

On Some Abstract
Termination Criteria

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• Termination by weak termination + ...
(globally / locally)

$$P?(\rightarrow) \Rightarrow SN(\rightarrow)$$

ARS_s
(+TRS_s)

• Termination Criteria based on
non-overlapping and/or non-erasing properties

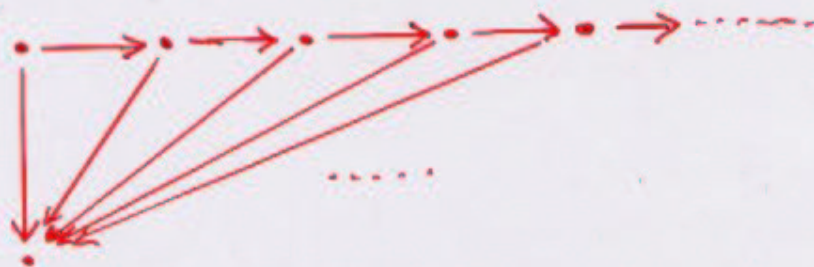
TRS_s

ARS: (A, \rightarrow) with $\rightarrow \subseteq A \times A$

$WN(\rightarrow) \not\Rightarrow SN(\rightarrow)$

$\left. \begin{array}{l} WN(\rightarrow) \\ WCR(\rightarrow) \end{array} \right\} \not\Rightarrow SN(\rightarrow)$

$\left. \begin{array}{l} WN(\rightarrow) \\ CR(\rightarrow) \end{array} \right\} \not\Rightarrow SN(\rightarrow)$



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Theorem [Klop'80]

Given: $(A, \rightarrow), a \in A$.

If ① $\exists b \in NF(\rightarrow) : a \xrightarrow{*} b$

WN(a)

② $WCR(G(a))$

(cc. conf. below)

③ $\underbrace{\max\{k \mid a \xrightarrow{k} b\}}_{=: n}$ exists

boundedness

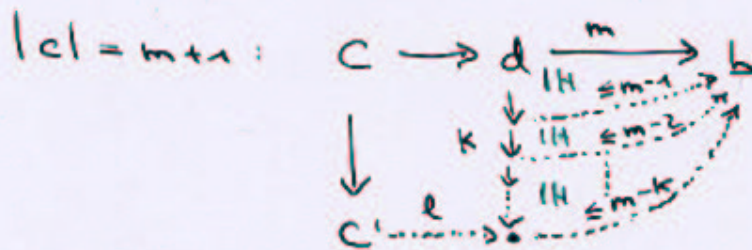
then: $SN(a)$

Proof: $B := \{c \mid a \xrightarrow{*} c \xrightarrow{*} b\}$

for $c \in B$: $|c| := \max\{k \mid c \xrightarrow{k} b\}$

(a) $\forall c, c' \in B : c \rightarrow c' \Rightarrow |c| > |c'|$

(b) $\left. \begin{array}{l} c \in B \\ c \rightarrow c' \end{array} \right\} \Rightarrow c' \in B$ by ind. on $|c|$



$$1+m \geq 1+l+m-k$$

$$\Rightarrow k \geq l$$

\Rightarrow all reductions from a have length at most n

$\Rightarrow SN(a) \wedge CR(G(a))$

Theorem [Klop'80]

Given: (A, \rightarrow) , $a \in A$.

If ① $\exists b \in NF(\rightarrow): a \xrightarrow{*} b$

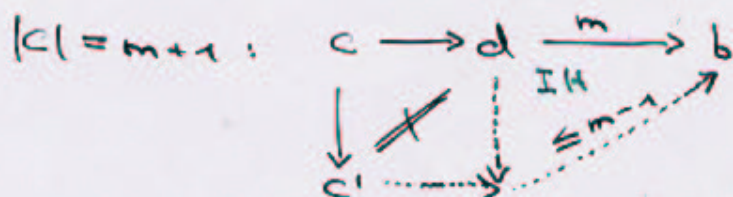
④ $WCR^{-1}(G(a))$

then: $SN(a)$

WCR⁻¹:



Proof: $\left. \begin{array}{l} c \in B \\ c \rightarrow c' \end{array} \right\} \Rightarrow \begin{array}{l} c' \in B \\ |c'| = |c| - 1 \end{array}$ by ind. on $|c|$



Theorem [Toyama's]

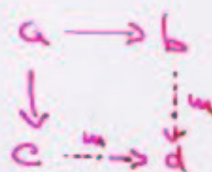
Given: $(A, \rightarrow), a \in A$.

If ① $\exists b \in NF(\rightarrow): a \xrightarrow{*} b$

② $BWCR(G(a))$

then: $SN(a)$

$BWCR(D): \forall a, b, c \in D \exists d \in D \exists m \in \mathbb{N}:$



Proof:

Theorem

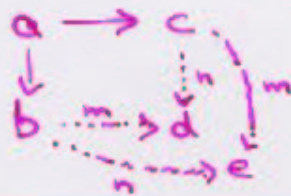
Given: (A, \rightarrow) , $a \in A$.

If ① $\exists b \in NF(a) : a \xrightarrow{*} b$

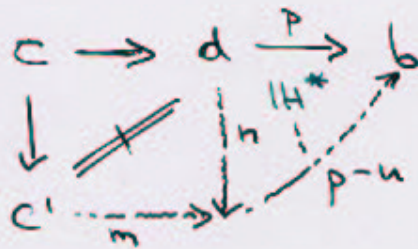
then: $SN(a)$

⑥ BSCR($G(a)$)

"BSCR"(D): $\forall a, b, c \in D \exists d, e \in D \exists m, u \in \mathbb{N}$:



Proof: $\left. \begin{matrix} c \in B \\ c \rightarrow c' \end{matrix} \right\} \Rightarrow \left. \begin{matrix} c' \in B \\ c' \xrightarrow{H^{-1}} b \end{matrix} \right\}$



$$\begin{aligned} 1 + m + p - n &\leq 1 + p \\ m &\leq n \end{aligned}$$

Symmetrically:

$$\begin{aligned} 1 + u + p - m &\leq 1 + p \\ n &\leq m \end{aligned}$$

hence: $m = u$

$$\Rightarrow c' \xrightarrow{p} b = |c| - 1$$

Theorem [Klop'80]

Given: (A, \rightarrow) , $a \in A$

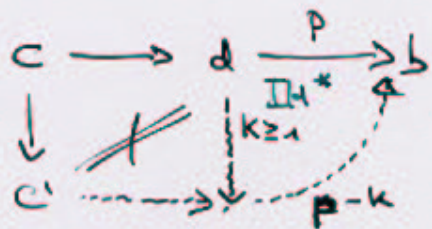
If ① $\exists b \in NF(\rightarrow): a \xrightarrow{*} b$

② $\underbrace{\max\{k \mid a \xrightarrow{k} b\}}_{=: n}$ exists

$$\textcircled{7} \text{SCR}^{\geq 1}(G(a))$$

$$\text{SCR}^{\geq 1}(D) :$$

Proof: $\left. \begin{array}{l} c \in B \\ c \rightarrow c' \end{array} \right\} \Rightarrow \begin{array}{l} c' \in B \\ c' \xrightarrow{|c|-1} b \end{array}$



$$\begin{aligned} p - k + 2 &\leq p + 1 \\ 1 &\leq k \end{aligned}$$

Symmetrically:

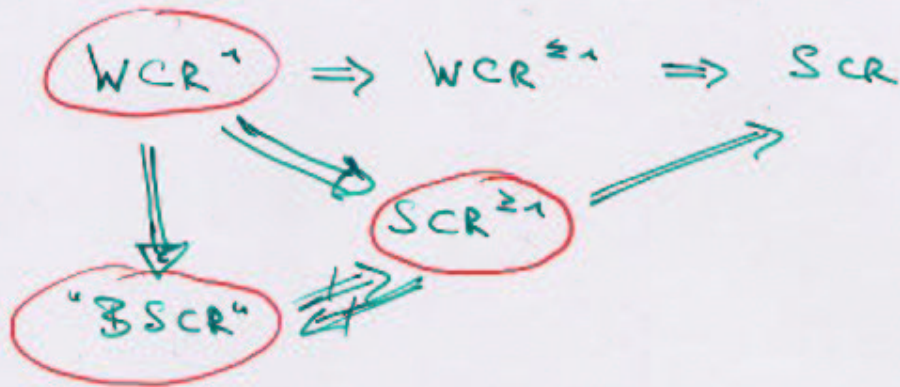
$$\begin{aligned} p - 1 + k + 1 &\leq p + 1 \\ k &\leq 1 \end{aligned}$$

$$\Rightarrow k = 1$$

$$c' \xrightarrow{p} b = |c| - 1$$

$$WN \Rightarrow [WCR^1 \Leftrightarrow BWCR \Leftrightarrow BSCR \Leftrightarrow SCR^{2-1}]$$

however:



→ maybe useful for refined analysis
if $WN(\rightarrow)$ does not hold globally

Non-overlapping (NO) /

Non-erasing (NE)

⊗ $NO \wedge NE \Rightarrow [WN \Leftrightarrow SN]$
 $NO \wedge WIN \Rightarrow SN$

[192, '95]
— —

Theorem $R = R_1 \cup R_2$

[Dankowitz '95]

- ① R_1, R_2 SN
 - ② R_1 preserves R_2 -NFs
 - ③ $R_1 \cup R_2$ NO
 - ④ $R_1 \cup R_2$ NE
- } $\Rightarrow R$ SN

necessary?

Non-overlapping (No) /

Non-erasing

* $NO \wedge NE \Rightarrow [WN \Leftrightarrow SN]$
 $NO \wedge WIN \Rightarrow SN$

[92, '95]

Theorem $R = R_1 \cup R_2$

- (1) R_1, R_2 WIN
(2) R_1 preserves R_2 -NFs
(3) $R_1 \cup R_2$ No } $\Rightarrow R$ SN

Proof: Suppose: $\infty_R(s)$

Claim: $\exists s'$: $\begin{matrix} \infty_R \\ s' \end{matrix} \xrightarrow{i(R_2)} \begin{matrix} \infty_R \\ s' \end{matrix} \in NF(R_2)$

(2) (1) $\Rightarrow \exists s''$: $\begin{matrix} \infty_R \\ s' \end{matrix} \xrightarrow{i(R_1)} \begin{matrix} \infty_R \\ s'' \end{matrix} \in \underbrace{NF(R_1) \cap NF(R_2)}_{=NF(R)}$

(2) $\Rightarrow \begin{matrix} \infty_R \\ s' \end{matrix} \xrightarrow{i(R)} \begin{matrix} \infty_R \\ s'' \end{matrix}$

$\Rightarrow WIN(s', R)$

(3) (*) $\Rightarrow SN(s', R)$

$\Rightarrow \Downarrow$

Non-Erasing Transformations

$$R \left\{ \begin{array}{l} \dots \\ p(s(x)) = x \\ \text{fact}(\emptyset) = S(\emptyset) \\ \text{fact}(S(x)) = S(x) \cdot \text{fact}(p(S(x))) \\ \dots \end{array} \right.$$

Simply!

$$SN(R) \Rightarrow SN(R')$$

↕

The transformation $R \rightarrow R'$ is sound (wrt. SN) if R is non-overlapping and R' is obtained from R

by replacing $l \rightarrow r \in R$ with $l \rightarrow r' \in R'$

where: $r \xrightarrow{*} r'$
 \swarrow
 R

using only non-erasing rules

This can be generalized
(a little bit)

Conclusion

abstract relationships / criteria

+

structural knowledge

may be

very useful

for termination proofs