma_1^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: B]] \quad \Gamma_2^-, [[\mathbf{N}, \mathbf{N}: A, B]] \quad \Gamma_3^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A, B]] \quad \Gamma_4^-, [[\mathbf{N}, \mathbf{NB}: A, B]] \quad \Gamma_5^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NB}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_t B]}{\Gamma_0^+, \dots, \Gamma_5^+} \quad \Gamma_1^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: B] \quad \Gamma_2^+, [\mathbf{T}, \mathbf{NT}: A, B] \quad \Gamma_3^+ \quad \Gamma_4^+, [\mathbf{NT}, \mathbf{NTB}: A, B] \quad \Gamma_5^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A]} \wedge_t \mathbf{E}: \mathbf{N} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_t B]] \quad \Gamma_1^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: B]] \quad \Gamma_2^-, [[\mathbf{N}, \mathbf{F}, \mathbf{FT}, \mathbf{A}, B]] \quad \Gamma_3^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FB}: A, B]] \quad \Gamma_4^-, [[\mathbf{F}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A, B]] \quad \Gamma_5^-, [[\mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{FB}, \mathbf{FTB}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_t B]}{\Gamma_0^+, \dots, \Gamma_5^+} \quad \Gamma_1^+, [\mathbf{T}, \mathbf{TB}: B] \quad \Gamma_2^+, [\mathbf{T}: A, B] \quad \Gamma_3^+ \quad \Gamma_4^+ \quad \Gamma_5^+, [\mathbf{T}, \mathbf{TB}: A]} \wedge_t \mathbf{E}: \mathbf{F} \\
\frac{\Gamma_1^- \quad \Gamma_2^- \quad \Gamma_3^-}{\Gamma_1^+, [\mathbf{T}, \mathbf{TB}: B] \quad \Gamma_2^+, [\mathbf{T}: A, B] \quad \Gamma_3^+, [\mathbf{T}, \mathbf{TB}: A]} \wedge_t \mathbf{I}: \mathbf{T} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_t B]] \quad \Gamma_1^-, [[\mathbf{B}: B]] \quad \Gamma_2^-, [[\mathbf{B}: A, B]] \quad \Gamma_3^-, [[\mathbf{B}: A]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_t B]}{\Gamma_0^+, \dots, \Gamma_3^+} \quad \Gamma_1^+, [\mathbf{TB}: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [\mathbf{TB}: A]} \wedge_t \mathbf{E}: \mathbf{B} \\
\frac{\Gamma_0^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{FT}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A \wedge_t B]] \quad \Gamma_1^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{NFT}: A, B]] \quad \Gamma_2^-, [[\mathbf{N}, \mathbf{N}, \mathbf{F}, \mathbf{B}, \mathbf{NF}, \mathbf{NB}, \mathbf{FB}, \mathbf{NFB}: A, B]] \quad \Gamma_3^-, [[\mathbf{N}, \mathbf{NF}, \mathbf{NB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{A}: A, B]] \quad \Gamma_4^-, [[\mathbf{F}, \mathbf{NF}, \mathbf{FT}, \mathbf{FB}, \mathbf{NFT}, \mathbf{NFB}, \mathbf{FTB}, \mathbf{A}: A, B]]}{\frac{\Gamma_0^+, [\mathbf{T}, \mathbf{NT}, \mathbf{TB}, \mathbf{NTB}: A \wedge_t B]}{\Gamma_0^+, \dots, \Gamma_4^+} \quad \Gamma_1^+, [\mathbf{T}, \mathbf{NT}: A, B] \quad \Gamma_2^+ \quad \Gamma_3^+, [\mathbf{NT}, \mathbf{NTB}: A, B] \quad \Gamma_4^+} \wedge_t \mathbf{E}: \mathbf{NF}
\end{array}$$



$$\frac{\Gamma_0^-, [[N, N, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \vee_t B]] \quad \Gamma_1^-, [[F, NF: B]] \quad \Gamma_2^-, [[F: A, B]] \quad \Gamma_3^-, [[F, NF: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+} \vee_t E: \mathbf{F}$$

$$\Gamma_0^+, \dots, \Gamma_3^+$$

$$\frac{\Gamma_1^-, [[N, N, F, NF, FT, NFT: B]] \quad \Gamma_2^-, [[N, N: A, B]] \quad \Gamma_3^-, [[N, F, FT: A, B]] \quad \Gamma_4^-, [[FT, NFT: A, B]] \quad \Gamma_5^-, [[N, N, F, NF, FT, NFT: A]]}{\Gamma_1^+, [T, NT: B] \quad \Gamma_2^+, [T, NT: A, B] \quad \Gamma_3^+, [T: A, B] \quad \Gamma_4^+, [T, NT: A, B] \quad \Gamma_5^+, [T, NT: A]} \vee_t I: \mathbf{T}$$

$$\Gamma_1^+, \dots, \Gamma_5^+, [T: A \vee_t B]$$

$$\frac{\Gamma_0^-, [[N, N, F, NF, FT, NB, FB, NFT, NFB, FTB, A: A \vee_t B]] \quad \Gamma_1^-, [[N, N, F, B, NF, NB, FB, NFB: B]] \quad \Gamma_2^-, [[N, N, B, NB: A, B]] \quad \Gamma_3^-, [[N, F, B, FB: A, B]] \quad \Gamma_4^-, [[B, NB, FB, NFB: A, B]] \quad \Gamma_5^-, [[N, N, F, B, NF, NB, FB, NFB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+ \quad \Gamma_5^+} \vee_t E: \mathbf{B}$$

$$\Gamma_0^+, \dots, \Gamma_5^+$$

$$\frac{\Gamma_0^-, [[N, N, F, B, FT, NB, FB, NFT, NFB, FTB, A: A \vee_t B]] \quad \Gamma_1^-, [[NF: B]] \quad \Gamma_2^-, [[NF: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+} \vee_t E: \mathbf{NF}$$

$$\Gamma_0^+, \Gamma_1^+, \Gamma_2^+$$

$$\frac{\Gamma_1^-, [[N, NF, NFT: B]] \quad \Gamma_2^-, [[N: A, B]] \quad \Gamma_3^-, [[NFT: A, B]] \quad \Gamma_4^-, [[N, NF, NFT: A]]}{\Gamma_1^+, [NT: B] \quad \Gamma_2^+, [NT: A, B] \quad \Gamma_3^+, [NT: A, B] \quad \Gamma_4^+, [NT: A]} \vee_t I: \mathbf{NT}$$

$$\Gamma_1^+, \dots, \Gamma_4^+, [NT: A \vee_t B]$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, NB, FB, NFT, NFB, FTB, A: A \vee_t B]] \quad \Gamma_1^-, [[F, NF, FT, NFT: B]] \quad \Gamma_2^-, [[F, FT: A, B]] \quad \Gamma_3^-, [[FT, NFT: A, B]] \quad \Gamma_4^-, [[F, NF, FT, NFT: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+} \vee_t E: \mathbf{FT}$$

$$\Gamma_0^+, \dots, \Gamma_4^+$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, FB, NFT, NFB, FTB, A: A \vee_t B]] \quad \Gamma_1^-, [[N, NF, NB, NFB: B]] \quad \Gamma_2^-, [[N, NB: A, B]] \quad \Gamma_3^-, [[NB, NFB: A, B]] \quad \Gamma_4^-, [[N, NF, NB, NFB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+} \vee_t E: \mathbf{NB}$$

$$\Gamma_0^+, \dots, \Gamma_4^+$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, NFT, NFB, FTB, A: A \vee_t B]] \quad \Gamma_1^-, [[F, NF, FB, NFB: B]] \quad \Gamma_2^-, [[F, FB: A, B]] \quad \Gamma_3^-, [[FB, NFB: A, B]] \quad \Gamma_4^-, [[F, NF, FB, NFB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+} \vee_t E: \mathbf{FB}$$

$$\Gamma_0^+, \dots, \Gamma_4^+$$

$$\frac{\Gamma_1^-, [[N, N, B, NB: A, B]] \quad \Gamma_2^-, [[N, F, B, FT, FB, FTB: A, B]] \quad \Gamma_3^-, [[FT, NFT, FTB, A: A, B]] \quad \Gamma_4^-, [[B, NB, FB, NFB, FTB, A: A, B]]}{\Gamma_1^+, [T, NT, TB, NTB: A, B] \quad \Gamma_2^+, [T, TB: A, B] \quad \Gamma_3^+, [T, NT, TB, NTB: A, B] \quad \Gamma_4^+, [TB, NTB: A, B]} \vee_t I: \mathbf{TB}$$

$$\Gamma_1^+, \dots, \Gamma_4^+, [TB: A \vee_t B]$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFB, FTB, A: A \vee_t B]] \quad \Gamma_1^-, [[NF, NFT: B]] \quad \Gamma_2^-, [[NFT: A, B]] \quad \Gamma_3^-, [[NF, NFT: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+} \vee_t E: \mathbf{NFT}$$

$$\Gamma_0^+, \dots, \Gamma_3^+$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFT, FTB, A: A \vee_t B]] \quad \Gamma_1^-, [[NF, NFB: B]] \quad \Gamma_2^-, [[NFB: A, B]] \quad \Gamma_3^-, [[NF, NFB: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+} \vee_t E: \mathbf{NFB}$$

$$\Gamma_0^+, \dots, \Gamma_3^+$$

$$\frac{\Gamma_1^-, [[N, NF, NB, NFT, NFB, A: B]] \quad \Gamma_2^-, [[N, NB: A, B]] \quad \Gamma_3^-, [[NFT, A: A, B]] \quad \Gamma_4^-, [[NB, NFB, A: A, B]] \quad \Gamma_5^-, [[N, NF, NB, NFT, NFB, A: A]]}{\Gamma_1^+, [NT, NTB: B] \quad \Gamma_2^+, [NT, NTB: A, B] \quad \Gamma_3^+, [NT, NTB: A, B] \quad \Gamma_4^+, [NTB: A, B] \quad \Gamma_5^+, [NT, NTB: A]} \vee_t I: \mathbf{NTB}$$

$$\Gamma_1^+, \dots, \Gamma_5^+, [NTB: A \vee_t B]$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFT, NFB, A: A \vee_t B]] \quad \Gamma_1^-, [[F, NF, FT, FB, NFT, NFB, FTB, A: B]] \quad \Gamma_2^-, [[F, FT, FB, FTB: A, B]] \quad \Gamma_3^-, [[FT, NFT, FTB, A: A, B]] \quad \Gamma_4^-, [[FB, NFB, FTB, A: A, B]] \quad \Gamma_5^-, [[F, NF, FT, FB, NFT, NFB, FTB, A: A]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+ \quad \Gamma_5^+} \vee_t E: \mathbf{FTB}$$

$$\Gamma_0^+, \dots, \Gamma_5^+$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A \vee_t B]] \quad \Gamma_1^-, [[NF, NFT, NFB, A: B]] \quad \Gamma_2^-, [[NFT, A: A, B]] \quad \Gamma_3^-, [[NFB, A: A, B]] \quad \Gamma_4^-, [[NF, NFT, NFB, A: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \vee_t B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+}{\Gamma_0^+, \dots, \Gamma_4^+} \vee_t E: A}$$

The introduction and elimination rules for connective  $\wedge_f$  are given by

$$\frac{\Gamma_0^-, [[N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[N, F, B, FB: B]] \quad \Gamma_2^-, [[N, F: A, B]] \quad \Gamma_3^-, [[N, B: A, B]] \quad \Gamma_4^-, [[N, F, B, FB: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+}{\Gamma_0^+, \dots, \Gamma_4^+} \wedge_f E: N}$$

$$\frac{\Gamma_0^-, [[N, F, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[N, N, F, B, NF, NB, FB, NFB: B]] \quad \Gamma_2^-, [[N, N, F, NF: A, B]] \quad \Gamma_3^-, [[N, N, B, NB: A, B]] \quad \Gamma_4^-, [[N, NF, NB, NFB: A, B]] \quad \Gamma_5^-, [[N, N, F, B, NF, NB, FB, NFB: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+ \quad \Gamma_5^+}{\Gamma_0^+, \dots, \Gamma_5^+} \wedge_f E: N}$$

$$\frac{\Gamma_0^-, [[N, N, B, NF, FT, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[F, FB: B]] \quad \Gamma_2^-, [[F: A, B]] \quad \Gamma_3^-, [[F, FB: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+}{\Gamma_0^+, \dots, \Gamma_3^+} \wedge_f E: F}$$

$$\frac{\Gamma_1^-, [[N, F, B, FT, FB, FTB: B]] \quad \Gamma_2^-, [[N, F, FT: A, B]] \quad \Gamma_3^-, [[N, B: A, B]] \quad \Gamma_4^-, [[FT, FTB: A, B]] \quad \Gamma_5^-, [[N, F, B, FT, FB, FTB: A]]}{\frac{\Gamma_1^+, [T, TB: B] \quad \Gamma_2^+, [T: A, B] \quad \Gamma_3^+, [T, TB: A, B] \quad \Gamma_4^+, [T, TB: A, B] \quad \Gamma_5^+, [T, TB: A]}{\Gamma_1^+, \dots, \Gamma_5^+, [T: A \wedge_f B]} \wedge_f I: T}$$

$$\frac{\Gamma_0^-, [[N, N, F, NF, FT, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[B, FB: B]] \quad \Gamma_2^-, [[B: A, B]] \quad \Gamma_3^-, [[B, FB: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+}{\Gamma_0^+, \dots, \Gamma_3^+} \wedge_f E: B}$$

$$\frac{\Gamma_0^-, [[N, N, F, B, FT, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[F, NF, FB, NFB: B]] \quad \Gamma_2^-, [[F, NF: A, B]] \quad \Gamma_3^-, [[NF, NFB: A, B]] \quad \Gamma_4^-, [[F, NF, FB, NFB: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+}{\Gamma_0^+, \dots, \Gamma_4^+} \wedge_f E: NF}$$

$$\frac{\Gamma_1^-, [[N, N, F, NF, FT, NFT: A, B]] \quad \Gamma_2^-, [[N, N, B, NB: A, B]] \quad \Gamma_3^-, [[N, NF, NB, NFT, NFB, A: A, B]] \quad \Gamma_4^-, [[FT, NFT, FTB, A: A, B]]}{\frac{\Gamma_1^+, [T, NT: A, B] \quad \Gamma_2^+, [T, NT, TB, NTB: A, B] \quad \Gamma_3^+, [NT, NTB: A, B] \quad \Gamma_4^+, [T, NT, TB, NTB: A, B]}{\Gamma_1^+, \dots, \Gamma_4^+, [NT: A \wedge_f B]} \wedge_f I: NT}$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, NB, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[F, FT, FB, FTB: B]] \quad \Gamma_2^-, [[F, FT: A, B]] \quad \Gamma_3^-, [[FT, FTB: A, B]] \quad \Gamma_4^-, [[F, FT, FB, FTB: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+}{\Gamma_0^+, \dots, \Gamma_4^+} \wedge_f E: FT}$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, FB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[B, NB, FB, NFB: B]] \quad \Gamma_2^-, [[B, NB: A, B]] \quad \Gamma_3^-, [[NB, NFB: A, B]] \quad \Gamma_4^-, [[B, NB, FB, NFB: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+}{\Gamma_0^+, \dots, \Gamma_4^+} \wedge_f E: NB}$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, NFT, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[FB: B]] \quad \Gamma_2^-, [[FB: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+}{\Gamma_0^+, \Gamma_1^+, \Gamma_2^+} \wedge_f E: FB}$$

$$\frac{\Gamma_1^-, [[B, FB, FTB: B]] \quad \Gamma_2^-, [[B: A, B]] \quad \Gamma_3^-, [[FTB: A, B]] \quad \Gamma_4^-, [[B, FB, FTB: A]]}{\frac{\Gamma_1^+, [TB: B] \quad \Gamma_2^+, [TB: A, B] \quad \Gamma_3^+, [TB: A, B] \quad \Gamma_4^+, [TB: A]}{\Gamma_1^+, \dots, \Gamma_4^+, [TB: A \wedge_f B]} \wedge_f I: TB}$$

$$\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFB, FTB, A: A \wedge_f B]] \quad \Gamma_1^-, [[F, NF, FT, NFB, NFB, FTB, A: B]] \quad \Gamma_2^-, [[F, NF, FT, NFB: A, B]] \quad \Gamma_3^-, [[NF, NFB, NFB, A: A, B]] \quad \Gamma_4^-, [[FT, NFB, FTB, A: A, B]] \quad \Gamma_5^-, [[F, NF, FT, FB, NFB, FTB, A: A]]}{\frac{\Gamma_0^+, [T, NT, TB, NTB: A \wedge_f B] \quad \Gamma_1^+ \quad \Gamma_2^+ \quad \Gamma_3^+ \quad \Gamma_4^+ \quad \Gamma_5^+}{\Gamma_0^+, \dots, \Gamma_5^+} \wedge_f E: NFB}$$



$$\begin{array}{c}
\frac{\Gamma_0^-, [[N, N, F, B, NF, FT, NB, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[N, N, F, B, NF, NB, NFB, A: B]] \quad \Gamma_2^-, [[N, F, B, FT, FB, FTB: A, B]] \quad \Gamma_3^-, [[F, NF, FT, FB, NFB, FTB, A: A, B]] \quad \Gamma_4^-, [[B, NB, FB, NFB, FTB, A: A, B]]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [T, NT, TB, NTB: B] \quad \Gamma_2^+, [T, TB: A, B] \quad \Gamma_3^+, [TB, NTB: A, B] \quad \Gamma_4^+, [T, NT, TB, NTB: A]} \vee_f E: \mathbf{TB} \\
\frac{\Gamma_1^-, [T, NT, TB, NTB: B] \quad \Gamma_2^-, [T, TB: A, B] \quad \Gamma_3^-, [TB, NTB: A, B] \quad \Gamma_4^-, [T, NT, TB, NTB: A]}{\Gamma_1^+, [T, NT, TB, NTB: A \vee_f B]} \vee_f I: \mathbf{TB} \\
\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[NFT: B]] \quad \Gamma_2^-, [[NFT: A, B]] \quad \Gamma_3^-, [[NFT: A]] \\
\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [NT: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [NT: A] \\
\frac{\Gamma_1^+, [NT: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [NT: A]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B]} \vee_f E: \mathbf{NFT} \\
\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[NT, NTB: B]] \quad \Gamma_2^-, [[NF, NFB, A: B]] \quad \Gamma_3^-, [[NB, NFB, A: A, B]] \quad \Gamma_4^-, [[N, NF, NB, NFB, A: A]] \\
\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [NT, NTB: B] \quad \Gamma_2^+, [NF, NFB, A: B] \quad \Gamma_3^+, [NB, NFB, A: A, B] \quad \Gamma_4^+, [N, NF, NB, NFB, A: A] \\
\frac{\Gamma_1^+, [NT, NTB: B] \quad \Gamma_2^+, [NF, NFB, A: B] \quad \Gamma_3^+, [NB, NFB, A: A, B] \quad \Gamma_4^+, [N, NF, NB, NFB, A: A]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B]} \vee_f E: \mathbf{NFB} \\
\Gamma_1^-, [NT, NTB: B] \quad \Gamma_2^-, [NTB: A, B] \quad \Gamma_3^-, [NT, NTB: A] \\
\frac{\Gamma_1^-, [NT, NTB: B] \quad \Gamma_2^-, [NTB: A, B] \quad \Gamma_3^-, [NT, NTB: A]}{\Gamma_1^+, [NT, NTB: B] \quad \Gamma_2^+, [NTB: A, B] \quad \Gamma_3^+, [NT, NTB: A]} \vee_f I: \mathbf{NTB} \\
\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[FT, NFB, A: B]] \quad \Gamma_2^-, [[FT, NTB: A, B]] \quad \Gamma_3^-, [[FT, NFB, FTB, A: A, B]] \quad \Gamma_4^-, [[FTB, A: A, B]] \quad \Gamma_5^-, [[FT, NFB, FTB, A: A]] \\
\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [FT, NFB, A: B] \quad \Gamma_2^+, [FT, NTB: A, B] \quad \Gamma_3^+, [FT, NFB, FTB, A: A, B] \quad \Gamma_4^+, [FTB, A: A, B] \quad \Gamma_5^+, [FT, NFB, FTB, A: A] \\
\frac{\Gamma_1^+, [FT, NFB, A: B] \quad \Gamma_2^+, [FT, NTB: A, B] \quad \Gamma_3^+, [FT, NFB, FTB, A: A, B] \quad \Gamma_4^+, [FTB, A: A, B] \quad \Gamma_5^+, [FT, NFB, FTB, A: A]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B]} \vee_f E: \mathbf{FTB} \\
\Gamma_0^-, [[N, N, F, B, NF, FT, NB, FB, NFB, FTB, A: A \vee_f B]] \quad \Gamma_1^-, [[NFT, A: B]] \quad \Gamma_2^-, [[NFT, A: A, B]] \quad \Gamma_3^-, [[A: A, B]] \quad \Gamma_4^-, [[NFT, A: A]] \\
\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [NFT, A: B] \quad \Gamma_2^+, [NFT, A: A, B] \quad \Gamma_3^+, [A: A, B] \quad \Gamma_4^+, [NFT, A: A] \\
\frac{\Gamma_1^+, [NFT, A: B] \quad \Gamma_2^+, [NFT, A: A, B] \quad \Gamma_3^+, [A: A, B] \quad \Gamma_4^+, [NFT, A: A]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B]} \vee_f E: \mathbf{A} \\
\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B] \quad \Gamma_1^+, [NT, NTB: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [NTB: A, B] \quad \Gamma_4^+, [NT, NTB: A] \\
\frac{\Gamma_1^+, [NT, NTB: B] \quad \Gamma_2^+ \quad \Gamma_3^+, [NTB: A, B] \quad \Gamma_4^+, [NT, NTB: A]}{\Gamma_0^+, [T, NT, TB, NTB: A \vee_f B]} \vee_f I: \mathbf{A}
\end{array}$$

**Definition 22.** A *natural deduction derivation* is defined inductively as follows:

1. Let  $A$  be any formula. Then

$$\frac{[V^-: A]}{[V^+: A]}$$

is a derivation of  $A$  from the assumption  $[V^-: A]$  (an *initial derivation*).

2. If  $D_k$  are derivations of  $\Gamma_k^+, \Delta_k^+$  from the assumptions  $\Gamma_k^-, \hat{\Delta}_k^-$ , and

$$\frac{\left\langle \frac{\Gamma_k^-, [\Delta_k^-]}{\Gamma_k^+, \Delta_k^+} \right\rangle_{k \in K}}{\Pi^+}$$

is an instance of a deduction rule with  $\hat{\Delta}_k^-$  a subsequent of  $\Delta_k^-$ , and all eigenvariable conditions are satisfied, then

$$\frac{\langle D_k \rangle_{k \in K}}{\Pi^+}$$

is a derivation of  $\Pi^+$  from the assumptions  $\bigcup_{k \in K} \Gamma_k^-$ . The formulas in  $\hat{\Delta}_k^-$  which do not occur in  $\bigcup_{k \in K} \Gamma_k^-$  are said to be *discharged* at this inference.

**Theorem 23.** A partial sequent  $\Gamma^+$  can be derived from the assumptions  $\Gamma^-$  in the natural deduction system for Shramko-Wansing logic iff, for every interpretation  $\mathfrak{J}$ , either some formula in  $\Gamma^-$  ( $v \in V^-$ ) evaluates to the truth value  $v$ , or there is a  $w \in V^+$  and a formula in  $\Gamma^+$  that evaluates to  $w$ .

*Proof.* See Theorems 4.7 and 5.4 of Baaz et al. [3] or Theorems 4.2.8 and 4.3.4 of Zach [15].  $\square$

**Corollary 24.**  $\Gamma \models A$  iff there is a natural deduction derivation of  $[V^+:A]$  from  $[V^-:\Gamma]$ .

## 6 Resolution and clause formation rules for Shramko-Wansing logic

The many-valued resolution calculus of Baaz and Fermüller [1] applies to Shramko-Wansing logic. We present the framework here, as well as a language preserving clause translation system for Shramko-Wansing logic.

**Definition 25** (Signed formula). A *signed formula* is an expression of the form  $A^v$ , where  $A$  is a formula and  $v \in V$ . If  $A$  is a propositional variable,  $A^v$  is a *signed atom*.

**Definition 26** (Signed clause). A (signed) *clause*  $C = \{A_1^{v_1}, \dots, A_n^{v_n}\}$  is a finite set of signed atoms (proper clause). The empty clause is denoted by  $\square$ .

An *extended clause* is a finite set of signed formulas.

**Definition 27** (Semantics of clause sets). Let  $\mathfrak{J}$  be an interpretation.  $\mathfrak{J}$  *satisfies* a clause  $C$  iff there is some signed formula  $A^v \in C$ , so that  $\text{val}_{\mathfrak{J}}(A) = v$ .  $\mathfrak{J}$  satisfies a clause set  $\mathcal{C}$  iff it satisfies every clause in  $\mathcal{C}$ .  $\mathcal{C}$  is called *satisfiable* iff some structure satisfies it, and *unsatisfiable* otherwise.

The clause formation rules for connective  $\neg_t$  are given by

$$\begin{array}{l}
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{N}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{N}}\}\}} \neg_t:\mathbf{N} \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{N}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{T}}\}\}} \neg_t:\mathbf{N} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{F}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{B}}\}\}} \neg_t:\mathbf{F} \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{T}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{N}}\}\}} \neg_t:\mathbf{T} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{B}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{F}}\}\}} \neg_t:\mathbf{B} \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{NF}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{TB}}\}\}} \neg_t:\mathbf{NF} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{NT}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NT}}\}\}} \neg_t:\mathbf{NT} \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{FT}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{NB}}\}\}} \neg_t:\mathbf{FT} \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{NB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{FT}}\}\}} \neg_t:\mathbf{NB} \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{\mathbf{FB}}\}\}}{\mathcal{C} \cup \{C \cup \{A^{\mathbf{FB}}\}\}} \neg_t:\mathbf{FB}
\end{array}$$



$$\begin{array}{c}
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{TB}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NF}\}\}} \neg_t:TB \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{NFT}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NTB}\}\}} \neg_t:NFT \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{NFB}\}\}}{\mathcal{C} \cup \{C\}} \neg_t:NFB \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{NTB}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NFT}, A^{FTB}\}\}} \neg_t:NTB \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^{FTB}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NFB}\}\}} \neg_t:FTB \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_t A)^A\}}{\mathcal{C} \cup \{C \cup \{A^A\}\}} \neg_t:A
\end{array}$$

The clause formation rules for connective  $\neg_f$  are given by

$$\begin{array}{c}
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^N\}\}}{\mathcal{C} \cup \{C \cup \{A^N\}\}} \neg_f:N \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^N\}\}}{\mathcal{C} \cup \{C \cup \{A^F\}\}} \neg_f:N \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^F\}\}}{\mathcal{C} \cup \{C \cup \{A^N\}\}} \neg_f:F \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^T\}\}}{\mathcal{C} \cup \{C \cup \{A^B\}\}} \neg_f:T \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^B\}\}}{\mathcal{C} \cup \{C \cup \{A^T\}\}} \neg_f:B \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{NF}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NF}\}\}} \neg_f:NF \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{NT}\}\}}{\mathcal{C} \cup \{C \cup \{A^{FB}\}\}} \neg_f:NT \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{FT}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NB}\}\}} \neg_f:FT \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{NB}\}\}}{\mathcal{C} \cup \{C \cup \{A^{FT}\}\}} \neg_f:NB \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{FB}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NT}\}\}} \neg_f:FB \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{TB}\}\}}{\mathcal{C} \cup \{C \cup \{A^{TB}\}\}} \neg_f:TB \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{NFT}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NFT}, A^{NFB}\}\}} \neg_f:NFT \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{NFB}\}\}}{\mathcal{C} \cup \{C\}} \neg_f:NFB \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{NTB}\}\}}{\mathcal{C} \cup \{C \cup \{A^{FTB}\}\}} \neg_f:NTB \\
\frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^{FTB}\}\}}{\mathcal{C} \cup \{C \cup \{A^{NTB}\}\}} \neg_f:FTB \quad \frac{\mathcal{C} \cup \{C \cup \{(\neg_f A)^A\}}{\mathcal{C} \cup \{C \cup \{A^A\}\}} \neg_f:A
\end{array}$$

The clause formation rules for connective  $\wedge_t$  are given by

$$\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^N\}\}}{\mathcal{C} \cup \{C \cup \{B^N, B^T, B^B, B^{TB}\}, C \cup \{A^N, A^T, B^N, B^T\}, C \cup \{A^N, A^B, B^N, B^B\}, C \cup \{A^N, A^T, A^B, A^{TB}\}\}} \wedge_t:N$$

$$\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^N\}\}}{\mathcal{C} \cup \{C \cup \{B^N, B^T, B^B, B^{NT}, B^{NB}, B^{TB}, B^{NTB}\}, C \cup \{A^N, A^N, A^T, A^{NT}, B^N, B^N, B^T, B^{NT}\}, C \cup \{A^N, A^N, A^B, A^{NB}, B^N, B^N, B^B, B^{NB}\}, C \cup \{A^N, A^{NT}, A^{NB}, A^{NTB}, B^N, B^{NT}, B^{NB}, B^{NTB}\}, C \cup \{A^N, A^N, A^T, A^B, A^{NT}, A^{NB}, A^{TB}, A^{NTB}\}\}} \wedge_t:N$$

$$\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^F\}\}}{\mathcal{C} \cup \{C \cup \{B^N, B^F, B^T, B^B, B^{FT}, B^{FB}, B^{TB}, B^{FTB}\}, C \cup \{A^N, A^F, A^T, A^{FT}, B^N, B^F, B^T, B^{FT}\}, C \cup \{A^N, A^F, A^B, A^{FB}, B^N, B^F, B^B, B^{FB}\}, C \cup \{A^F, A^{FT}, A^{FB}, A^{FTB}, B^F, B^{FT}, B^{FB}, B^{FTB}\}, C \cup \{A^N, A^F, A^T, A^B, A^{FT}, A^{FB}, A^{TB}, A^{FTB}\}\}} \wedge_t:F$$

$$\frac{\mathcal{C} \cup \{C \cup \{(A \wedge_t B)^T\}\}}{\mathcal{C} \cup \{C \cup \{B^T, B^{TB}\}, C \cup \{A^T, B^T\}, C \cup \{A^T, A^{TB}\}\}} \wedge_t:T$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^B\}\}}{\mathcal{C}U\{CU\{B^B, B^{TB}\}, CU\{A^B, B^B\}, CU\{A^B, A^{TB}\}\}} \wedge_t: B$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^{NB}\}\}}{\mathcal{C}U\{CU\{B^B, B^{NB}, B^{TB}, B^{NTB}\}, CU\{A^B, A^{NB}, B^B, B^{NB}\}, CU\{A^{NB}, A^{NTB}, B^{NB}, B^{NTB}\}, CU\{A^B, A^{NB}, A^{TB}, A^{NTB}\}\}} \wedge_t: NB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^{NT}\}\}}{\mathcal{C}U\{CU\{B^T, B^{NT}, B^{TB}, B^{NTB}\}, CU\{A^T, A^{NT}, B^T, B^{NT}\}, CU\{A^{NT}, A^{NTB}, B^{NT}, B^{NTB}\}, CU\{A^T, A^{NT}, A^{TB}, A^{NTB}\}\}} \wedge_t: NT$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^{FT}\}\}}{\mathcal{C}U\{CU\{B^T, B^{FT}, B^{TB}, B^{FTB}\}, CU\{A^T, A^{FT}, B^T, B^{FT}\}, CU\{A^{FT}, A^{FTB}, B^{FT}, B^{FTB}\}, CU\{A^T, A^{FT}, A^{TB}, A^{FTB}\}\}} \wedge_t: FT$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^{NB}\}\}}{\mathcal{C}U\{CU\{B^B, B^{NB}, B^{TB}, B^{NTB}\}, CU\{A^B, A^{NB}, B^B, B^{NB}\}, CU\{A^{NB}, A^{NTB}, B^{NB}, B^{NTB}\}, CU\{A^B, A^{NB}, A^{TB}, A^{NTB}\}\}} \wedge_t: NB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^{FB}\}\}}{\mathcal{C}U\{CU\{B^B, B^{FB}, B^{TB}, B^{FTB}\}, CU\{A^B, A^{FB}, B^B, B^{FB}\}, CU\{A^{FB}, A^{FTB}, B^{FB}, B^{FTB}\}, CU\{A^B, A^{FB}, A^{TB}, A^{FTB}\}\}} \wedge_t: FB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^{TB}\}\}}{\mathcal{C}U\{CU\{B^{TB}\}, CU\{A^{TB}\}\}} \wedge_t: TB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^{NTB}\}\}}{\mathcal{C}U\{CU\{B^T, B^{NT}, B^{FT}, B^{TB}, B^{NTB}, B^{FTB}, B^A\}, CU\{A^T, A^{NT}, A^{FT}, A^{NTB}, B^T, B^{NT}, B^{FT}, B^{NTB}\}, CU\{A^{NT}, A^{NTB}, A^{FTB}, A^A, B^{NT}, B^{NTB}, B^A\}, CU\{A^{FT}, A^{FTB}, A^{FTB}, A^A, B^{FT}, B^{FTB}, B^A\}, CU\{A^T, A^{NT}, A^{FT}, A^{TB}, A^{NTB}, A^{FTB}, A^A\}\}} \wedge_t: NTB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^{FTB}\}\}}{\mathcal{C}U\{CU\{B^T, B^{NB}, B^{FB}, B^{TB}, B^{NTB}, B^{FTB}, B^A\}, CU\{A^T, A^{NB}, A^{FB}, A^{NTB}, B^T, B^{NB}, B^{FB}, B^{NTB}\}, CU\{A^{NB}, A^{NTB}, A^{FTB}, A^A, B^{NB}, B^{NTB}, B^A\}, CU\{A^{FB}, A^{FTB}, A^{FTB}, A^A, B^{FB}, B^{FTB}, B^A\}, CU\{A^T, A^{NB}, A^{FB}, A^{TB}, A^{NTB}, A^{FTB}, A^A\}\}} \wedge_t: FTB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^{NTB}\}\}}{\mathcal{C}U\{CU\{B^{TB}, B^{NTB}\}, CU\{A^{NTB}, B^{NTB}\}, CU\{A^{TB}, A^{NTB}\}\}} \wedge_t: NTB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^{FTB}\}\}}{\mathcal{C}U\{CU\{B^{TB}, B^{FTB}\}, CU\{A^{FTB}, B^{FTB}\}, CU\{A^{TB}, A^{FTB}\}\}} \wedge_t: FTB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_t B)^A\}\}}{\mathcal{C}U\{CU\{B^{TB}, B^{NTB}, B^{FTB}, B^A\}, CU\{A^{NTB}, A^A, B^{NTB}, B^A\}, CU\{A^{FTB}, A^A, B^{FTB}, B^A\}, CU\{A^{TB}, A^{NTB}, A^{FTB}, A^A\}\}} \wedge_t: A$$

The clause formation rules for connective  $\vee_t$  are given by

$$\frac{\mathcal{C}U\{CU\{(A \vee_t B)^N\}\}}{\mathcal{C}U\{CU\{B^N, B^N, B^F, B^{NF}\}, CU\{A^N, A^N, B^N, B^N\}, CU\{A^N, A^F, B^N, B^F\}, CU\{A^N, A^N, A^F, A^{NF}\}\}} \vee_t: N$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_t B)^N\}\}}{\mathcal{C}U\{CU\{B^N, B^{NF}\}, CU\{A^N, B^N\}, CU\{A^N, A^{NF}\}\}} \vee_t: N$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_t B)^F\}\}}{\mathcal{C}U\{CU\{B^F, B^{NF}\}, CU\{A^F, B^F\}, CU\{A^F, A^{NF}\}\}} \vee_t: F$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_t B)^T\}\}}{\mathcal{C}U\{CU\{B^N, B^N, B^F, B^T, B^{NF}, B^{NT}, B^{FT}, B^{NTT}\}, CU\{A^N, A^N, A^T, A^{NT}, B^N, B^N, B^T, B^{NT}\}, CU\{A^N, A^F, A^T, A^{FT}, B^N, B^T, B^{FT}\}, CU\{A^T, A^{NT}, A^{FT}, A^{NTT}, B^T, B^{NT}, B^{FT}, B^{NTT}\}, CU\{A^N, A^N, A^F, A^T, A^{NF}, A^{NT}, A^{FT}, A^{NTT}\}\}} \vee_t: T$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_t B)^B\}\}}{\mathcal{C}U\{CU\{B^N, B^N, B^F, B^B, B^{NF}, B^{NB}, B^{FB}, B^{NFB}\}, CU\{A^N, A^N, A^B, A^{NB}, B^N, B^N, B^B, B^{NB}\}, CU\{A^N, A^F, A^B, A^{FB}, B^N, B^F, B^B, B^{FB}\}, CU\{A^B, A^{NB}, A^{FB}, A^{NFB}, B^B, B^{NB}, B^{FB}, B^{NFB}\}, CU\{A^N, A^N, A^F, A^B, A^{NF}, A^{NB}, A^{FB}, A^{NFB}\}\}} \vee_t: B$$



$$\frac{\mathcal{C}U\{CU\{(A \wedge_f B)^{FT}\}\}}{\mathcal{C}U\{CU\{B^F, B^{FT}, B^{FB}, B^{FTB}\}, CU\{A^F, A^{FT}, B^F, B^{FT}\}, CU\{A^{FT}, A^{FTB}, B^{FT}, B^{FTB}\}, CU\{A^F, A^{FT}, A^{FB}, A^{FTB}\}\}} \wedge_f:FT$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_f B)^{NB}\}\}}{\mathcal{C}U\{CU\{B^B, B^{NB}, B^{FB}, B^{NFB}\}, CU\{A^B, A^{NB}, B^B, B^{NB}\}, CU\{A^{NB}, A^{NFB}, B^{NB}, B^{NFB}\}, CU\{A^B, A^{NB}, A^{FB}, A^{NFB}\}\}} \wedge_f:NB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_f B)^{FB}\}\}}{\mathcal{C}U\{CU\{B^{FB}\}, CU\{A^{FB}\}\}} \wedge_f:FB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_f B)^{TB}\}\}}{\mathcal{C}U\{CU\{B^B, B^{FB}, B^{TB}, B^{FTB}\}, CU\{A^B, A^{TB}, B^B, B^{TB}\}, CU\{A^{TB}, A^{FTB}, B^{TB}, B^{FTB}\}, CU\{A^B, A^{FB}, A^{TB}, A^{FTB}\}\}} \wedge_f:TB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_f B)^{NFB}\}\}}{\mathcal{C}U\{CU\{B^F, B^{NF}, B^{FB}, B^{NFB}, B^{FTB}, B^A\}, CU\{A^F, A^{NF}, A^{FT}, A^{NFT}, B^F, B^{NF}, B^{FT}, B^{NFT}\}, CU\{A^{NF}, A^{NFT}, A^{NFB}, A^A, B^{NF}, B^{NFT}, B^{NFB}, B^A\}, CU\{A^{FT}, A^{NFT}, A^{FTB}, A^A, B^{FT}, B^{FTB}, B^A\}, CU\{A^F, A^{NF}, A^{FT}, A^{FB}, A^{NFT}, A^{NFB}, A^{FTB}, A^A\}\}} \wedge_f:NFB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_f B)^{NFB}\}\}}{\mathcal{C}U\{CU\{B^{FB}, B^{NFB}\}, CU\{A^{NFB}, B^{NFB}\}, CU\{A^{FB}, A^{NFB}\}\}} \wedge_f:NFB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_f B)^{NTB}\}\}}{\mathcal{C}U\{CU\{B^B, B^{NB}, B^{FB}, B^{NTB}, B^{NFB}, B^{FTB}, B^A\}, CU\{A^B, A^{NB}, A^{TB}, A^{NTB}, B^B, B^{NB}, B^{TB}, B^{NTB}\}, CU\{A^{NB}, A^{NFB}, A^{NTB}, A^A, B^{NB}, B^{NFB}, B^{NTB}, B^A\}, CU\{A^{TB}, A^{NTB}, A^{FTB}, A^A, B^{TB}, B^{NTB}, B^{FTB}, B^A\}, CU\{A^B, A^{NB}, A^{FB}, A^{TB}, A^{NFB}, A^{NTB}, A^{FTB}, A^A\}\}} \wedge_f:NTB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_f B)^{FTB}\}\}}{\mathcal{C}U\{CU\{B^{FB}, B^{FTB}\}, CU\{A^{FTB}, B^{FTB}\}, CU\{A^{FB}, A^{FTB}\}\}} \wedge_f:FTB$$

$$\frac{\mathcal{C}U\{CU\{(A \wedge_f B)^A\}\}}{\mathcal{C}U\{CU\{B^{FB}, B^{NFB}, B^{FTB}, B^A\}, CU\{A^{NFB}, A^A, B^{NFB}, B^A\}, CU\{A^{FTB}, A^A, B^{FTB}, B^A\}, CU\{A^{FB}, A^{NFB}, A^{FTB}, A^A\}\}} \wedge_f:A$$

The clause formation rules for connective  $\vee_f$  are given by

$$\frac{\mathcal{C}U\{CU\{(A \vee_f B)^N\}\}}{\mathcal{C}U\{CU\{B^N, B^N, B^T, B^{NT}\}, CU\{A^N, A^N, B^N, B^N\}, CU\{A^N, A^T, B^N, B^T\}, CU\{A^N, A^N, A^T, A^{NT}\}\}} \vee_f:N$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_f B)^N\}\}}{\mathcal{C}U\{CU\{B^N, B^{NT}\}, CU\{A^N, B^N\}, CU\{A^N, A^{NT}\}\}} \vee_f:N$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_f B)^F\}\}}{\mathcal{C}U\{CU\{B^N, B^N, B^F, B^T, B^{NF}, B^{NT}, B^{FT}, B^{NFT}\}, CU\{A^N, A^N, A^F, A^{NF}, B^N, B^N, B^F, B^{NF}\}, CU\{A^N, A^F, A^T, A^{FT}, B^N, B^F, B^T, B^{FT}\}, CU\{A^F, A^{NF}, A^{FT}, A^{NFT}, B^F, B^{NF}, B^{FT}, B^{NFT}\}, CU\{A^N, A^N, A^F, A^T, A^{NF}, A^{NT}, A^{FT}, A^{NFT}\}\}} \vee_f:F$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_f B)^T\}\}}{\mathcal{C}U\{CU\{B^T, B^{NT}\}, CU\{A^T, B^T\}, CU\{A^T, A^{NT}\}\}} \vee_f:T$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_f B)^B\}\}}{\mathcal{C}U\{CU\{B^N, B^N, B^T, B^B, B^{NT}, B^{NB}, B^{TB}, B^{NTB}\}, CU\{A^N, A^N, A^B, A^{NB}, B^N, B^N, B^B, B^{NB}\}, CU\{A^N, A^T, A^B, A^{TB}, B^N, B^T, B^B, B^{TB}\}, CU\{A^B, A^{NB}, A^{TB}, A^{NTB}, B^B, B^{NB}, B^{TB}, B^{NTB}\}, CU\{A^N, A^N, A^T, A^B, A^{NT}, A^{NB}, A^{TB}, A^{NTB}\}\}} \vee_f:B$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_f B)^{NF}\}\}}{\mathcal{C}U\{CU\{B^N, B^{NF}, B^{NT}, B^{NFT}\}, CU\{A^N, A^{NF}, B^N, B^{NF}\}, CU\{A^{NF}, A^{NFT}, B^{NF}, B^{NFT}\}, CU\{A^N, A^{NF}, A^{NT}, A^{NFT}\}\}} \vee_f:NF$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_f B)^{NT}\}\}}{\mathcal{C}U\{CU\{B^{NT}\}, CU\{A^{NT}\}\}} \vee_f:NT$$

$$\frac{\mathcal{C}U\{CU\{(A \vee_f B)^{FT}\}\}}{\mathcal{C}U\{CU\{B^T, B^{NT}, B^{FT}, B^{NFT}\}, CU\{A^T, A^{FT}, B^T, B^{FT}\}, CU\{A^{FT}, A^{NFT}, B^{FT}, B^{NFT}\}, CU\{A^T, A^{NT}, A^{FT}, A^{NFT}\}\}} \vee_f:FT$$

$$\begin{array}{c}
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{NB}\}\}}{\mathcal{C} \cup \{C \cup \{B^N, B^{NT}, B^{NB}, B^{NTB}\}, C \cup \{A^N, A^{NB}, B^N, B^{NB}\}, C \cup \{A^{NB}, A^{NTB}, B^{NB}, B^{NTB}\}, C \cup \{A^N, A^{NT}, A^{NB}, A^{NTB}\}\}} \forall_f:NB \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{TB}\}\}}{\mathcal{C} \cup \{C \cup \{B^T, B^{NT}, B^{TB}, B^{NTB}\}, C \cup \{A^T, A^{TB}, B^T, B^{TB}\}, C \cup \{A^{TB}, A^{NTB}, B^{TB}, B^{NTB}\}, C \cup \{A^T, A^{NT}, A^{TB}, A^{NTB}\}\}} \forall_f:TB \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{NFT}\}\}}{\mathcal{C} \cup \{C \cup \{B^{NT}, B^{NFT}\}, C \cup \{A^{NFT}, B^{NFT}\}, C \cup \{A^{NT}, A^{NFT}\}\}} \forall_f:NFT \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^{NTB}\}\}}{\mathcal{C} \cup \{C \cup \{B^{NT}, B^{NTB}\}, C \cup \{A^{NTB}, B^{NTB}\}, C \cup \{A^{NT}, A^{NTB}\}\}} \forall_f:NTB \\
\frac{\mathcal{C} \cup \{C \cup \{(A \vee_f B)^A\}\}}{\mathcal{C} \cup \{C \cup \{B^{NT}, B^{NFT}, B^{NTB}, BA\}, C \cup \{A^{NFT}, AA, B^{NFT}, BA\}, C \cup \{A^{NTB}, AA, B^{NTB}, BA\}, C \cup \{A^{NT}, A^{NFT}, A^{NTB}, AA\}\}} \forall_f:A
\end{array}$$

**Theorem 28.** Let  $T(\mathcal{C})$  be the result of exhaustively applying the translation rules to a clause set  $\mathcal{C}$ . Then  $T(\mathcal{C})$  is a set of proper clauses, i.e.,  $T(\mathcal{C})$  contains only signed atoms (all connectives are eliminated). Furthermore,  $T(\mathcal{C})$  is satisfiable iff  $\mathcal{C}$  is.

**Proposition 29.** Let  $A$  be a sentence and  $\Delta$  be a set of sentences. Then

1.  $\models A$  iff  $\{\{A^v \mid v \in V^-\}\}$  is unsatisfiable.
2.  $\Delta \models A$  iff

$$\{\{B^w \mid w \in V^+\} \mid B \in \Delta\} \cup \{\{A^v \mid v \in V^-\}\}$$

is unsatisfiable.

**Definition 30.** A clause  $R$  is a *resolvent* of clauses  $C_1, C_2$  if  $R = (C_1 \setminus \{A^{v_1}\}) \cup (C_2 \setminus \{A^{v_2}\})$  where  $v_1 \neq v_2$

**Definition 31.** A *resolution refutation* of a clause set  $\mathcal{C}$  is a sequence of clauses  $C_1, \dots, C_n$  so that for every  $i$ ,  $C_i \in \mathcal{C}$  or  $C_i$  is a resolvent of  $C_j, C_k$  with  $j, k < i$ , and  $C_n = \emptyset$ .

**Theorem 32.** A clause set  $\mathcal{C}$  is unsatisfiable iff it has a resolution refutation.

*Proof.* See Theorems 3.14 and 3.19 of Baaz and Fermüller [1] or Theorems 2.5.5 and 2.5.8 of Zach [15].  $\square$

**Corollary 33.**  $\Delta \models A$  iff

$$T(\{\{B^w \mid w \in V^+\} \mid B \in \Delta\} \cup \{\{A^v \mid v \in V^-\}\})$$

has a resolution refutation.

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## A shramko-wansing.lgc – specification of Shramko-Wansing logic

```

logic "Shramko-Wansing".
truth_values { e , n , f , t , b , nf , nt , ft , nb , fb , tb ,
              nft , nfb , ntb , ftb , a }.
designated_truth_values { t , nt , tb , ntb }.
operator(negt/1, mapping {
    (e ) : e ,
    (n ) : t ,
    (f ) : b ,
    (t ) : n ,
    (b ) : f ,
    (nf ) : tb ,
    (nt ) : nt ,
    (nb ) : ft ,
    (ft ) : nb ,
    (fb ) : fb ,
    (tb ) : nf ,
    (nft ) : ntb ,
    (nfb ) : ftb ,
    (ntb ) : nft ,
    (ftb ) : ntb ,
    (a ) : a
}
).
operator(negf/1, mapping {
    (e ) : e ,
    (n ) : f ,
    (f ) : n ,
    (t ) : b ,
    (b ) : t ,
    (nf ) : nf ,
    (nt ) : fb ,
    (nb ) : ft ,
    (ft ) : nb ,
    (fb ) : nt ,
    (tb ) : tb ,
    (nft ) : nft ,
    (nfb ) : nft ,
    (ntb ) : ftb ,

```

```

        (ftb) : ntb,
        (a ) : a
    }
).
ordering(truth, "nf < { f < { fb < {ftb < tb, b < tb}, ft < {ftb, t
    < tb}, e < {b,t}}, nfb < { fb, a < {ftb, ntb}, nb < {b, ntb
    }}, nft < { ft, a, nt < {t, ntb < tb}}, n < {e, nb, nt}").
ordering(falsity, "nt < { n < { nf < {f < fb, nfb < fb}, e < {f, b
    < fb}, nb < {nfb, b}}, nft < { nf, ft < {f, ftb}, a < {nfb,
    ftb}}, t < { e, ft, tb < {b, ftb < fb}}, ntb < {nb, a, tb}").
operator(andt /2, inf(truth)).
operator(ort /2, sup(truth)).
operator(andf /2, inf(falsity)).
operator(orf /2, sup(falsity)).

```

## B shramko-wansing.cfg – L<sup>A</sup>T<sub>E</sub>X translations

```

texName(n, "\\boldsymbol{N}").
texName(f, "\\boldsymbol{F}").
texName(t, "\\boldsymbol{T}").
texName(b, "\\boldsymbol{B}").
texName(e, "\\mathbf{N}").
texName(a, "\\mathbf{A}").
texName(nf, "\\mathbf{NF}").
texName(nt, "\\mathbf{NT}").
texName(nb, "\\mathbf{NB}").
texName(ft, "\\mathbf{FT}").
texName(fb, "\\mathbf{FB}").
texName(tb, "\\mathbf{TB}").
texName(nft, "\\mathbf{NFT}").
texName(nfb, "\\mathbf{NFB}").
texName(ntb, "\\mathbf{NTB}").
texName(ftb, "\\mathbf{FTB}").

texName(andt, "\\land_t" ).
texName(ort, "\\lor_t" ).
texName(negt, "\\lnot_t" ).
texInfix(andt).
texInfix(ort).
texPrefix(negt).
texName(andf, "\\land_f" ).
texName(orf, "\\lor_f" ).
texName(negf, "\\lnot_f" ).
texInfix(andf).
texInfix(orf).
texPrefix(negf).

texName(forall, "\\forall_t").
texName(exists, "\\exists_t").

texExtra("ShortName", "\\mathbf{SIXTEEN}_{2,\\mathit{des}}").
texExtra("Semantics", "The connectives  $\\land_t$  and  $\\lor_t$  of
     $\\ShortName$  are defined as the inf and sup of the  $\\le_t$ 
    ordering given in \\cref{fig}. The connectives  $\\land_f$  and  $\\lor_f$  correspond to inf and sup of  $\\le_f$ . For simplicity,
    we leave out the  $\\le_i$  ordering and other operators and

```



```
inversions defined in \cite{ml}.\begin{figure}\centering\includegraphics[width=.5\textwidth]{shramko-wansing.png}\caption{The  $\leq_t$  and  $\leq_f$  orderings}\label{fig}\end{figure}).  
texExtra("Preamble", "\usepackage{graphicx,amsbsy}").  
texExtra("Link", "https://logic.at/multlog/shramko-wansing.pdf").
```